

1 Basic Theories

1.1 Boolean Theory

Operators Some boolean operators are supported by L^AT_EX, but they have names suggesting shape rather than content, e.g., `a \rightarrow b`. It would be nice if they were given informative, short names without clashing with existing L^AT_EX commands.

$a \Rightarrow b$	<code>a \imp b</code>
$a \Rightarrow\Rightarrow b$	<code>a \Imp b</code>
$a \Leftarrow b$	<code>a \pmi b</code> or <code>a \impby b</code>
$a \Leftarrow\Leftarrow b$	<code>a \Pmi b</code> or <code>a \Impby b</code>
$a = b$	<code>a \Eq b</code>
<code>if c then x else y fi</code>	<code>\cond{c}{x}{y}</code>

Other boolean symbols (`=`, `\bot` for \perp , `\lor` for \vee , etc.) have reasonable names in L^AT_EX, and I will not show them.

Two more variants of if-then-else-fi:

- `\condb{c}{x}{y}`:
`if c then x`
`else y fi`
- `\condbb{c}{x}{y}`:
`if c then`
`x`
`else`
`y`
`fi`

Proof format Using the `align*` environment provided by $\mathcal{A}\mathcal{M}\mathcal{S}$ -L^AT_EX (package name `amsmath`), a calculational proof with hints can be typeset easily. The first proof in the textbook:

$a \wedge b \Rightarrow c$	Material Implication
$= \neg(a \wedge b) \vee c$	Duality
$= \neg a \vee \neg b \vee c$	Material Implication
$= a \Rightarrow \neg b \vee c$	Material Implication
$= a \Rightarrow (b \Rightarrow c)$	

Its code:

```
\begin{align*}
&\&\text{Blank } a \ \&\text{et } b \ \&\text{imp } c && \&\& \text{\text{Material Implication}} \ \&\& \\
&\&\text{Eq } \neg(a \ \&\text{et } b) \ \&\text{vel } c && \&\& \text{\text{Duality}} \ \&\& \\
&\&\text{Eq } \neg a \ \&\text{vel } \neg b \ \&\text{vel } c && \&\& \text{\text{Material Implication}} \ \&\& \\
&\&\text{Eq } \neg a \ \&\text{vel } \neg b \ \&\text{vel } c && \&\& \text{\text{Material Implication}} \ \&\& \\
&\&\text{Eq } a \ \&\text{imp } (\neg b \ \&\text{vel } c) && \&\& \text{\text{Material Implication}} \ \&\& \\
&\&\text{Eq } a \ \&\text{imp } (b \ \&\text{imp } c) && \&\& \text{\text{Material Implication}} \ \&\&
\end{align*}
```

```

&\Eq a \imp \neg b \vel c      && \text{Material Implication} \\
&\Eq a \imp (b \imp c)
\end{align*}

```

The command `\Blank` is a blank relation symbol I invented; it is necessary in that position to keep `align*` happy. Its definition is simply: `\mathrel{}`.

2 Basic Data Structures

Bunch Theory

A, B	<code>A, B</code>
$A'B$	<code>A'B</code>
$null$	<code>\nul</code>
$\#A$	<code>\card A</code>
$0, \dots, 10$	<code>0 \bto 10</code>
$nat, xnat$	<code>\nat, \xnat</code>
$int, xint$	<code>\int, \xint</code>
$rat, xrat$	<code>\rat, \xrat</code>

String Theory

nil	<code>\nil</code>
n^*S	<code>n^*S</code>
$*S$	<code>\{ }^*S</code>
$0; \dots, 10$	<code>0 \sto 10</code>

List Theory

L^+M	<code>L^+M</code>
$n \rightarrow i \mid L$	<code>n \to i \ow L</code>
Ln	<code>L \ap n</code>

3 Function Theory

The `\fun` and `\fn` commands produce functions; `\fun` requires a domain and `\fn` omits the domain.

$\langle x: nat \rightarrow x + 1 \rangle$	<code>\fun{x}{\nat}{x+1}</code>
$\langle x \rightarrow x + 1 \rangle$	<code>\fn{x}{x+1}</code>

The `\bind` and `\bnd` commands help you produce quantified expressions. They just have the quantifier missing, and you just put it back. `\bind` requires a domain and `\bnd` omits the domain. Some examples:

$\forall x \cdot x = x$	<code>\forallall\bnd{x}{x=x}</code>
$\Sigma i: 0, \dots, 10 \cdot i^2$	<code>\Sigma\bind{i}{0 \bto 10}{i^2}</code>
$\S x: nat \cdot x/2 : nat$	<code>\S\bind{x}{\nat}{x/2:\nat}</code>

Two quantifiers are not already available in \LaTeX : MAX and MIN . I have defined them as \backslashMAX and \backslashMIN , respectively.

Both application and composition are \backslashap . You can think of it as standing for “apposition”. Selective union is \backslashow , standing for “otherwise”. You have seen them in List Theory. More examples:

$MAX v: x \cdot n$	$\text{\backslashMAX}\text{\backslashbind}\{v\}\{x\}\{n\}$
$MIN v: x \cdot n$	$\text{\backslashMIN}\text{\backslashbind}\{v\}\{x\}\{n\}$
$f g$	$f \text{\backslashow} g$
$h f x g y$	$h \text{\backslashap} f \text{\backslashap} x \text{\backslashap} g \text{\backslashap} y$

4 Program Theory

ok	\backslashok
$S . R$	$S \text{\backslashdc} R$
$x := e$	$x \text{\backslashget} e$

5 Programming Language

Two forms of while-do-od:

- $\text{\backslashwhile}\{c\}\{P\}$:

while c do P od
--
- $\text{\backslashwhileb}\{c\}\{P\}$:

while c
do P od