# Alternating Direction Implicit methods for Black-Scholes Equations 

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## 1 Introduction

In this project we consider numerically solving two- and three-dimensional Black-Scholes parabolic partial differential equations (PDEs) using finite difference (FD) methods. The main purpose is to contrast the performance of Alternating Direction Implicit (ADI) methods and Crank-Nicolson (CN) methods in multiple dimensions as measured by time complexity. A secondary purpose is to present an efficient implementation of the ADI method.

We will discuss the problem formulation in Section 2, describe the algorithms in Section 3, present numerical results of our algorithms in Section 4, discuss the order of convergence and the runtime in Section 5 , and make our conclusions in Section 6.

## 2 Problem Formulation

Although the focus of the project is on multidimensional Black-Scholes PDEs, we will start by introducing the one-dimensional Black-Scholes PDE.

$$
\begin{equation*}
V(\tau, S)_{\tau}=\mathcal{L}(V)=\frac{1}{2} \sigma^{2} S^{2} V_{S S}+(r-q) S V_{S}-r V \tag{1}
\end{equation*}
$$

where $V$ is the price of the option, $S$ is the price of the underlying asset, $\tau$ is backwards time, $\sigma$ is the volatility, $r$ is the interest rate, and $q$ is the dividend rate. We start at some time $\tau=0$ until the end time $\tau=T$, given some initial conditions $V^{*}(S)=V(\tau=0, S)$ known as the payoff. The payoffs most useful to us will be the call and put, given by

$$
\begin{align*}
V_{\text {call }}^{*}(S) & =\max (S-K, 0)  \tag{2}\\
V_{\text {put }}^{*}(S) & =\max (K-S, 0) \tag{3}
\end{align*}
$$

The solution of Equation (1) at $\tau=T$ gives the Black-Scholes price of the option.
Next, we will consider the $d$-dimensional European options with no dividend, which are given by

$$
\begin{equation*}
V\left(\tau, S_{i}\right)_{\tau}=\mathcal{L}^{(d)}(V)=\frac{1}{2} \sum_{i, j=1}^{d} \rho_{i, j} \sigma_{i} \sigma_{j} V_{S_{i}} V_{S_{j}}+\sum_{i=1}^{d} r S_{i} V_{S_{i}}-r V, \tag{4}
\end{equation*}
$$

where $S_{i}$ are the stock prices, $\sigma_{i}, \rho_{i, j}$ denote volatility of stock prices and correlation between stock prices (with $\rho_{j, j}=1$ ).

For the problem to be well-posed, we need to have initial conditions, which correspond to the type of payoff at the option's maturity. The test problems we use to illustrate our methods are geometric put and call which are convenient because they can be simplified to one-dimensional option pricing problems that either have exact solutions or can be solved with efficient algorithms [6]. The payoff functions for $d$-dimensional geometric average call and put are given by:

$$
\begin{align*}
V_{\text {call }}^{*}\left(S_{i}\right) & =\max \left(\left(\Pi S_{i}\right)^{1 / d}-K, 0\right)  \tag{5}\\
V_{\text {put }}^{*}\left(S_{i}\right) & =\max \left(K-\left(\Pi S_{i}\right)^{1 / d}, 0\right) \tag{6}
\end{align*}
$$

Note that when $d=1$, the above definitions simplify to the one-dimensional call and put payoffs.
Geometric average option pricing problems are special because that they are one of the few option pricing problems that have exact solutions no matter what number of underlying assets (dimensions) we have. Therefore it is a popular and convenient test problem for many algorithms designed for high-dimensional PDE problems.

With a geometric average call/put initial condition, the problem (4) can be transformed to a one-dimensional option pricing problem

$$
\begin{equation*}
V(\tau, S)_{\tau}=\frac{\sigma^{2} S^{2}}{2} V_{S S}+(r-q) S-r S \tag{7}
\end{equation*}
$$

where

$$
\begin{align*}
\sigma & =\frac{1}{n} \sqrt{\sum_{i, j=1}^{d} \rho_{i, j} \sigma_{i} \sigma_{j}}  \tag{8}\\
q & =\frac{1}{2 n} \sum_{i=1}^{d} \sigma_{i}^{2}-\frac{1}{2} \sigma^{2} \tag{9}
\end{align*}
$$

and we can either use the Black-Scholes formula to compute the solution in the case of European options or use known algorithms such as the penalty method [6] in the case of American options.

## 3 PDE discretization

Consider the stock prices $S_{i}$ which belong to the semi-infinite domain $[0, \infty)$. For convenience, we will use $x, y$, and $z$ to represent the (up to) three stock prices $S_{1}, S_{2}$, and $S_{3}$ in our problem. For computation purposes, we truncate the domain to $\left[0, x_{\max }\right]$. Here, we pick $S_{1_{\max }}=10 \mathrm{~K}$. As is convention, let $\tau^{k}$, $k=0, \ldots, N_{\tau}$ denote a (possibly non-uniform) partition of the time interval $[0, T]$ in ascending order, and $\Delta \tau^{k}=\tau^{k}-\tau^{k-1}$. Let $x_{0}, x_{1}, \ldots, x_{N_{x}}$ denote a (possibly non-uniform) partition of the interval [ $\left.0, x_{\text {max }}\right]$, and $y_{i}, z_{i}$ defined likewise. In our numerical experiments, we use a uniform timestepping scheme for European options and a variable timestepping scheme presented in [6] for American options. For spatial discretization, we use a nonuniform grid in [1] with a more dense grid (more gridpoints are allocated near the region of interest) for two dimensions and a slightly less dense grid for three dimensions. Finally, let $v_{h, i, j}^{k}$ denote a computed solution of $u\left(\tau^{k}, x_{h}, y_{i}, z_{j}\right)$ and $v_{i, j}^{k}$ denote a computed solution of $u\left(\tau^{k}, x_{i}, y_{j}\right)$. For the purposes of algorithm analysis, we assume that $N_{x}=N_{y}=N_{z}$ which we denote as $N$.

### 3.1 Space discretization

We use standard second-order FD for the spatial derivatives, and denote the $d$-dimensional discretized operator as a square $(N+1)^{d}$ matrix $L^{(d)}$ which encompasses the boundary conditions.
Note that $L^{(d)}$ can be decomposed into $L^{(d)}=L_{0}+L_{1}+L_{2}+\cdots+L_{d}$, where $L_{j}(j>0)$ are the matrices resulting from just the FD with respect to the $j$-th spatial variable, sharing equally the $-r u$ (no-derivative) term and $L_{0}$ is the matrix resulting from only the cross-derivative discretizations.

### 3.2 Crank-Nicolson Method

Let $\theta \in[0,1]$. To solve the solution at time $\tau^{k}$ given the solution at $\tau^{k-1}$, the $\theta$-timestepping discretization solves the linear system

$$
\begin{equation*}
\left(I-\theta \Delta \tau^{k} L^{(d)}\right) v^{k}=\left(I+(1-\theta) \Delta \tau^{k} L^{(d)}\right) v^{k-1} \tag{10}
\end{equation*}
$$

where $I$ denotes the identity matrix of size $(N+1)^{d}$. When $\theta=1$, we have the Fully Implicit/Backward Euler (BE) method, and when $\theta=1 / 2$, we have the Crank-Nicolson (CN) method. Although CN is second order convergent, it may produce spurious oscillations when applied to non-smooth initial conditions. On the other hand, BE is more stable, but only first order convergent. To obtain both the accuracy of CN and the stability of BE, we use the Rannacher smoothing [9] technique which applies BE to the first two timesteps partitioned into four smaller timesteps, and in the subsequent timesteps uses CN and the usual stepsize. This results in smooth data for the Crank-Nicolson method, and the use of a few iterations of the first-order BE method does not affect the overall second order convergence of the CN method.

Solving Equation (10) requires the solution of a $N^{d} \times N^{d}$ sparse, nonsymmetric linear system with bandwidth $\mathcal{O}\left(N^{d-1}\right)$ each iteration. Recall that banded LU factorization requires $\mathcal{O}(n l u)$ where $n$ is the size of the matrix, $l$ is the lower bandwidth and $u$ is the upper bandwidth. Banded LU will take $\mathcal{O}\left(N^{d(d-1)^{2}}\right)$; in both the two-dimensional or three-dimensional case MATLAB's backslash operator eventually resorts to using a general sparse LU solver with full pivoting [3] which is designed to reduce fill-in.

### 3.3 ADI and Craig-Sneyd

ADI methods such as $[2,4,8]$ minimize the bandwidth of the matrices being solved; If we solve a tridiagonal system at each iteration, then the cost reduces from $\mathcal{O}\left(N^{d(d-1)^{2}}\right)$ to $\mathcal{O}\left(N^{d}\right)$, which is optimal since we have $\mathcal{O}\left(N^{d}\right)$ unknowns.

The original ADI method [5] does not consider mixed derivatives (i.e. $v_{x y}$ ). The Craig-Sneyd [2] and more general ADI methods [8] address this and we will focus on the Modified Craig-Sneyd (MCS) method in this study.

In the MCS method, we solve the following set of equations/update the following vectors for a $d$-dimensional problem, where the restrictions $1 / 3 \leq \theta, \sigma=\theta, \mu=1 / 2-\theta$ are imposed upon the parameters to ensure second order convergence.

Note that for solving the systems of equations, we re-order the entries of the right-hand-side vector so that the matrix $I-\theta \Delta \tau^{k} L_{j}$ is always tridiagonal. This ensures $\mathcal{O}\left(N^{d}\right)$ performance per iteration.

### 3.4 Matricized Craig-Sneyd

This only applies to the two-dimensional case, although a similar process can be used for the threedimensional case.

Instead of writing the previous algorithm in vector form, we can take advantage of the special structure of the matrices $L_{1}$ and $L_{2}$. Note that in the default ordering,

$$
\begin{align*}
& L_{0}=\rho D_{x} \otimes D_{y}  \tag{11}\\
& L_{1}=L_{x} \otimes I  \tag{12}\\
& L_{2}=I \otimes L_{y} \tag{13}
\end{align*}
$$

```
Algorithm 1 MCS ADI timestepping in \(d\) dimensions
    Take an explicit step:
    \(Y_{0}=v^{k-1}+\Delta \tau^{k}\left(L v^{k-1}\right)\)
    Take \(d\) implicit corrector steps, one in each dimension:
    for \(j=1, \ldots, d\) do
        compute \(g=Y_{j-1}-\theta \Delta \tau^{k} L_{j} v^{k-1}\)
        reorder the entries in \(g\) such that \(L_{j}\) is tridiagonal
        Solve the linear system \(\left(I-\theta \Delta \tau^{k} L_{j}\right) Y_{j}=Y_{j-1}-\theta \Delta \tau^{k} L_{j} v^{k-1}\)
    end for
    Take two more explicit updates:
    \(\tilde{Y}_{0}=Y_{0}+\sigma \Delta \tau^{k} L_{0}\left(Y_{2}-v^{k-1}\right)\)
    \(\hat{Y}_{0}=\tilde{Y}_{0}+\mu \Delta \tau^{k} L\left(Y_{2}-v^{k-1}\right)\)
    Take \(d\) implicit corrector steps in each dimension:
    for \(j=1, \ldots, d\) do
        compute \(g=\tilde{Y}_{j-1}-\theta \Delta \tau^{k} L_{j} v^{k-1}\)
        reorder the entries in \(g\) such that \(L_{j}\) is tridiagonal
        Solve the linear system \(\left(I-\theta \Delta \tau^{k} L_{j}\right) \hat{Y}_{j}=g\)
    end for
    reorder \(\hat{Y}_{d}\) to the original ordering
    \(\operatorname{assign} v^{k}\) to \(\hat{Y}_{d}\)
```

where $D_{x}$ and $D_{y}$ are matrices resulting from the discretization of the $\sigma_{x} x v_{x}$ and $\sigma_{y} y v_{y}$ terms, $I$ is the identity matrix of size $N+1$, and $L_{x}, L_{y}$ are matrices resulting from the discretization of just the $x$ and just the $y$ terms, with the no-derivative term $(-r u)$ shared equally between $L_{x}$ and $L_{y}$.

Let $V^{k}$ be an $N+1 \times N+1$ matrix containing the computed solution at time $\tau^{k}$; so $u=\operatorname{vec}(U)$ where vec denotes the vectorization operator. Then we will use the following useful identity:

$$
\begin{equation*}
(A \otimes B) v=\operatorname{vec}\left(B V A^{T}\right) \tag{14}
\end{equation*}
$$

where $V=\operatorname{vec}^{-1}(v)$ is the inverse of the vectorization operator applied to $v$.
Then,

$$
\begin{align*}
L_{0} u & =\rho D_{y} U D_{x}^{T}  \tag{15}\\
L_{1} u & =U L_{x}^{T}  \tag{16}\\
L_{2} u & =L_{y} U \tag{17}
\end{align*}
$$

In addition (after reordering the entries in $v$ appropriately),

$$
\begin{align*}
& I-\theta \Delta \tau L_{1}=I \otimes I-I \otimes \theta \Delta \tau L_{x}=I \otimes\left(I-\theta \Delta \tau L_{x}\right)  \tag{18}\\
& I-\theta \Delta \tau L_{2}=I \otimes I-I \otimes \theta \Delta \tau L_{y}=I \otimes\left(I-\theta \Delta \tau L_{y}\right) \tag{19}
\end{align*}
$$

Therefore,

$$
\begin{equation*}
\left(I-\theta \Delta \tau L_{1}\right)^{-1}=\left(I \otimes\left(I-\theta \Delta \tau L_{x}\right)\right)^{-1}=I^{-1} \otimes\left(I-\theta \Delta \tau L_{x}\right)^{-1}=I \otimes\left(I-\theta \Delta \tau L_{x}\right)^{-1} \tag{20}
\end{equation*}
$$

and we can apply identity (11) to (17) as well. Note that we never actually compute the inverse of any matrix, but (17) is used to illustrate how we can avoid solving an $\mathcal{O}\left(N^{2}\right)$ system and instead solve an $\mathcal{O}(N)$ system. Note that in Algorithm 2, all the update equations only involve multiplying with or solving linear systems of size $N$.

```
Algorithm 2 Matrix-form of Craig-Sneyd ADI timestepping in two dimensions
    Take an explicit step:
    \(Y_{0}^{T}=U^{T}+\Delta \tau^{k}\left(\rho D_{y} U D_{x}^{T}+L_{x} U^{T}+\left(U L_{y}\right)^{T}\right)\)
    Take two implicit corrector steps, one in each dimension
    Solve the linear system \(\left(I-\theta \Delta \tau^{k} L_{x}\right) Y_{1}^{T}=Y_{0}-\theta \Delta \tau^{k} L_{x} U^{T}\)
    Solve the linear system \(\left(I-\theta \Delta \tau^{k} L_{y}\right) Y_{2}=Y_{1}-\theta \Delta \tau^{k} L_{y} U\)
    Take two more explicit updates:
    \(\tilde{Y}_{0}=Y_{0}+\sigma \rho \Delta \tau^{k} D_{y} U D_{x}^{T}\)
    \(\hat{Y}_{0}=\tilde{Y}_{0}+\mu \Delta \tau^{k}\left[\rho D_{y}\left(Y_{2}-U\right) D_{x}^{T}+L_{x}\left(Y_{2}-U\right)^{T}+\left(\left(Y_{2}-U\right) L_{y}\right)^{T}\right]\)
    Take two implicit corrector steps, one in each dimension
    Solve the linear system \(\left(I-\theta \Delta \tau^{k} L_{x}\right) \tilde{Y}_{1}^{T}=\tilde{Y}_{0}^{T}-\theta \Delta \tau^{k} L_{x} U^{T}\)
    Solve the linear system \(\left(I-\theta \Delta \tau^{k} L_{y}\right) U^{k}=\tilde{Y}_{1}-\theta \Delta \tau^{k} L_{y} U\)
```

For three and higher dimensions, we cannot use the matrix identity (14) in the same way, however, we can still take advantage of the special structure of the matrix. The tridiagonal matrix that we solve in lines 7 and 16 in Algorithm 1 can be decomposed as follows:

$$
\begin{equation*}
I-\theta \Delta \tau^{k} L_{j}=I \otimes I \otimes \cdots \otimes I \otimes\left(I-\theta \Delta \tau^{k} L_{x}\right) \tag{21}
\end{equation*}
$$

Therefore, when solving the linear system

$$
\begin{equation*}
\left(I-\theta \Delta \tau^{k} L_{j}\right) v=g \tag{22}
\end{equation*}
$$

for $u$, we can rearrange the $(N+1)^{d}$ length vector $g$ into an $N+1 \times(N+1)^{d-1}$ matrix $G$ in column major order. Then instead of solving Equation (22) we only need to solve the $N+1$ by $N+1$ linear system

$$
\begin{equation*}
\left(I-\theta \Delta \tau^{k} L_{x}\right) V=G \tag{23}
\end{equation*}
$$

After solving, $V$ can be rearranged into a vector again for explicit updates.

### 3.5 American options

In contrast to European options, American options can be exercised at any point prior to the expiry date. Therefore, the Black-Scholes model (4) becomes more complicated and can be written in the form of a linear complementarity problem (LCP)

$$
\begin{equation*}
\left(\left(v_{\tau}-L v>0\right) \wedge\left(v=V^{*}\right)\right) \vee\left(\left(v_{\tau}=L v\right) \wedge\left(v-V^{*}>0\right)\right) \tag{24}
\end{equation*}
$$

which leads to the PDE

$$
\begin{equation*}
v_{\tau}=\mathcal{L} v+p \max \left(V^{*}-v, 0\right) \tag{25}
\end{equation*}
$$

for some large value $p$, typically the reciprocal of any accuracy required.

The PDE (25) can be solved by the penalty method introduced in [6]. For ADI methods, [7] incorporate the penalty method into the Craig-Sneyd methods, leading to efficient algorithms for two-dimensional problems.

The penalty term $p \max \left(V^{*}-v, 0\right)$ is discretized as $P^{k}\left(V^{*}-v^{k}\right)$. The matrix $P^{k}$ is a diagonal matrix defined as

$$
P_{i, i}^{k}= \begin{cases}p & \text { if } v_{i}^{k}<V_{i}^{*}  \tag{26}\\ 0 & \text { otherwise }\end{cases}
$$

For discretization, the other terms in the PDE (25) are discretized in the same way as the previous section.

The penalty iteration applied to $\theta$-timestespping is given by Algorithm 3

```
Algorithm 3 Penalty Iteration for \(\theta\)-timestepping
    let \(v^{k, 0}=v^{k-1}\) and compute \(P^{k, 0}=P\left(v^{k-1}\right)\).
    for \(m=1, \ldots\) do
        solve \(\left(I-\theta \Delta \tau^{k} L^{(d)}+P^{k, m-1} v^{k, m}=\left(I+(1-\theta) \Delta \tau^{k} L^{(d)}\right) v^{k-1}+P^{k, m-1} V^{*}\right.\).
        compute \(P^{k, m}=P\left(v^{k, m}\right)\)
        if \(\max _{j}\left\{\frac{\left|v_{j}^{k, m}-v_{j}^{k, m-1}\right|}{\max \left\{1,\left|v_{j}^{k, m}\right|\right\}}\right\}\) or \(P^{k, m-1}=P^{k, m}\) then
            break
        end if
    end for
    \(v^{k},=v^{k, m}\)
```

For ADI methods, the MCS method remains largely unchanged, except Algorithm 3 is incorporated into the computation of the last term:

Unfortunately, since the penalty matrix is not the result of a kronecker (tensor) product, it is not straightforward to efficiently solve the linear system in Algorithm 2. However, we can still speed up the computation by using our more efficient algorithm for the other solves in each timestep.

Finally, we have the matrix form of MCS ADI timestepping for American options in $d$ dimensions:

## 4 Numerical results

We show numerical results for three algorithms (Crank-Nicolson, Vectorized MCS, and Matrcized MCS) applied to four problems: European and American Geometric Average Put in two and three dimensions. This section will first show the European problems in two and three dimensions in Tables 1-8, then show the computed reference values for American problems (using a one-dimensional penalty iteration) in Tables 9 and 10, and then show the American problems in Tables 11-16.
The tables for European options have 6 columns, showing the number of subintervals, the number of timesteps taken, the computed value, the signed error, the computed convergence rate, and the time recorded.

```
Algorithm 4 MCS ADI timestepping in \(d\) dimensions
    Take an explicit step:
    \(Y_{0}=v^{k-1}+\Delta \tau^{k}\left(L v^{k-1}\right)\)
    Take \(d\) implicit corrector steps, one in each dimension:
    for \(j=1, \ldots, d\) do
        compute \(g=Y_{j-1}-\theta \Delta \tau^{k} L_{j} v^{k-1}\)
        reorder the entries in \(g\) such that \(L_{j}\) is tridiagonal
        Solve the linear system \(\left(I-\theta \Delta \tau^{k} L_{j}\right) Y_{j}=Y_{j-1}-\theta \Delta \tau^{k} L_{j} v^{k-1}\)
    end for
    Take two more explicit updates:
    \(\tilde{Y}_{0}=Y_{0}+\sigma \Delta \tau^{k} L_{0}\left(Y_{2}-v^{k-1}\right)\)
    \(\hat{Y}_{0}=\tilde{Y}_{0}+\mu \Delta \tau^{k} L\left(Y_{2}-v^{k-1}\right)\)
    Take \(d\) implicit corrector steps in each dimension:
    for \(j=1, \ldots, d-1\) do
        compute \(g=\tilde{Y}_{j-1}-\theta \Delta \tau^{k} L_{j} v^{k-1}\)
        reorder the entries in \(g\) such that \(L_{j}\) is tridiagonal
        Solve the linear system \(\left(I-\theta \Delta \tau^{k} L_{j}\right) \hat{Y}_{j}=g\)
    end for
    compute \(g=\tilde{Y}_{d-1}-\theta \Delta \tau^{k} L_{d} v^{k-1}\)
    reorder the entries in \(g\) such that \(L_{d}\) is tridiagonal
    Apply Algorithm 3 to the system \(\left(I-\theta \Delta \tau^{k} L_{d}\right) P^{k} v^{k}=g+P^{k} V^{*}\)
    reorder \(v^{k}\) to the original ordering
```

```
Algorithm 5 Matrix-form of MCS ADI timestepping for American options in two dimensions
```

Algorithm 5 Matrix-form of MCS ADI timestepping for American options in two dimensions
Take an explicit step:
Take an explicit step:
$Y_{0}^{T}=U^{T}+\Delta \tau^{k}\left(\rho D_{y} U D_{x}^{T}+L_{x} U^{T}+\left(U L_{y}\right)^{T}\right)$
$Y_{0}^{T}=U^{T}+\Delta \tau^{k}\left(\rho D_{y} U D_{x}^{T}+L_{x} U^{T}+\left(U L_{y}\right)^{T}\right)$
Take two implicit corrector steps, one in each dimension
Take two implicit corrector steps, one in each dimension
Solve the linear system $\left(I-\theta \Delta \tau^{k} L_{x}\right) Y_{1}^{T}=Y_{0}-\theta \Delta \tau^{k} L_{x} U^{T}$
Solve the linear system $\left(I-\theta \Delta \tau^{k} L_{x}\right) Y_{1}^{T}=Y_{0}-\theta \Delta \tau^{k} L_{x} U^{T}$
Solve the linear system $\left(I-\theta \Delta \tau^{k} L_{y}\right) Y_{2}=Y_{1}-\theta \Delta \tau^{k} L_{y} U$
Solve the linear system $\left(I-\theta \Delta \tau^{k} L_{y}\right) Y_{2}=Y_{1}-\theta \Delta \tau^{k} L_{y} U$
Take two more explicit updates:
Take two more explicit updates:
$\tilde{Y}_{0}=Y_{0}+\sigma \rho \Delta \tau^{k} D_{y} U D_{x}^{T}$
$\tilde{Y}_{0}=Y_{0}+\sigma \rho \Delta \tau^{k} D_{y} U D_{x}^{T}$
$\hat{Y}_{0}=\tilde{Y}_{0}+\mu \Delta \tau^{k}\left[\rho D_{y}\left(Y_{2}-U\right) D_{x}^{T}+L_{x}\left(Y_{2}-U\right)^{T}+\left(\left(Y_{2}-U\right) L_{y}\right)^{T}\right]$
$\hat{Y}_{0}=\tilde{Y}_{0}+\mu \Delta \tau^{k}\left[\rho D_{y}\left(Y_{2}-U\right) D_{x}^{T}+L_{x}\left(Y_{2}-U\right)^{T}+\left(\left(Y_{2}-U\right) L_{y}\right)^{T}\right]$
Take two implicit corrector steps, one in each dimension
Take two implicit corrector steps, one in each dimension
Solve the linear system $\left(I-\theta \Delta \tau^{k} L_{x}\right) \tilde{Y}_{1}^{T}=\tilde{Y}_{0}^{T}-\theta \Delta \tau^{k} L_{x} U^{T}$
Solve the linear system $\left(I-\theta \Delta \tau^{k} L_{x}\right) \tilde{Y}_{1}^{T}=\tilde{Y}_{0}^{T}-\theta \Delta \tau^{k} L_{x} U^{T}$
compute $B=\tilde{Y}_{1}-\theta \Delta \tau^{k} L_{y} U$
compute $B=\tilde{Y}_{1}-\theta \Delta \tau^{k} L_{y} U$
Apply the penalty iteration (i.e. Algorithm 3) to the system
Apply the penalty iteration (i.e. Algorithm 3) to the system
$\left(I-\theta \Delta \tau^{k} L_{2}\right)+P^{k} v^{k}=b+P^{k} V^{*}$

```
    \(\left(I-\theta \Delta \tau^{k} L_{2}\right)+P^{k} v^{k}=b+P^{k} V^{*}\)
```

```
Algorithm 6 MCS ADI timestepping in \(d\) dimensions
    Take an explicit step:
    \(Y_{0}=v^{k-1}+\Delta \tau^{k}\left(L v^{k-1}\right)\)
    Take \(d\) implicit corrector steps, one in each dimension:
    for \(j=1, \ldots, d\) do
        compute \(g=Y_{j-1}-\theta \Delta \tau^{k} L_{j} v^{k-1}\)
        compute \(G\) according to Equation (23)
        Solve the linear system \(\left(I-\theta \Delta \tau^{k} L_{j}\right) Y_{j}=G\)
    end for
    Take two more explicit updates:
    \(\tilde{Y}_{0}=Y_{0}+\sigma \Delta \tau^{k} L_{0}\left(Y_{2}-v^{k-1}\right)\)
    \(\hat{Y}_{0}=\tilde{Y}_{0}+\mu \Delta \tau^{k} L\left(Y_{2}-v^{k-1}\right)\)
    Take \(d\) implicit corrector steps in each dimension:
    for \(j=1, \ldots, d-1\) do
        compute \(g=\tilde{Y}_{j-1}-\theta \Delta \tau^{k} L_{j} v^{k-1}\)
        compute \(G\) according to Equation (23)
        Solve the linear system \(\left(I-\theta \Delta \tau^{k} L_{j}\right) \hat{Y}_{j}=G\)
    end for
    compute \(g=\tilde{Y}_{d-1}-\theta \Delta \tau^{k} L_{d} v^{k-1}\)
    reorder the entries in \(g\) such that \(L_{d}\) is tridiagonal
    Apply Algorithm 3 to the system \(\left(I-\theta \Delta \tau^{k} L_{d}\right) P^{k} v^{k}=g+P^{k} V^{*}\)
    reorder \(v^{k}\) to the original ordering
```

The tables for American options have 7 columns, showing the number of subintervals, the number of timesteps taken, the number of penalty iterations taken, the computed value, an estimation of the error using successive differences, an estimation of the rate of convergence using the ratio of the differences, and the time recorded.

The error and convergence rate is computed point-wise at the expiry date and the strike price, which is typically near the region of interest in financial applications, since stock prices decreasing to zero or increasing many multiple times are unlikely occurances.

### 4.1 Two-Dimensional European Geometric Average Put

| $N$ | timesteps | Value | Error | Rate | time $(\mathrm{s})$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 10 | 12 | 8.490342 | $-1.32 \mathrm{e}-01$ | - | $5.56 \mathrm{e}-02$ |
| 20 | 22 | 8.590928 | $-3.17 \mathrm{e}-02$ | 2.06 | $1.86 \mathrm{e}-01$ |
| 40 | 42 | 8.614176 | $-8.49 \mathrm{e}-03$ | 1.90 | $8.70 \mathrm{e}-01$ |
| 80 | 82 | 8.620456 | $-2.21 \mathrm{e}-03$ | 1.94 | $5.39 \mathrm{e}+00$ |
| 160 | 162 | 8.622131 | $-5.35 \mathrm{e}-04$ | 2.05 | $4.55 \mathrm{e}+01$ |
| 320 | 322 | 8.622530 | $-1.36 \mathrm{e}-04$ | 1.98 | $4.02 \mathrm{e}+02$ |
| 640 | 642 | 8.622632 | $-3.37 \mathrm{e}-05$ | 2.01 | $4.17 \mathrm{e}+03$ |

Table 1: Two dimensional European Geometric Average Put problem, $\sigma_{i}=0.4, \rho=0.2, K_{i}=100$, $S_{i \max }=1000, r=0.1$. Exact solution is 8.622665388263 . Solved with Crank-Nicolson.

| $N$ | timesteps | Value | Error | Rate | time(s) |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 10 | 12 | 8.502136 | $-1.21 \mathrm{e}-01$ | - | $3.30 \mathrm{e}-02$ |
| 20 | 22 | 8.585842 | $-3.68 \mathrm{e}-02$ | 1.71 | $5.19 \mathrm{e}-02$ |
| 40 | 42 | 8.613079 | $-9.59 \mathrm{e}-03$ | 1.94 | $1.51 \mathrm{e}-01$ |
| 80 | 82 | 8.620397 | $-2.27 \mathrm{e}-03$ | 2.08 | $6.82 \mathrm{e}-01$ |
| 160 | 162 | 8.622116 | $-5.49 \mathrm{e}-04$ | 2.05 | $4.44 \mathrm{e}+00$ |
| 320 | 322 | 8.622524 | $-1.41 \mathrm{e}-04$ | 1.96 | $3.54 \mathrm{e}+01$ |
| 640 | 642 | 8.622631 | $-3.48 \mathrm{e}-05$ | 2.02 | $3.06 \mathrm{e}+02$ |

Table 2: Two dimensional European Geometric Average Put problem, $\sigma_{i}=0.4, \rho=0.2, K_{i}=100$, $S_{i \max }=1000, r=0.1$. Exact solution is 8.622665388263 . Solved with vectorized Craig-Sneyd.

| $N$ | timesteps | Value | Error | Rate | time(s) |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 10 | 12 | 8.502133 | $-1.21 \mathrm{e}-01$ | - | $9.81 \mathrm{e}-01$ |
| 20 | 22 | 8.585842 | $-3.68 \mathrm{e}-02$ | 1.71 | $8.22 \mathrm{e}-02$ |
| 40 | 42 | 8.613079 | $-9.59 \mathrm{e}-03$ | 1.94 | $1.19 \mathrm{e}-01$ |
| 80 | 82 | 8.620397 | $-2.27 \mathrm{e}-03$ | 2.08 | $2.67 \mathrm{e}-01$ |
| 160 | 162 | 8.622116 | $-5.49 \mathrm{e}-04$ | 2.05 | $1.37 \mathrm{e}+00$ |
| 320 | 322 | 8.622524 | $-1.41 \mathrm{e}-04$ | 1.96 | $5.07 \mathrm{e}+00$ |
| 640 | 642 | 8.622631 | $-3.48 \mathrm{e}-05$ | 2.02 | $2.33 \mathrm{e}+01$ |
| 1280 | 1282 | 8.622657 | $-8.50 \mathrm{e}-06$ | 2.03 | $2.05 \mathrm{e}+02$ |
| 2560 | 2562 | 8.622663 | $-2.15 \mathrm{e}-06$ | 1.99 | $2.57 \mathrm{e}+03$ |

Table 3: Two dimensional European Geometric Average Put problem, $\sigma_{i}=0.4, \rho=0.2, K_{i}=100$, $S_{i \max }=1000, r=0.1$. Exact solution is 8.622665388263. Solved with matricized Craig-Sneyd.

### 4.2 Three-Dimensional European Geometric Average Put

| $N$ | timesteps | Value | Error | Rate | time(s) |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 8 | 10 | 7.344038 | $-3.30 \mathrm{e}-01$ | - | $8.77 \mathrm{e}-02$ |
| 16 | 18 | 7.635836 | $-3.84 \mathrm{e}-02$ | 3.10 | $2.48 \mathrm{e}+00$ |
| 32 | 34 | 7.669053 | $-5.16 \mathrm{e}-03$ | 2.89 | $2.44 \mathrm{e}+02$ |
| 64 | 66 | 7.673314 | $-9.00 \mathrm{e}-04$ | 2.52 | $1.76 \mathrm{e}+04$ |

Table 4: Three dimensional European Geometric Average Put problem, $\sigma_{i}=0.4, \rho=0.2, K_{i}=100$, $S_{i \max }=1000, r=0.1$. Exact solution is 7.674214. Solved with Crank-Nicolson.

| $N$ | timesteps | Value | Error | Rate | time(s) |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 8 | 10 | 7.477913 | $-1.96 \mathrm{e}-01$ | - | $1.65 \mathrm{e}-01$ |
| 16 | 18 | 7.644642 | $-2.96 \mathrm{e}-02$ | 2.73 | $2.97 \mathrm{e}-01$ |
| 32 | 34 | 7.670165 | $-4.05 \mathrm{e}-03$ | 2.87 | $2.22 \mathrm{e}+00$ |
| 64 | 66 | 7.673582 | $-6.33 \mathrm{e}-04$ | 2.68 | $3.11 \mathrm{e}+01$ |
| 128 | 130 | -47.315418 | $-5.50 \mathrm{e}+01$ | -16.41 | $5.48 \mathrm{e}+02$ |

Table 5: Three dimensional European Geometric Average Put problem, $\sigma_{i}=0.4, \rho=0.2, K_{i}=100$, $S_{i \max }=1000, r=0.1$. Exact solution is 7.674214. Solved with vectorized Craig-Sneyd. $\theta=1 / 2$.

| $N$ | timesteps | Value | Error | Rate | time $(\mathrm{s})$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 8 | 10 | 7.477653 | $-1.97 \mathrm{e}-01$ | - | $1.46 \mathrm{e}-01$ |
| 16 | 18 | 7.644627 | $-2.96 \mathrm{e}-02$ | 2.73 | $3.16 \mathrm{e}-01$ |
| 32 | 34 | 7.670165 | $-4.05 \mathrm{e}-03$ | 2.87 | $1.27 \mathrm{e}+00$ |
| 64 | 66 | 7.673582 | $-6.33 \mathrm{e}-04$ | 2.68 | $1.13 \mathrm{e}+01$ |
| 128 | 130 | -43.979294 | $-5.17 \mathrm{e}+01$ | -16.32 | $1.90 \mathrm{e}+02$ |

Table 6: Three dimensional European Geometric Average Put problem, $\sigma_{i}=0.4, \rho=0.2, K_{i}=100$, $S_{i_{\max }}=1000, r=0.1$. Exact solution is 7.674214. Solved with matricized Craig-Sneyd. $\theta=1 / 2$.

We see that when $\theta=1 / 2$, we do not have stability in convergence since for $N=128$, the value is clearly wrong, indicating that the method is not stable. However, this contradicts the [8] which states that the MCS ADI method is stable in three dimensions if $\theta>6 / 13$.

Since larger values of $\theta$ result in more stability, we attempt to remedy this by increasing $\theta$ to $2 / 3$, which achieves stability of the Craig-Sneyd method. We will use $\theta=2 / 3$ for American options in the subsequent section.

| $N$ | timesteps | Value | Error | Rate | time $(\mathrm{s})$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 8 | 10 | 7.485199 | $-1.89 \mathrm{e}-01$ | - | $1.88 \mathrm{e}-01$ |
| 16 | 18 | 7.644765 | $-2.94 \mathrm{e}-02$ | 2.68 | $3.12 \mathrm{e}-01$ |
| 32 | 34 | 7.670106 | $-4.11 \mathrm{e}-03$ | 2.84 | $2.29 \mathrm{e}+00$ |
| 64 | 66 | 7.673565 | $-6.49 \mathrm{e}-04$ | 2.66 | $2.99 \mathrm{e}+01$ |
| 128 | 130 | 7.674071 | $-1.44 \mathrm{e}-04$ | 2.17 | $6.16 \mathrm{e}+02$ |

Table 7: Three dimensional European Geometric Average Put problem, $\sigma_{i}=0.4, \rho=0.2, K_{i}=100$, $S_{i \max }=1000, r=0.1$. Exact solution is 7.674214. Solved with vectorized Craig-Sneyd. $\theta=2 / 3$.

| $N$ | timesteps | Value | Error | Rate | time(s) |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 8 | 10 | 7.484872 | $-1.89 \mathrm{e}-01$ | - | $1.89 \mathrm{e}-01$ |
| 16 | 18 | 7.644748 | $-2.95 \mathrm{e}-02$ | 2.68 | $3.22 \mathrm{e}-01$ |
| 32 | 34 | 7.670106 | $-4.11 \mathrm{e}-03$ | 2.84 | $1.12 \mathrm{e}+00$ |
| 64 | 66 | 7.673565 | $-6.49 \mathrm{e}-04$ | 2.66 | $9.87 \mathrm{e}+00$ |
| 128 | 130 | 7.674071 | $-1.44 \mathrm{e}-04$ | 2.17 | $2.69 \mathrm{e}+02$ |
| 256 | 258 | 7.674168 | $-4.67 \mathrm{e}-05$ | 1.62 | $4.90 \mathrm{e}+03$ |

Table 8: Three dimensional European Geometric Average Put problem, $\sigma_{i}=0.4, \rho=0.2, K_{i}=100$, $S_{i \max }=1000, r=0.1$. Exact solution is 7.674214. Solved with matricized Craig-Sneyd. $\theta=2 / 3$.

### 4.3 Reference values for American option pricing problems

| $N$ | timesteps | Iterations | Value | Error | Rate | time $(\mathrm{s})$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 100 | 39 | 59 | 9.4647541515 | - | - | $9.32 \mathrm{e}-03$ |
| 200 | 73 | 114 | 9.4684135388 | $3.66 \mathrm{e}-03$ | - | $1.60 \mathrm{e}-02$ |
| 400 | 141 | 226 | 9.4692758928 | $8.62 \mathrm{e}-04$ | 2.09 | $4.46 \mathrm{e}-02$ |
| 800 | 277 | 456 | 9.4694864416 | $2.11 \mathrm{e}-04$ | 2.03 | $1.42 \mathrm{e}-01$ |
| 1600 | 547 | 912 | 9.4695394777 | $5.30 \mathrm{e}-05$ | 1.99 | $5.05 \mathrm{e}-01$ |
| 3200 | 1087 | 1778 | 9.4695525711 | $1.31 \mathrm{e}-05$ | 2.02 | $1.94 \mathrm{e}+00$ |
| 6400 | 2167 | 3465 | 9.4695558064 | $3.24 \mathrm{e}-06$ | 2.02 | $7.06 \mathrm{e}+00$ |
| 12800 | 4326 | 6943 | 9.4695566300 | $8.24 \mathrm{e}-07$ | 1.97 | $2.76 \mathrm{e}+01$ |
| 25600 | 8643 | 14199 | 9.4695568442 | $2.14 \mathrm{e}-07$ | 1.94 | $1.12 \mathrm{e}+02$ |

Table 9: One-dimensional American Put problem, with parameters chosen to match the two-dimensional problem.

| $N$ | timesteps | Iterations | Value | Error | Rate | time $(\mathrm{s})$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 100 | 37 | 56 | 8.4046980424 | - | - | $7.55 \mathrm{e}-03$ |
| 200 | 68 | 107 | 8.4076796836 | $2.98 \mathrm{e}-03$ | - | $1.51 \mathrm{e}-02$ |
| 400 | 132 | 213 | 8.4084829388 | $8.03 \mathrm{e}-04$ | 1.89 | $4.17 \mathrm{e}-02$ |
| 800 | 257 | 426 | 8.4086725800 | $1.90 \mathrm{e}-04$ | 2.08 | $1.33 \mathrm{e}-01$ |
| 1600 | 507 | 851 | 8.4087203268 | $4.77 \mathrm{e}-05$ | 1.99 | $4.71 \mathrm{e}-01$ |
| 3200 | 1006 | 1665 | 8.4087321856 | $1.19 \mathrm{e}-05$ | 2.01 | $1.83 \mathrm{e}+00$ |
| 6400 | 2004 | 3230 | 8.4087351562 | $2.97 \mathrm{e}-06$ | 2.00 | $6.58 \mathrm{e}+00$ |
| 12800 | 3999 | 6482 | 8.4087359061 | $7.50 \mathrm{e}-07$ | 1.99 | $2.60 \mathrm{e}+01$ |
| 25600 | 7989 | 13173 | 8.4087360994 | $1.93 \mathrm{e}-07$ | 1.96 | $1.04 \mathrm{e}+02$ |

Table 10: One-dimensional American Put problem, with parameters chosen to match the threedimensional problem.

### 4.4 Two-Dimensional American Geometric Average Put

| $N$ | timesteps | Iterations | Value | Difference | Rate | time(s) |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 10 | 27 | 30 | 9.255044 | - | - | $3.50 \mathrm{e}-02$ |
| 20 | 62 | 75 | 9.415944 | $-1.61 \mathrm{e}-01$ | - | $1.45 \mathrm{e}-01$ |
| 40 | 134 | 191 | 9.455322 | $-3.94 \mathrm{e}-02$ | 2.03 | $1.38 \mathrm{e}+00$ |
| 80 | 274 | 460 | 9.465989 | $-1.07 \mathrm{e}-02$ | 1.88 | $1.54 \mathrm{e}+01$ |
| 160 | 552 | 1053 | 9.468645 | $-2.66 \mathrm{e}-03$ | 2.01 | $2.84 \mathrm{e}+02$ |

Table 11: Two-Dimensional American Geometric Average Put problem, $\sigma_{i}=0.4, \rho=0.2, K_{i}=100$, $S_{i_{\max }}=1000, r=0.1$. Solved with Crank-Nicolson.

| $N$ | timesteps | Iterations | Value | Difference | Rate | time $(\mathrm{s})$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 10 | 27 | 34 | 9.229663 | - | - | $1.44 \mathrm{e}-02$ |
| 20 | 62 | 84 | 9.404770 | $-1.75 \mathrm{e}-01$ | - | $4.64 \mathrm{e}-02$ |
| 40 | 134 | 218 | 9.449318 | $-4.45 \mathrm{e}-02$ | 1.97 | $3.01 \mathrm{e}-01$ |
| 80 | 274 | 500 | 9.462960 | $-1.36 \mathrm{e}-02$ | 1.71 | $2.32 \mathrm{e}+00$ |
| 160 | 552 | 1105 | 9.467121 | $-4.16 \mathrm{e}-03$ | 1.71 | $1.90 \mathrm{e}+01$ |
| 320 | 1102 | 2221 | 9.468552 | $-1.43 \mathrm{e}-03$ | 1.54 | $1.57 \mathrm{e}+02$ |

Table 12: Two-Dimensional American Geometric Average Put problem, $\sigma_{i}=0.4, \rho=0.2, K_{i}=100$, $S_{i \text { max }}=1000, r=0.1$. Solved with Craig-Sneyd.

| $N$ | timesteps | Iterations | Value | Difference | Rate | time $(\mathrm{s})$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 10 | 27 | 34 | 9.229664 | - | - | $2.28 \mathrm{e}-02$ |
| 20 | 62 | 84 | 9.404770 | $-1.75 \mathrm{e}-01$ | - | $5.07 \mathrm{e}-02$ |
| 40 | 134 | 218 | 9.449318 | $-4.45 \mathrm{e}-02$ | 1.97 | $1.93 \mathrm{e}-01$ |
| 80 | 274 | 501 | 9.462960 | $-1.36 \mathrm{e}-02$ | 1.71 | $1.34 \mathrm{e}+00$ |
| 160 | 552 | 1107 | 9.467121 | $-4.16 \mathrm{e}-03$ | 1.71 | $1.10 \mathrm{e}+01$ |
| 320 | 1102 | 2222 | 9.468552 | $-1.43 \mathrm{e}-03$ | 1.54 | $8.48 \mathrm{e}+01$ |

Table 13: Two-Dimensional American Geometric Average Put problem, $\sigma_{i}=0.4, \rho=0.2, K_{i}=100$, $S_{i_{\max }}=1000, r=0.1$. Solved with Matricized Craig-Sneyd.

### 4.5 Three-Dimensional American Geometric Average Put

| $N_{x}$ | timesteps | Iterations | Value | Error | Rate | time(s) |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8 | 10 | 13 | 8.229706 | - | - | $8.66 \mathrm{e}-02$ |
| 16 | 18 | 30 | 8.369964 | $-1.40 \mathrm{e}-01$ | - | $2.69 \mathrm{e}+00$ |
| 32 | 34 | 66 | 8.395212 | $-2.52 \mathrm{e}-02$ | 2.47 | $3.42 \mathrm{e}+02$ |
| 64 | 66 | 131 | 8.403383 | $-8.17 \mathrm{e}-03$ | 1.63 | $3.81 \mathrm{e}+04$ |

Table 14: Three dimensional American Geometric Average Put problem, $\sigma_{i}=0.4, \rho=0.2, K_{i}=100$, $S_{i_{\max }}=1000, r=0.1$. Solved with Crank-Nicolson.

| $N$ | timesteps | Iterations | Value | Error | Rate | time(s) |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8 | 10 | 27 | 8.230826 | - | - | $5.96 \mathrm{e}-02$ |
| 16 | 18 | 57 | 8.355286 | $-1.24 \mathrm{e}-01$ | - | $1.83 \mathrm{e}-01$ |
| 32 | 34 | 122 | 8.387251 | $-3.20 \mathrm{e}-02$ | 1.96 | $2.37 \mathrm{e}+00$ |
| 64 | 66 | 226 | 8.399088 | $-1.18 \mathrm{e}-02$ | 1.43 | $3.59 \mathrm{e}+01$ |
| 128 | 130 | 448 | 8.404125 | $-5.04 \mathrm{e}-03$ | 1.23 | $7.71 \mathrm{e}+02$ |

Table 15: Three dimensional American Geometric Average Put problem, $\sigma_{i}=0.4, \rho=0.2, K_{i}=100$, $S_{i_{\max }}=1000, r=0.1$. Solved with vectorized Craig-Sneyd. $\theta=2 / 3$.

| $N$ | timesteps | Iterations | Value | Error | Rate | time(s) |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8 | 10 | 35 | 8.231054 | - | - | $7.24 \mathrm{e}-02$ |
| 16 | 18 | 70 | 8.355287 | $-1.24 \mathrm{e}-01$ | - | $1.30 \mathrm{e}-01$ |
| 32 | 34 | 143 | 8.387251 | $-3.20 \mathrm{e}-02$ | 1.96 | $1.56 \mathrm{e}+00$ |
| 64 | 66 | 274 | 8.399088 | $-1.18 \mathrm{e}-02$ | 1.43 | $2.29 \mathrm{e}+01$ |
| 128 | 130 | 539 | 8.404125 | $-5.04 \mathrm{e}-03$ | 1.23 | $5.69 \mathrm{e}+02$ |

Table 16: Three dimensional American Geometric Average Put problem, $\sigma_{i}=0.4, \rho=0.2, K_{i}=100$, $S_{i_{\max }}=1000, r=0.1$. Solved with matricized Craig-Sneyd. $\theta=2 / 3$.

## 5 Convergence and runtime analysis

We will now analyze the convergence and runtime of our algorithms. Recall that both MCS and CN are supposed to be second order accurate, and this is shown in Figure 1. For two and three dimensional European options, it is clear that the convergence is second order. For American options, it is known from [6] that the penalty method for the one-dimensional problem does not exhibit second order convergence unless variable timesteps are used. We adapted the variable timestep method to the two-dimensional problem to help fix the rate of convergence. However, more work is required for it to work with three dimensions so we left that as future work.

In Figure 2 we show the time complexity of our methods. Note that we expect that the runtime is $\mathcal{O}(N)$ for the solving of a linear system in each timestep in the MCS method for two dimensions, however, we still have to operate on $\mathcal{O}\left(N^{2}\right)$ number of inputs so the runtime is expected to be $\mathcal{O}\left(N^{3}\right)$ since we are taking $\mathcal{O}(N)$ timesteps. For the CN method, the solving of a banded linear system is $\mathcal{O}\left(N^{4}\right)$, however matlab's backslash operator uses LU with full pivoting which reduces fill-in and therefore reduces the runtime of the solving of the linear system. However, we still see that CN is significantly slower than the ADI methods for a two-dimensional European option.
For the two-dimensional American option, we see that the runtime of the CN seems to be $\mathcal{O}\left(N^{4}\right)$ and the ADI method is $\mathcal{O}\left(N^{3}\right)$ as expected.
For three-dimensional problems, we expect the runtime of CN to be $\mathcal{O}\left(N^{7}\right)$ per iteration for the banded solve, therefore the total runtime is $\mathcal{O}\left(N^{8}\right)$. We also expect the runtime of the ADI method to be $\mathcal{O}\left(N^{3}\right)$ per solve, which leads to $\mathcal{O}\left(N^{4}\right)$ total runtime for the ADI method. From the bottom half of Figure 2 it seems that both the European and the American problems have approximately the expected runtime. As mentioned before, the use of an ADI method improves the time complexity per iteration from $\mathcal{O}\left(N^{d(d-1)^{2}}\right)$


Figure 1: Convergence study for European (left) and American (right) options, in both two (top) and three (bottom) dimensions.


Figure 2: Runtime analysis for European (left) and American (right) options, in both two (top) and three (bottom) dimensions.
to $\mathcal{O}\left(N^{d}\right)$ (assuming that a banded solver is used) by reducing the bandwidth to the minimum required. Our experimental results support this theory, although with the use of a general sparse LU solver and other expensive operations such as computing some dense matrix-matrix products the experimental results don't agree entirely with the theory, as is expected.

## 6 Conclusions

We see that the ADI method is a powerful tool which allows us to minimize the complexity of computing the solution at the next timestep. The time complexity of the MCS ADI method is linear in the number of unknowns that we are required to solve, which is asymptotically optimal and cannot be reduced further. We also see that when solving the linear systems, we can exploit the structure of the matrix (it is block diagonal with identical tridiagonal blocks) to reduce the runtime significantly in the case of European options. In the case of American options, the penalty matrix destroys the block-diagonal structure with identical blocks, so we cannot avoid solving a "full size" matrix. Finally, we see that the variable timestepping method introduced in [6] doesn't extend that well to two-dimensions when using ADI methods and do not work in three-dimensions. Additionally, our understanding of the methods according to [8] is incomplete, since according to the analysis in the papers $\theta \geq 6 / 13$ is required for stability of the MCS ADI method in three dimensions, however we do not see stability with $\theta=1 / 2$ and need to increase $\theta$ to $2 / 3$ to see stability with the three-dimensional problem.

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