

Total marks 50 - Total time is 50 minutes. Answer all 4 questions.

Instructions: This midterm is open book, open notes. Non-programmable calculators allowed. No electronic communication devices allowed.

The line numbers in grammars are for reference only and are not part of the grammars. Ellipses ... indicate omitted, correct text. ϵ indicates an empty string.

If you need to make any assumptions in order to answer a question, state the assumptions clearly in your answer book.

Remember: an explanation of your reasoning is always more important than the (correct) Yes/No answer.

I. [10 marks] Regular expressions and automata.

1. [4 marks] Write a regular expression defining the language of comments consisting of a string surrounded by `/*` and `*/`, unless it is inside double quotes (`"`). For example, `/* aaa */` is a comment and so is `/* a"*/"a */`, but `/* aa"*/` is not. Assume your alphabet is Σ .
2. [4 marks] Design a finite automaton for this language.
3. [2 marks] Is your automaton deterministic or non-deterministic? Explain your answer.

II. [16 marks] LL grammars and parsing.

Consider the following context-free grammar:

$$\begin{array}{lcl} 1: & S & ::= S (S) S \\ 2: & & \mid \epsilon \end{array}$$

1. [2 marks] Give either a leftmost or a rightmost derivation for string `((()))`. Explain which one you produced.
2. [3 marks] Is this grammar ambiguous? Justify your answer.
3. [3 marks] What is the language generated by this grammar?
4. [4 marks] Is this grammar LL(1)? If not, construct an equivalent one that is.
5. [4 marks] Prove that it is LL(1).

III. [14 marks] LL and LR parsing.

1. [8 marks] Show that the following grammar is SLR(1):

```
1:  S  ::= S A
2:      |  A
3:  A  ::= a
```

2. [3 marks] Without building the table for LALR(1), explain why this grammar is LALR(1) as well. Just saying that $L(\text{SLR}(1)) \subseteq L(\text{LALR}(1))$ gets 0 marks.
3. [3 marks] Explain the intuition why LR parsing recognizes more grammars than LL.

IV. [10 marks] Run-time analysis.

Show the *run-time* checks that have to be made to make sure that the following program does not have type and range errors.

```
var j, j:  2..50
    m, n: integer
    A: array [3..75] of -10..-5

i = i-1;
m = i+j;
A[m] = n;
A[3m-7] = -2j;
```