1. [15 marks] The grammar below describes a simplified version of declarations in C.

Convert this grammar to an LL(1) grammar.
2. [20 marks] Describe the semantic analysis checks that a competent compiler would make on the following fragment of Java.

```java
public class OrderedList {
    private boolean empty;
    private OrderedList left, right;
    private Comparable val;

    public void addElement(Comparable EL) throws NullPointerException, DuplicateException, ClassCastException {
        if (val == null) throw new NullPointerException("addElement");
        if (empty) {
            left = new OrderedList();
            right = new OrderedList();
            val = EL;
            empty = false;
            return;
        }
        int n = EL.compareTo(val);
        if (n == 0) throw new DuplicateException("addElement");
        if (n < 0)
            left.addElement(val);
        else
            right.addElement(val);
    }
}
```

3. [15 marks] In the Postscript programming language, arrays may be created dynamically using the `[ ]` operators. Arrays in Postscript are heterogeneous. This means that the individual elements stored in an array can each be of a different type (e.g. an integer, a real, a boolean, a string, a name (reference to a variable), the body of a function, a dictionary (a table like structure), or even another array.

Examples:

```
[ 3, 3.14, true ]
% 3 elements: integer, real, boolean

[ /ABC, 2.78, 23, {DUP ADD POP}, [ 3, 4, 5 ] ]
% 5 elements: variable, real, integer, function, array
```

Describe a set of run time mechanisms and data structures that would be sufficient to support this form of dynamic, heterogeneous array. You should assume that you also need to support strict run time subscript range checking for indexing into arrays. Checking on the appropriate use of an array element is done by the Postscript interpreter and is not part of this question.
4. [15 marks] A proposal has been made to extend a Turing/Pascal like programming language to allow arrays and strings with run time determined sizes in record type declarations. For example:

```plaintext
1     type flexRecord
2         record
3             int  A ;
4             real B [ E1 ] ;
5             boolean C ;
6             string D ( E2 ) ;
7             int  E [ 12 ] ;
8         end ;
```

Where $E1$ and $E2$ are expressions that are evaluated at runtime. The language designers have asked for your opinion as an ace compiler writer.

a) What implementation difficulties do you see in this proposal?

b) There are two alternatives for when the expressions in a type declaration like this get evaluated:
   b.1 Once when the type declaration is encountered during execution
   b.2 Each time that the type is used to create a variable.
Which alternative would be more difficult to implement? Why?

c) What language restrictions would you suggest to make this construct easier to implement?

5. [15 marks] A lexical analyzer for Fortran has to deal with several kinds of tokens:

- **Identifiers**: a non-empty sequence of letters and digits starting with a letter
- **Integer constants**: $D$
- **Boolean constants**: .FALSE. .TRUE.
- **Boolean operators**: .AND. .OR. .NOT.
- **Comparison operators**: .LT. .LE. .EQ. .NE. .GE. .GT.

Where $D$ stands for a non-empty sequence of decimal digits and $E$ is used to mark the exponent part of real numbers (e.g. $1.0E10$). Integer and real constants may be used in the same expression with automatic conversion as required. Fortran is a *blank insensitive* language, except inside strings, the presence or absence of blanks *anywhere* doesn’t change the meaning of the program.

Discuss the problems that this *combination of lexical tokens* will cause in the design of a lexical analyzer for Fortran. You do not have to give the design of a lexical analyzer.
6. **[15 marks]** Describe how the C struct and union declarations shown below would be laid out in memory using the structure mapping algorithm described in lecture.

```c
typedef union{
    long A ;
    float B1, B2 ;
    char C[ 17 ] ;
    struct {
        short D1 ;
        double D2 ;
    } D ;
} U ;

struct Mstruct {
    short N ;
    long P ;
    /* U from typedef on left */
    U Q[ 2 ] ;
    int R[ 3 ] ;
    double S ;
    M ;
    U ;
}
```

You should assume the sizes and alignment constraints:

<table>
<thead>
<tr>
<th>Type</th>
<th>Size</th>
<th>Align</th>
</tr>
</thead>
<tbody>
<tr>
<td>char</td>
<td>8</td>
<td>8</td>
</tr>
<tr>
<td>short</td>
<td>16</td>
<td>16</td>
</tr>
<tr>
<td>int</td>
<td>32</td>
<td>32</td>
</tr>
<tr>
<td>long</td>
<td>64</td>
<td>32</td>
</tr>
<tr>
<td>float</td>
<td>32</td>
<td>32</td>
</tr>
<tr>
<td>double</td>
<td>64</td>
<td>64</td>
</tr>
</tbody>
</table>

7. **[10 marks]** Describe the symbol and type table entries that a typical compiler might make for the data structures in the previous question.

8. **[15 marks]** Describe the optimizations that a good optimizing compiler would perform on the fragment of C code shown below. You may assume that int and floats take 4 bytes of storage. You can show only the final results of the optimization if you make it clear what optimizations have been performed.

```c
float A[ 100 ] , C[ 200 ] ;
double B[ 100 ][ 200 ] ;
int I , J , K ;
/* Assume that A, B and C are given values here */

for( I = 0 ; I < 200 ; I++ )
    B[ 0 ][ I ] = C[ I ] * C[ I ] ;
B[ 0 ][ 0 ] = A[ 0 ] ;
for( I = 1 ; I < 100 ; I++ )
for( I = 1 ; I < 100 ; I++ )
    for( J = 1 ; J < 200 ; J++ )
```