## CSC2556

# Algorithms for Collective Decision Making 

Nisarg Shah

## Introduction

- People
> Instructor: Nisarg Shah (/~nisarg, nisarg@cs)
> TA: Evi Micha (emicha@cs), Soroush Ebadian (soroush@cs)
- Info
> Course Page: cs.toronto.edu/~nisarg/teaching/2556s22/
> Discussion Board: piazza.com/utoronto.ca/winter2022/csc2556
- Meeting
> Lectures: Online (Zoom) until at least Jan 31
> Questions? Schedule 1-1 meeting by emailing me


## What?

- Collective decision making by groups of agents
> Literature of computational social choice at the intersection of computer science and economics
- Single-agent problems
> E.g., the traveling salesman problem
> A single agent wants to find the optimal route
- Multi-agent problems
> What if multiple traveling salesmen want to share a bus?
> Each agent has a different optimal route
> Tradeoff $\rightarrow$ fairness, efficiency, strategic manipulations, ...


## What?

- Models will differ in various considerations, e.g.:
> Can agents form binding contracts?
o Entering in contracts allows agents to hedge uncertainties
> Can agents exchange/pay/receive money?
- Maybe we make a decision that is less preferable to an agent but pay the agent to compensate
> What is the structure of the outcome space?
- Is there a common decision that affects everyone (e.g., voting) or does each agent receive something (e.g., resource allocation)?


## Logistics

## Optional Reference Textbooks

- Handbook of Computational Social Choice
> By Felix Brandt, Vincent Conitzer, Ulle Endriss, Jérôme Lang, and Ariel D. Procaccia
- Algorithmic Game Theory
> By Noam Nisan, Tom Roughgarden, Eva Tardos and Vijay Vazirani
- Networks, Crowds and Markets
> By David Easley and Jon Kleinberg
- Online versions available on the course web page


## Grading Policy

- 2 assignments: 40\%
- Final project: 50\%
- Embedded Ethics Module: 5\%
> Pre-module survey: $1 \%$
> Post-module survey: $1 \%$
> Post-module assignment: 3\%
- Class participation: 5\%


## Assignments

- Theoretical
> They will require deriving intricate proofs
- We will assume...
> Strong familiarity with abstract reasoning and proof techniques
> Adequate familiarity of CS concepts (e.g., algorithm design, worstcase approximation, NP-hardness)
> Adequate familiarity of math concepts (e.g., probability, statistics, linear algebra, calculus)
> No prior background in economics


## Assignments

- Individual assignments
> Free to discuss with classmates or read online material
> Must write solutions in your own words
o Easier if you do not take any pictures/notes from the discussions
> Plagiarism will be dealt with strictly!
- Citation
> For each question, you must cite the peer (write the name) or the online sources (provide links) referred, if any
> Failing to do this is also plagiarism!


## Other Policies

- "No Garbage" Policy
> Borrowed from: Prof. Allan Borodin (citation!)

1. Partial marks for viable approaches
2. Zero marks if the answer makes no sense
3. $20 \%$ marks if you admit to not knowing how to solve

- $20 \%>0 \%$ !


## Course Timeline

- (Approximate dates)
- $\approx$ Feb 1: HW1 posted
- $\approx$ Feb 15: HW1 due
- $\approx$ Mar 1: HW2 posted, project proposal due
- $\approx$ Mar 15: HW2 due
- $\approx$ Week of Mar 15: Mid-project check-in
- Last 1.5-2 lectures: Project presentations
- $\approx$ April 10: Project reports due


## Course Project

- How?
> Groups of 1-3
- Larger groups are better
- Find partners early, but maybe after the enrollment stabilizes
- What?
> Empirical: Quantitative analysis of algorithms presented in class (or your own) using simulations or real data
> Theoretical: Prove new observations about the algorithms or design new algorithms for a problem
> Ideal: A bit of both


## Project Topic

- From your own research area of interest
> We'll introduce broad concepts that you may be able to apply to your own research area in order to find a project topic
> E.g., fairness, allocation efficiency, preference elicitation, ...
- From the course
> I'll mention some open problems as we go along
> Later, I'll also post sample projects from previous years as well as sample project ideas for this year
> You can also study realistic variants of problems that we see in class


## Course Project: Timeline

- Find partners and think about a project idea
- Submission 1: Project proposal
> Ideally 1 page but up to 2 pages excluding references
> Outline of the idea, prior work, reasonable goals
- Mid-project meetings
> Optional, 1-1 with me, 30-minute
- Class presentations
- Submission 2: Final project report
> Up to 5 pages excluding references and appendix
> Focus on quality academic writing


## Introductions

## Brief Introductions

- What to say?
> Which program?
> Which year?
> Who are you working with (if any)?
> What is your area of interest (if any)?
> Anything else you'd like to share


## Overview of the Course

## Social Choice, Mechanism Design

- Social choice
> Given the preferences of the agents, which collective decision is the most desirable?
> Fairness, welfare, ethics, resource utilization, ...
- Mechanism design
> Agents have private information, which they may lie about
> How to design the "rules of the game" such that selfish agent behavior results in desirable outcomes
> We call this "implementing" the social choice rule


## Mechanism Design

- With money
> Principal can "charge" the agents (require payments)
> Helps significantly
> Example: auctions
- Without money
> Monetary transfers are not allowed
> Incentives must be balanced otherwise
> Often impossible without sacrificing the objective a little
> Example: elections, kidney exchange


## Example: Auction

Objective: The one who really needs it more should have it.
Rule 1: Each would tell me his/her value. I'll give it to the one with the higher value.


## Example: Auction

Objective: The one who really needs it more should have it.
Rule 2: Each would tell me his/her value. l'll give it to the one with the higher value, but they must pay me that value.


## Example: Auction

Objective: The one who really needs it more should have it.
Question: Can I make it easier so that each can just truthfully tell me how much they value it?


## Real-World Applications

- Auctions form a significant part of mechanism design with money
- Auctions are ubiquitous in the real world!
> A significant source of revenue for many large organizations (including Facebook and Google)
> Often run billions of tiny auctions everyday
> Need the algorithms to be fast


## Example: Facility Location



Cost to each agent: Distance from the hospital
Objective: Minimize the sum of costs
Constraint: No money

## Example: Facility Location



Q: What is the optimal hospital location?
Q: If we decide to choose the optimal location, will the agents really tell us where they live?

## Example: Facility Location



Cost to each agent: Distance from the hospital
Objective: Minimize the maximum cost
Constraint: No money

## Example: Facility Location



Q: What is the optimal hospital location?
Q: If we decide to choose the optimal location, will the agents really tell us where they live?

## Real-World Applications



Roth


Gale


Shapley

## Matching

- National Resident Matching Program (NRMP)
- School Choice (New York, Boston)

Fair Division splíddít

Voting
(2) ROBOVOTE

## Voting Theory

## Social Choice Theory



- Mathematical theory for aggregating individual preferences into collective decisions


## Voting Theory

- Originated in ancient Greece
- Formal foundations
> $13^{\text {th }}$ Century (Ramon Llull)
$>18^{\text {th }}$ Century (Marquis de Condorcet and Jean-Charles de Borda)
> 19th Century: Charles Dodgson (a.k.a. Lewis Carroll)
> $20^{\text {th }}$ Century: Nobel prizes to Kenneth Arrow and Amartya Sen



## Voting Theory

- We want to select a collective decision based on (possibly different) individual preferences
> Presidential election, restaurant/movie selection for group activity, committee selection, facility location, ...
- Resource allocation is a special case
> You can think of all possible allocations as the different "outcomes"
- A very restricted case due to lots of ties
- An agent is indifferent among all allocations in which the resources she gets are the same
> We want to study the general case


## Voting Framework

- Set of voters $N=\{1, \ldots, n\}$
- Set of alternatives $A,|A|=m$
- Voter $i$ has a preference ranking $>_{i}$ over the alternatives
- Preference profile $\gg$ is the collection of all voters' rankings


## Voting Framework

- Social choice function $f$
> Takes as input a preference profile $\overrightarrow{>}$
> Returns an alternative $a \in A$
- Social welfare function $f$
> Takes as input a preference profile $\overrightarrow{>}$
> Returns a societal preference $>^{*}$
- For now, voting rule = social choice function

| $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ |
| :---: | :---: | :---: |
| a | c | b |
| b | a | a |
| c | b | c |

## Voting Rules

- Plurality
> Each voter awards one point to her top alternative
> Alternative with the most point wins
> Most frequently used voting rule
> Almost all political elections use plurality
- Problem?

| $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ |
| :---: | :---: | :---: | :---: | :---: |
| a | a | a | b | b |
| b | b | b | c | c |
| c | c | c | d | d |
| d | d | d | e | e |
| e | e | e | a | a |

Winner
a

## Voting Rules

- Borda Count
> Each voter awards $m-k$ points to alternative at rank $k$
> Alternative with the most points wins
> Proposed by Ramon Llull in the $13^{\text {th }}$ Century but named after $18^{\text {th }}$ Century work by Jean-Charles de Borda
> Used for elections to the national assembly of Slovenia

| $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ |
| :---: | :---: | :---: |
| $\mathrm{a}(2)$ | $\mathrm{c}(2)$ | $\mathrm{b}(2)$ |
| $\mathrm{b}(1)$ | $\mathrm{a}(1)$ | $\mathrm{a}(1)$ |
| $\mathrm{c}(0)$ | $\mathrm{b}(0)$ | $\mathrm{c}(0)$ |


| Total |
| :---: |
| a: $2+1+1=4$ |
| b: $1+0+2=3$ |
| c: $0+2+0=2$ |

Winner
a

## Current uses [edit

Political uses [edit]



 1991, tactical voting has been an important feature of the nominating process.

The modified Borda count has been used by the Green Party of Ireland to elect its chairperson. ${ }^{\text {(5IIT] }}$
 UDA.

## Other uses [edit]

The Borda count is used in elections by some educational institutions in the United States.

- University of Michigan
- Central Student Government
- Student Government of the College of Literature, Science and the Arts (LSASG)
- University of Missouri: officers of the Graduate-Professional Council
- University of California Los Angeles: officers of the Graduate Student Association
- Harvard University: officers of the Civil Liberties Union
- Southern Illinois University at Carbondale: officers of the Faculty Senate,
- Arizona State University: officers of the Department of Mathematics and Statistics assembly.
- Wheaton College, Massachusetts: faculty members of committees.
- College of William and Mary: members of the faculty personnel committee of the School of Business Administration (tie-breaker)


## Borda count in real life

The Borda count is used in elections by some professional and technical societies.

- International Society for Cryobiology: Board of Governors.
- Tempo sustainable design network: management committee.
- U.S. Wheat and Barley Scab Initiative: members of Research Area Committees.
- X.Org Foundation: Board of Directors

The OpenGL Architecture Review Board uses the Borda count as one of the feature-selection methods

The modified Borda count is used to elect the President for the United States member committee of AIESEC.
The Borda count, and points-based systems similar to it, are often used to determine awards in competitions.
The Borda count is a popular method for granting sports awards in the United States. Uses include:

- MLB Most Valuable Player Award (baseball)
- Heisman Trophy (college football) ${ }^{[3]}$
- Ranking of NCAA college teams
 the other eight entries getting points from 8 to 1 . Although designed to favor a clear winner, it has produced very close races and even a tie.
The People's Remix Competition uses a Borda variant where each voter ranks only the top three contestants.




## Voting Rules

- Positional Scoring Rules
> Defined by a score vector $\vec{s}=\left(s_{1}, \ldots, s_{m}\right)$
> Each voter gives $s_{k}$ points to alternative at rank $k$
- A family containing many important rules
> Plurality $=(1,0, \ldots, 0)$
$>$ Borda $=(m-1, m-2, \ldots, 0)$
$>k$-approval $=(1, \ldots, 1,0, \ldots, 0) \leftarrow$ top $k$ get 1 point each
$>$ Veto $=(0, \ldots, 0,-1)$
> ...


## Voting Rules

- Plurality with runoff
> First round: two alternatives with the highest plurality scores survive
> Second round: between these two alternatives, select the one that majority of voters prefer
- Similar to the French presidential election system
> Problem: vote division
> Happened in the 2002 French presidential election


## Voting Rules

- Single Transferable Vote (STV)
> $m-1$ rounds
$>$ In each round, the alternative with the least plurality votes is eliminated
> Alternative left standing is the winner
> Used in Ireland, Malta, Australia, New Zealand, ...
- STV has been strongly advocated for due to various reasons


## STV Example

| 2 voters | $\mathbf{2}$ voters | $\mathbf{1}$ voter |
| :---: | :---: | :---: |
| a | b | c |
| b | a | d |
| c | d | b |
| d | c | a |



| 2 voters | $\mathbf{2}$ voters | $\mathbf{1}$ voter |
| :---: | :---: | :---: |
| a | b | c |
| b | a | b |
| c | c | a |


| 2 voters | $\mathbf{2}$ voters | $\mathbf{1}$ voter |
| :---: | :---: | :---: |
| b | b | b |


| 2 voters | $\mathbf{2}$ voters | $\mathbf{1}$ voter |
| :---: | :---: | :---: |
| a | b | b |
| b | a | a |

## Voting Rules

- Kemeny’s Rule
> Social welfare function (selects a ranking)
> Let $n_{a>b}$ be the number of voters who prefer $a$ to $b$
> Select a ranking $\sigma$ of alternatives $=$ for every pair $(a, b)$ where $a>_{\sigma} b$, we make $n_{b>a}$ voters unhappy
> Total unhappiness $K(\sigma)=\sum_{(a, b): a>_{\sigma} b} n_{b>a}$
> Select the ranking $\sigma^{*}$ with minimum total unhappiness
- Social choice function
> Choose the top alternative in the Kemeny ranking


## Kemeny Example

| 2 voters | $\mathbf{2}$ voters | $\mathbf{1}$ voter |
| :---: | :---: | :---: |
| a | b | c |
| b | a | d |
| c | d | b |
| d | c | a |

- $K(a>b>c>d)$
$>0 \times$ first 2 voters
$>2 x$ next 2 voters
$>5 \times$ last voter
$>9$ in total
- $K(b>a>c>d)$
$>1 \times$ first 2 voters
$>1 \times$ next 2 voters
$>4 x$ last voter
$>8$ in total


## Condorcet Winner

- Definition
- Alternative $x$ defeats $y$ in a pairwise election if a strict majority of voters prefer $x$ to $y$
> Alternative $x$ is a Condorcet winner if it defeats every other alternative in a pairwise election
- Question
> Can there be two Condorcet winners?
- Condorcet paradox
> No Condorcet winner when the majority

| $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ |
| :---: | :---: | :---: |
| a | b | c |
| b | c | a |
| c | a | b |

Majority Preference

$$
a \succ b
$$

$b>c$ preference is cyclic

## Condorcet Consistency

- Condorcet consistency
> A voting rule is Condorcet consistent if it selects the Condorcet winner whenever one exists
> On preference profiles where there is no Condorcet winner, it is free to output any winner
- Among the rules we saw so far...
> NOT Condorcet consistent: all positional scoring rules (plurality, Borda, ...), plurality with runoff, STV
> Condorcet consistent: Kemeny (Why?)


## Majority Consistency

- Majority consistency
> If a strict majority of voters rank alternative $x$ first, then $x$ must be the winner.
- Question: What is the relation between majority consistency and Condorcet consistency?

1. Majority consistency $\Rightarrow$ Condorcet consistency
2. Condorcet consistency $\Rightarrow$ Majority consistency
3. Equivalent
4. Incomparable

## Condorcet Consistency

- Copeland
> Score $(x)=$ \# alternatives $x$ beats in pairwise elections
> Select $x^{*}$ with the maximum score
> Condorcet consistent (Why?)
- Maximin
$>\operatorname{Score}(x)=\min _{y} n_{x>y}$
> Select $x^{*}$ with the maximum score
> Also Condorcet consistent (Why?)

