# Record Linkage: Similarity Measures and Algorithms

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#### Presenters



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#### Outline

- Part I: Motivation, similarity measures (90 min)
  - Data quality, applications
  - Linkage methodology, core measures
  - Learning core measures
  - Linkage based measures
- Part II: Efficient algorithms for approximate join (60 min)
- Part III: Clustering/partitioning algorithms (30 min)

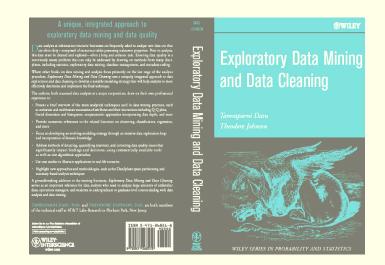
#### Data Quality: Status

- Pervasive problem in large databases
  - Inconsistency with reality: 2% of records obsolete in customer files in 1 month (deaths, name changes, etc) [DWI02]
  - Pricing anomalies : UA tickets selling for \$5, 1GB of memory selling for \$19.99 at amazon.com
- Massive financial impact
  - \$611B/year loss in US due to poor customer data [DWI02]
  - \$2.5B/year loss due to incorrect prices in retail DBs [E00]
- Commercial tools: specialized, rule-based, programmatic

#### How are Such Problems Created?

#### Human factors

- Incorrect data entry
- Ambiguity during data transformations
- Application factors
  - Erroneous applications populating databases
  - Faulty database design (constraints not enforced)
- Obsolence
  - Real-world is dynamic

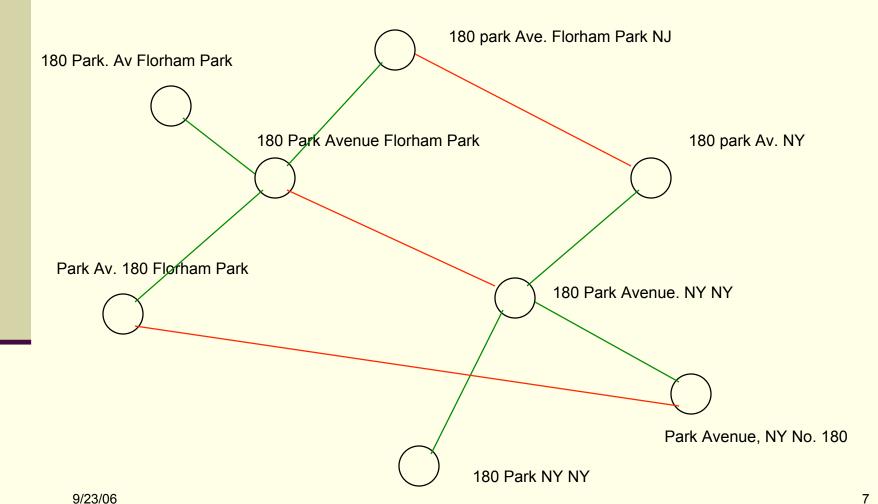


### **Application: Merging Lists**

- Application: merge address lists (customer lists, company lists) to avoid redundancy
- Current status: "standardize", different values treated as distinct for analysis
  - Lot of heterogeneity
  - Need approximate joins
- Relevant technologies
  - Approximate joins
  - Clustering/partitioning

ADDR
11810 WILLS RD ALPHARETTA GA 30076
11810 WILLS ROAD ALPHARETTA GA30004
FLR 110 RM 155 11810 WILLS RD ALPHARETTA GA 300042055
FLR 110 RM MAIN 11810 WILLIS RD ALPHARETTA GA 300042055
FLR 110 RM MAIN 11810 WILLS RD ALPHARETTA GA 30076
FLR BLDG 110 RM 155 11810 WILLS RD ALPHARETTA GA 300042055
FLR BLDG 110 RM MAIN 111810 WILLS RD ALPHARETTA GA 30076
FLR BLDG 110 RM MAIN 11810 WILLS RD ALPHARETTA GA 300042055
FLR MAIN RM 110 11810 WILLS RD ALPHARETTA GA 300042055
FLR MAIN RM BLGD 110 11810 WILLS RD ALPHARETTA GA 300042055
FLR NA RM NA 11810 WILLS RD ALPHARETTA GA 300042055
BLDG 110 FLR 1 RM 110 11810 WILLIS RD ALPHARETTA GA 30004205
BLDG 110 FLR 1 RM 110 11810 WILLS RD ALPHARETTA GA 300042055
BLDG 110 FLR 1 RM 110 11810 WILLS RD ALPHARETTA GA 300042081
BLDG 110 FLR 1 RM 110 11810 WILLS RD ALPHARETTA GA 30076
BLDG 110 FLR 1 RM RING 11810 WILLS RD ALPHARETTA GA 30004208
FLR 1 RM 1 11810 WILLS RD ALPHARETTA GA 300042055
FLR 1 RM 110 11801 WILLIS RD ALPHARETTA GA 30076
FLR 1 RM 110 11810 WILLIS RD ALPHARETTA GA 300042055
FLR 1 RM 110 11810 WILLIS RD ALPHARETTA GA 30076
FLR 1 RM 110 11810 WILLS RD ALPHARETTA GA 300042055
FLR 1 RM 110 11810 WILLS RD ALPHARETTA GA 300042081
FLR 1 RM 110 11810 WILLS RD ALPHARETTA GA 30076
FLR 1 RM 110 11810 WILLS RD ATLANTA GA 30076
FLR 1 RM 110 11810 WILLS ROAD ALPHARETTA GA 30076
FLR 1 RM BLDG 110 11810 WILLS RD ALPHARETTA GA 300042055
FLR 1 RM BLDG 11810 WILLS RD ALPHARETTA GA 300042055
FLR 1 RM COMPUTER 11810 WILLS RD ALPHARETTA GA 300042055

# Application: Merging Lists



#### **Application: Homeland Security**

- Application: correlate airline passenger data with homeland security data for no-fly lists
- Current status: "match" on name, deny boarding
  - Use more match attributes
  - Obtain more information
  - Relevant technologies
    - Schema mappings
    - Approximate joins

#### **TRAVEL**

#### 'No-fly list' keeps infants off planes

Tuesday, August 16, 2005; Posted: 4:49 a.m. EDT (08:49 GMT)

WASHINGTON (AP) -- Infants have been stopped from boarding planes at airports throughout the United States because their names are the same as or similar to those of possible terrorists on the government's "no-fly list."

It sounds like a joke, but it's not funny to parents who miss flights while scrambling to have babies' passports and other documents faxed.

Ingrid Sanden's 1-year-old daughter was stopped in Phoenix, Arizona, before boarding a flight home to Washington at Thanksgiving.

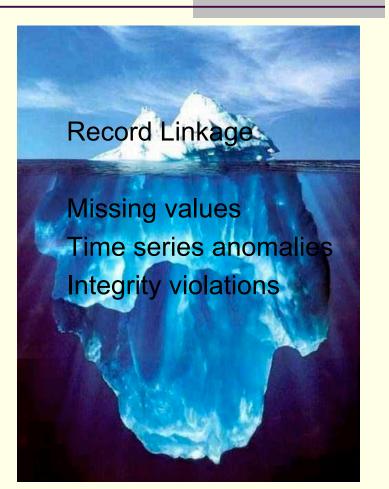
"I completely understand the war on



Ingrid Sanden holds her 1-year-old daughter, who was stopped before boarding a flight last Thanksgiving.

#### Record Linkage: Tip of the Iceberg

- An approximate join of R<sub>1</sub> and R<sub>2</sub> is
  - A subset of the cartesian product of R<sub>1</sub> and R<sub>2</sub>
  - "Matching" specified attributes of R<sub>1</sub> and R<sub>2</sub>
  - Labeled with a similarity score > t > 0
  - Clustering/partitioning of R: operates on the approximate join of R with itself.



#### The Fellegi-Sunter Model [FS69]

- Formalized the approach of Newcombe et al. [NKAJ59]
- Given two sets of records (relations) A and B perform an approximate join
  - A x B = {(a,b) |  $a \in A, b \in B$ } = M ∪ U
  - M = {(a,b) | a=b, a ∈ A, b ∈ B}; matched
  - U = {(a,b) | a <> b, a ∈ A, b ∈ B}; unmatched
- $\gamma(a,b) = (\gamma^i(a,b)) = 1..K$  comparison vector
  - Contains comparison features e.g., same last names, same SSN, etc.
  - $\Gamma$ : range of  $\gamma$ (a,b) the comparison space.

### The Fellegi-Sunter Model

- Seeking to characterize (a,b) as
  - $A_1$ : match;  $A_2$ : uncertain;  $A_3$ : non-match
- Function (linkage rule) from  $\Gamma$  to  $\{A_1 A_2 A_3\}$
- Distribution D over A x B

• m (
$$\gamma$$
) = P( $\gamma$ (a,b) | (a,b)  $\in$  M}

• 
$$u(\gamma) = P(\gamma(a,b) | (a,b) \in U)$$

#### Fellegi-Sunter Result

Sort vectors γ by m (γ)/u (γ) non increasing order; choose n < n'</p>

$$\mu = \sum_{i=1}^{n} u(\gamma_{i}) \qquad \lambda = \sum_{i=n'}^{N} m(\gamma_{i})$$

- Linkage rule with respect to minimizing  $P(A_2)$ , with  $P(A_1|U) = \mu$  and  $P(A_3|M) = \lambda$  is
  - $\gamma_1, \ldots, \gamma_n, \gamma_{n+1}, \ldots, \gamma_{n'-1}, \gamma_{n'}, \ldots, \gamma_N$

$$A_1 \qquad A_2 \qquad A_3$$

- Intuition
  - Swap i-th vector declared as A<sub>1</sub> with j-th vector in A<sub>2</sub>
  - If  $u(\gamma_i) = u(\gamma_j)$  then  $m(\gamma_j) < m(\gamma_l)$
  - After the swap, P(A<sub>2</sub>) is increased

#### Fellegi-Sunter Issues:

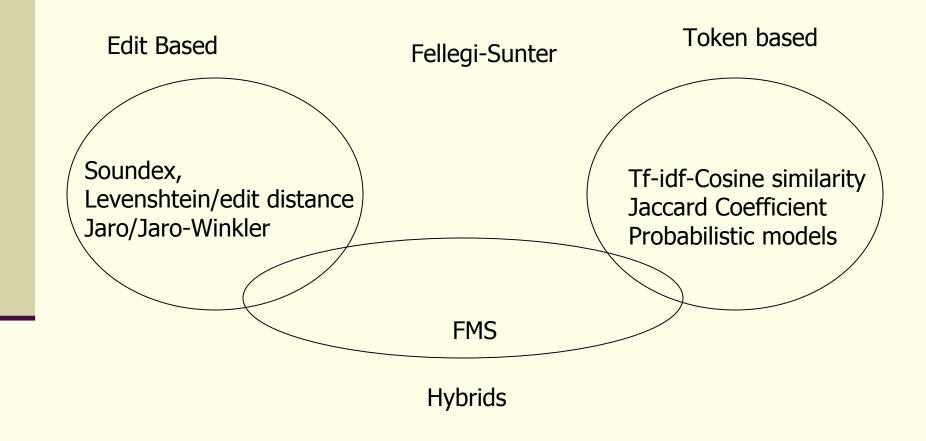
#### Tuning:

- Estimates for m (γ), u (γ) ?
  - Training data: active learning for M, U labels
  - Semi or un-supervised clustering: identify M U clusters
- Setting  $\mu$  ,  $\lambda$ ?
- Defining the comparison space  $\Gamma$ ?
  - Distance metrics between records/fields
- Efficiency/Scalability
  - Is there a way to avoid quadratic behavior (computing all |A|x|B| pairs)?

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#### Classification of the measures



#### Attribute Standardization

Several attribute fields in relations have loose or anticipated structure:

- Addresses, names
- Bibliographic entries (mainly for web data)
- Preprocessing to standardize such fields
  - Enforce common abbreviations, titles
  - Extract structure from addresses
- Part of ETL tools, commonly using field segmentation and dictionaries
- Recently machine learning approaches
  - HMM encode universe of states [CCZ02]

## Field Similarity

- Application notion of 'field'
  - Relational attribute, set of attributes, entire tuples.
- Basic problem: given two field values quantify their 'similarity' (wlog) in [0..1].
- If numeric fields, use numeric methods.
- Problem challenging for strings.

#### Soundex Encoding

- A phonetic algorithm that indexes names by their sounds when pronounced in english.
- Consists of the first letter of the name followed by three numbers. Numbers encode similar sounding consonants.
  - Remove all W, H
  - B, F, P, V encoded as 1, C,G,J,K,Q,S,X,Z as 2
  - D,T as 3, L as 4, M,N as 5, R as 6, Remove vowels
  - Concatenate first letter of string with first 3 numerals
  - Ex: great and grate become 6EA3 and 6A3E and then G63 More recent, metaphone, double metaphone etc.

### Edit Distance [G98]

- Character Operations: I (insert), D (delete), R (Replace).
- Unit costs.
- Given two strings, s,t, edit(s,t):
  - Minimum cost sequence of operations to transform s to t.
  - Example: edit(Error,Eror) = 1, edit(great,grate) = 2
- Folklore dynamic programming algorithm to compute edit();
  - Computation and decision problem: quadratic (on string length) in the worst case.

#### Edit Distance

- Several variants (weighted, block etc) -- problem can become NPcomplete easily.
- Operation costs can be learned from the source (more later)
  - String alignment = sequence of edit operations emitted by a memory-less process [RY97].
- Observations
  - May be costly operation for large strings
  - Suitable for common typing mistakes
    - Comprehensive vs Comprenhensive
  - Problematic for specific domains
    - AT&T Corporation vs AT&T Corp
    - IBM Corporation vs AT&T Corporation

## Edit Distance with affine gaps

- Differences between 'duplicates' often due to abbreviations or whole word insertions.
  - John Smith vs John Edward Smith vs John E. Smith
  - IBM Corp. vs IBM Corporation
- Allow sequences of mis-matched characters (gaps) in the alignment of two strings.
- Penalty: using the affine cost model
  - Cost(g) = s+e · I
  - s: cost of opening a gap
  - e: cost of extending the gap
  - I: length of a gap
  - Commonly e lower than s
- Similar dynamic programming algorithm

#### Jaro Rule [J89]

- Given strings  $s = a_1, ..., a_k$  and  $t = b_1, ..., b_L a_i$  in s is common to a character in t if there is a  $b_j$  in t such that  $a_i = b_j$  i-H  $\leq j \leq i$ +H where
  - H = min(|s|,|t|)/2
- Let s' = a<sub>1</sub>',...,a<sub>k</sub>' and t' = b<sub>1</sub>',...,b<sub>L</sub>' characters in s (t) common with t (s)
- A transposition for s',t' is a position i such that  $a_i' <> b_i'$ .
- Let T<sub>s',t</sub> be half the number of transpositions in s' and t'.

#### Jaro Rule

Jaro(s,t) = 
$$\frac{1}{3}(\frac{|s'|}{|s|} + \frac{|t'|}{|t|} + \frac{|s'| - T_{s',t'}}{|s'|})$$

Example:

Martha vs Marhta

- Jaro(Martha,Marhta) = 0.9722
- Jonathan vs Janathon
  - H = 4, s' = jnathn, t' = jnathn,  $T_{s',t'} = 0$
  - Jaro(Jonathan, Janathon) = 0.5

#### Jaro-Winkler Rule [W99]

- Uses the length P of the longest common prefix of s and t; P' = max(P,4)
- Jaro-Winkler(s,t) =  $Jaro(s,t) + \frac{P'}{10}(1 Jaro(s,t))$
- Example:
  - $\blacksquare$  JW(Martha,Marhta) = 0.9833
  - JW(Jonathan, Janathon) = 0.7
- Observations:
  - Both intended for small length strings (first, last names)

#### Term (token) based

- Varying semantics of 'term'
  - Words in a field
    - AT&T Corporation' -> 'AT&T', 'Corporation'
  - Q-grams (sequence of q-characters in a field)
    - {'AT&', 'T&T', '&T ', 'T C','
      - Co','orp','rpo','por','ora','rat','ati','tio','ion'} 3-grams
- Assess similarity by manipulating sets of terms.

#### Overlap metrics

- Given two sets of terms S, T
  - Jaccard coef.: Jaccard(S,T) = |S∩T|/|S∪T|
  - Variants
    - If scores (weights) available for each term (element in the set) compute Jaccard() only for terms with weight above a specific threshold.
- What constitutes a good choice of a term score?

# TF/IDF [S83]

- Term frequency (tf) inverse document frequency (idf).
- Widely used in traditional IR approaches.
- The tf/idf value of a 'term' in a document:
  - log (tf+1) \* log idf where
    - If : # of times 'term' appears in a document d
    - idf : number of documents / number of documents containing 'term'
  - Intuitively: rare 'terms' are more important

#### TF/IDF

- Varying semantics of 'term'
  - Words in a field
    - AT&T Corporation' -> 'AT&T', 'Corporation'
  - Qgrams (sequence of q-characters in a field)
    - { 'AT&','T&T','&T ', 'T C','
      - Co','orp','rpo','por','ora','rat','ati','tio','ion'} 3-grams
- For each 'term' in a field compute its corresponding tfidf score using the field as a document and the set of field values as the document collection.

#### Probabilistic analog (from FS model)

- Ps(j) : probability for j in set S
- γ<sup>j</sup>: event that values of corresponding fields are j in a random draw from sets A and B
- m  $(\gamma^j) = P(\gamma^j | M) = P_{A \cap B}(j)$
- $u(\gamma^j) = P(\gamma^j|U) = P_A(j)P_B(j)$
- Assume  $P_A(j) = P_B(j) = P_{A \cap B}(j)$ 
  - Provide more weight to agreement on rare terms and less weight to common terms
- IDF measure related to Fellegi-Sunter probabilistic notion:
  - Log(m( $\gamma^{str}$ )/u( $\gamma^{str}$ )) = log(P<sub>A∩B</sub>(str)/P<sub>A</sub> (str)P<sub>B</sub> (str)) = log(1/P<sub>A</sub>(str)) = IDF(str)

### Cosine similarity

- Each field value transformed via tfidf weighting to a (sparse) vector of high dimensionality d.
- Let a,b two field values and  $S_a$ ,  $S_b$  the set of terms for each. For w in  $S_a(S_b)$ , denote  $W(w,S_a)(W(w,S_b))$  its tfidf score.
  - For two such values:
    - Cosine(a,b) =  $\sum_{z \in Sa \cap Sb} W(z,Sa)W(z,Sb)$

### Cosine similarity

Suitable to assess closeness of

- AT&T Corporation', 'AT&T Corp' or 'AT&T Inc'
  - Low weights for 'Corporation','Corp','Inc'
  - Higher weight for 'AT&T'
  - Overall Cosine('AT&T Corp','AT&T Inc') should be high
- Via q-grams may capture small typing mistakes
  - 'Jaccard' vs 'Jacard' -> {'Jac', 'acc', 'cca', 'car', 'ard'} vs {'Jac', 'aca', 'car', 'ard'}
  - Common terms 'Jac', 'car', 'ard' would be enough to result in high value of Cosine('Jaccard','Jacard').

# Hybrids [CRF03]

Let S = {a<sub>1</sub>,...,a<sub>K</sub>}, T = {b<sub>1</sub>,...,b<sub>L</sub>} sets of terms:  
Sim(S,T) = 
$$\frac{1}{K} \sum_{i=1}^{K} \max_{j=1}^{L} sim'(a_i, b_j)$$
  
Sim'() some other similarity function  
C(t,S,T) = {w \in S s.t \exists v \in T, sim'(w,v) > t}  
D(w,T) = max<sub>v \in T</sub>sim'(w,v), w \in C(t,S,T)  
sTFIDF =  $\sum_{w \in C(t,S,T)} W(w,S) * W(w,T) * D(w,T)$ 

#### Other choices for term score?

- Several schemes proposed in IR
  - Okapi weighting
    - Model within document term frequencies as a mixture of two poisson distributions: one for relevant and one for irrelevant documents
  - Language models
    - Given Q=t<sub>1</sub>,...t<sub>n</sub> estimate p(Q|M<sub>d</sub>)
    - MLE estimate for term t : p(t|M<sub>d</sub>) = tf(t,d)/dl<sub>d</sub>
      - dl<sub>d</sub>:total number of tokens in d
    - Estimate p<sub>avg</sub>(t)
    - Weight it by a risk factor (modeled by a geometric distribution)
  - HMM

#### Fuzzy Match Similarity [CGGM03]

- Sets of terms S, T
- Main idea: cost of transforming S to T, tc(S,T).
- Transformation operations like edit distance.
  - Replacement cost: edit(s,t)\*W(s,S)
  - Insertion cost: c<sub>ins</sub> W(s,S) (c<sub>ins</sub> between 0,1)
  - Deletion cost: W(s,S)
- Computed by DP like edit()
- Generalized for multiple sets of terms

#### Fuzzy Match Similarity

#### Example

- 'Beoing Corporation','Boeing Company'
- S = {'Beoing', 'Corporation}, T = {'Boeing', Company'}
- tc(S,T) = 0.97 (unit weights for terms)
  - sum of
    - edit('Beoing','Boeing') = 2/6 (normalized)
    - edit('Corporation',Company') = 7/11

#### **Fuzzy Match Similarity**

- $W(S) = sum of W(s,S) for all s \in S$
- fms = 1-min((tc(S,T)/W(S),1))
- Approximating fms:
  - For  $s \in S$  let QG(s) set of qgrams of s
  - $d = (1-1/q) \frac{1}{W(S)} \sum_{s \in S} W(s, S) * \max_{t \in T} (\frac{2}{q} sim_{mh}(QG(s), QG(t)) + d)$
  - For suitable  $\delta$ ,  $\epsilon$  and size of min hash signature
    - $E(fms^{apx}(S,T)) \ge fms(S,T)$
    - $P(fms^{apx}(S,T) \le (1-\delta)fms(S,T)) \le \varepsilon$

#### Multi-attribute similarity measures

- Weighted sum of per attribute similarity
- Application of voting theory
- Rules (more of this later)

#### Voting theory application [GKMS04]

- Relations R with n attributes.
- In principle can apply a different similarity function for each pair of attributes into consideration.
- N orders of the relation tuples, ranked by a similarity score to a query.

# Voting Theory

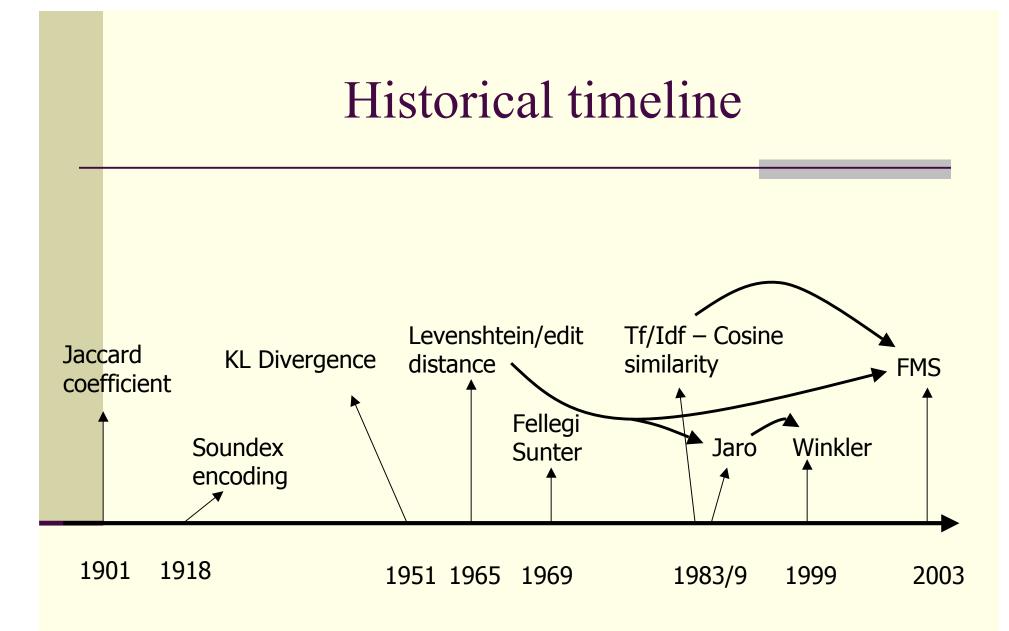
Tuple id T1 T2 T3 T4 T5	custname John smith Josh Smith Nicolas Smith Joseph Smith Jack Smith	address 800 Mountain Av spr 100 Mount Av Spring 800 spring Av Union 555 Mt. Road Spring 100 Springhill lake P	field 8,8 11,11 field 9,9
Query: John sm	hith 100 Mou	nt Rd. Springfield	5.1,5.1
custname	address	locati	on
T1 (1.0) T2 (0.8) T5 (0.7) T4 (0.6) T3 (0.4)	T2 (0.95 T1 (0.8) T4 (0.75 T3 (0.3) T5 (0.1)	T5 (0.	9) 7) 6)

#### Voting theory application

- Merge rankings to obtain a consensus
- Foot-rule distance
  - Let S,T orderings of the same domain D
  - S(i) (T(i)) the order position of the i-th element of D in S (T)

• 
$$F(S,T) = \sum_{i \in D} |S(i) - T(i)|$$

Generalized to distance between S and T1,..Tn F(S,T1,..Tn) =  $\sum_{j=1}^{n} F(S,Tj)$ 



9/23/06

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# Learning similarity functions

- Per attribute
  - Term based (vector space)
  - Edit based
    - Learning constants in character-level distance measures like levenshtein distances
    - Useful for short strings with systematic errors (e.g., OCRs) or domain specific error (e.g.,st., street)
- Multi-attribute records
  - Useful when relative importance of match along different attributes highly domain dependent
  - Example: comparison shopping website
    - Match on title more indicative in books than on electronics
    - Difference in price less indicative in books than electronics

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# Learning Distance Metrics [ST03]

- Learning a distance metric from relative comparisons:
  - A is closer to B than A is to C, etc

• 
$$d_{(A,W)}(x-y) = \sqrt{(x-y)^T A W A^T (x-y)}$$

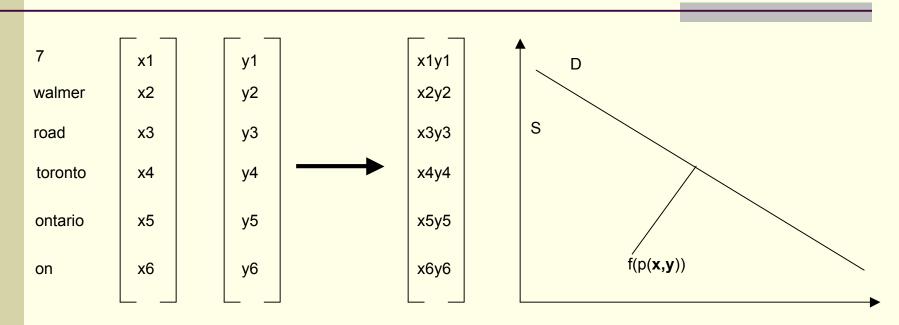
- A can be a real matrix: corresponds to a linear transform of the input
- W a diagonal matrix with non-negative entries (guarantees d is a distance metric)
- Learn entries of W such that to minimize training error
- Zero training error:
  - $\forall$  (i,j,k)  $\varepsilon$  Training set:  $d_{(A,W)}(x_i,x_k)-d_{(A,W)}(x_i,x_k) > 0$
- Select A,W such that d remains as close to an un-weighted euclidean metric as possible.

#### Learnable Vector Space Similarity

- Generic vector space similarity via tfidf
  - Tokens '11th' and 'square' in a list of addresses might have same IDF values
  - Addresses on same street more relevant than addresses on a square...
  - Can we make the distinction?
- Vectors  $\mathbf{x}, \mathbf{y}$ , Sim $(\mathbf{x}, \mathbf{y}) = \sum_{i=1}^{d} \frac{\mathbf{X}_{i} \mathbf{Y}_{i}}{||\mathbf{X}||||\mathbf{Y}||}$ Training data:
- Training data:

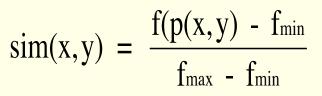
 $S = \{(x,y): x \text{ similar } y\}, D = \{(x,y) x \text{ different } y\}$ 

#### Learnable Vector Space Similarity



P(**x**,**y**)

7 walmer road toronto ontario7 walmer road toronto on

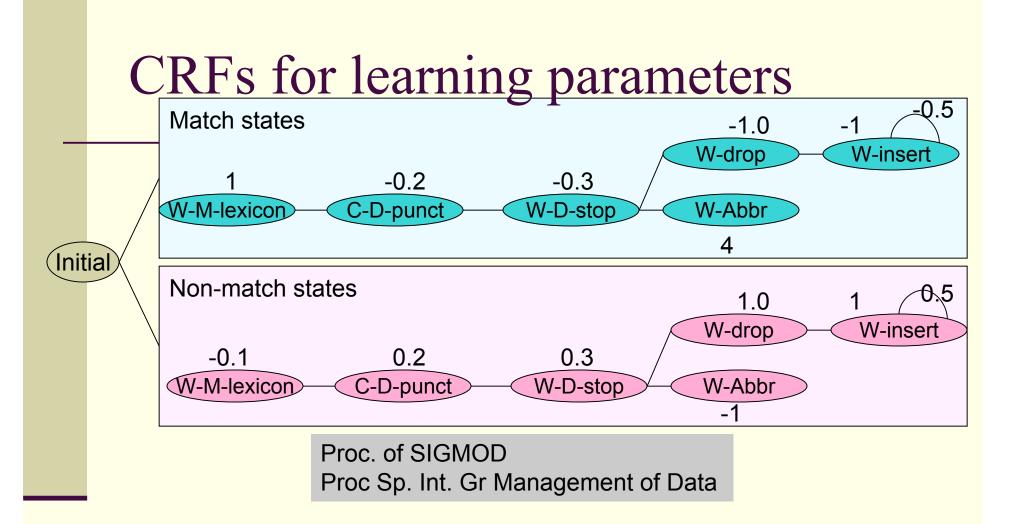


## Learning edit distance parameters

- Free to set relative weights of operations
- May learn weights from input [RY97] using an EM approach.
  - Input: similar pairs
  - Parameters: probability of edit operations
  - E: highest probability edit sequence
  - M: re-estimate probabilities using expectations of the E step
    - Pros: FSM representation (generative model)
    - Cons: fails to incorporate negative examples
- [BM03] extend to learn weights of edit operations with affine gaps
- [MBP05] use CRF approach (incorporates positive and negative input)

# Learning edit parameters using CRFs

- Sequence of edit operations
  - Standard character-level: Insert, Delete, Substitute
    - Costs depends on type: alphabet, number, punctuation
  - Word-level: Insert, Delete, Match, Abbreviation
    - Varying costs: stop words (Eg: The), lexicons (Eg: Corporation, Road)
- Given: examples of duplicate and non-duplicate strings
- Learner: Conditional Random Field
  - Allows for flexible overlapping feature sets
    - Ends with a dot and appears in a dictionary
  - Discriminative training ~ higher accuracy than earlier generative models



- State and transition parameters for match and non-match states
- Multiple paths through states summed over for each pair
- EM-like algorithm for training. 9/23/06

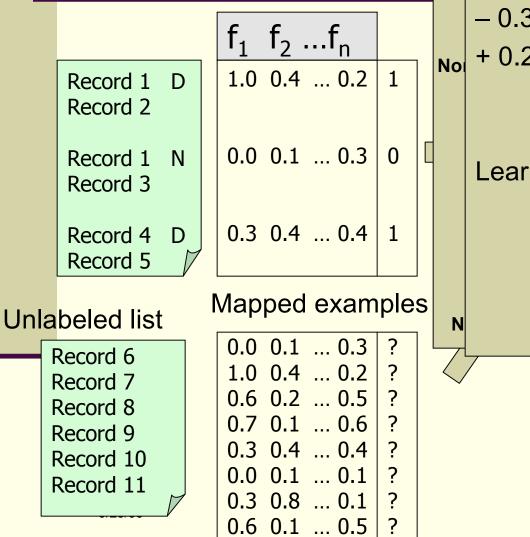
#### Results Citations Distance Metric Restaurant address Face Restaurant name Reasoning 0.952Edit Distance 0.290Earlier generative 0.6860.927Learned Edit Distance 0.3540.7120.9380.966approach (BM03) 0.3650.3800.8970.922Vector-space Word-level only, 0.875Learned Vector-space 0.4330.5320.924no order CRF Edit Distance 0.4480.7830.9180.964Initialized with manual weights (McCallum, Bellare, Pereira EMNLP 2005)

- Edit-distance is better than word-level measures
- CRFs trained with both duplicates and non-duplicates better than generative approaches using only duplicates
- Learning domain-specific edit distances could lead to higher accuracy than manually tuned weights

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#### Multi Attribute Similarity



All-Ngrams\*0.4 + AuthorTitleNgram\*0.2 – 0.3YearDifference + 1.0\*AuthorEditDis No + 0.2\*PageMatch – 3 > 0

#### Learners:

Support Vector Machines (SVM) Logistic regression, Linear regression,

Perceptron

## Learning approach

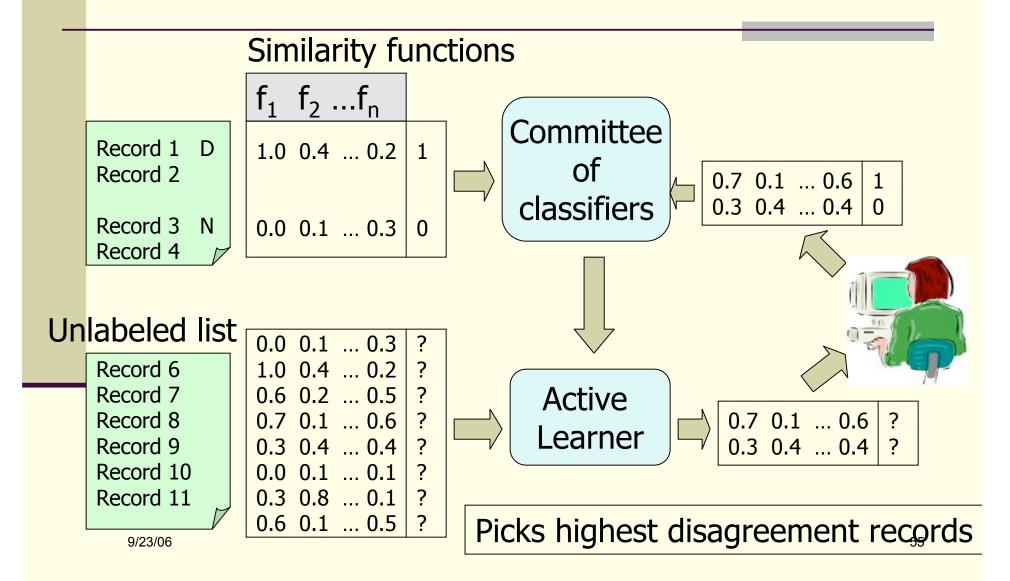
#### Learners used:

- SVMs: high accuracy with limited data,
- Decision trees:interpretable, efficient to apply
- Perceptrons: efficient incremental training (Bilenko et al 2005, Comparison shopping)
- Results:
  - Learnt combination methods better than both
    - Averaging of attribute-level similarities
    - String based methods like edit distance (Bilenko et al 2003)
  - Downside
    - Creating meaningful training data a huge effort

# Training data for learning approach

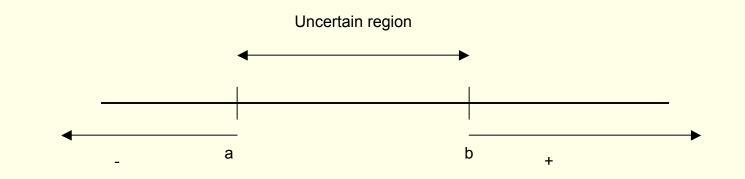
- Heavy manual search in preparing training data
  - Hard to spot challenging/covering duplicates in large lists
  - Even harder to find close non-duplicates that will capture the nuances
  - Need to seek out rare forms of errors in data
- A solution from machine learning Active learning
  - Given
    - Lots of unlabeled data pairs of records
    - Limited labeled data
  - Find examples most informative for classification
    - Highest uncertainty of classification from current data

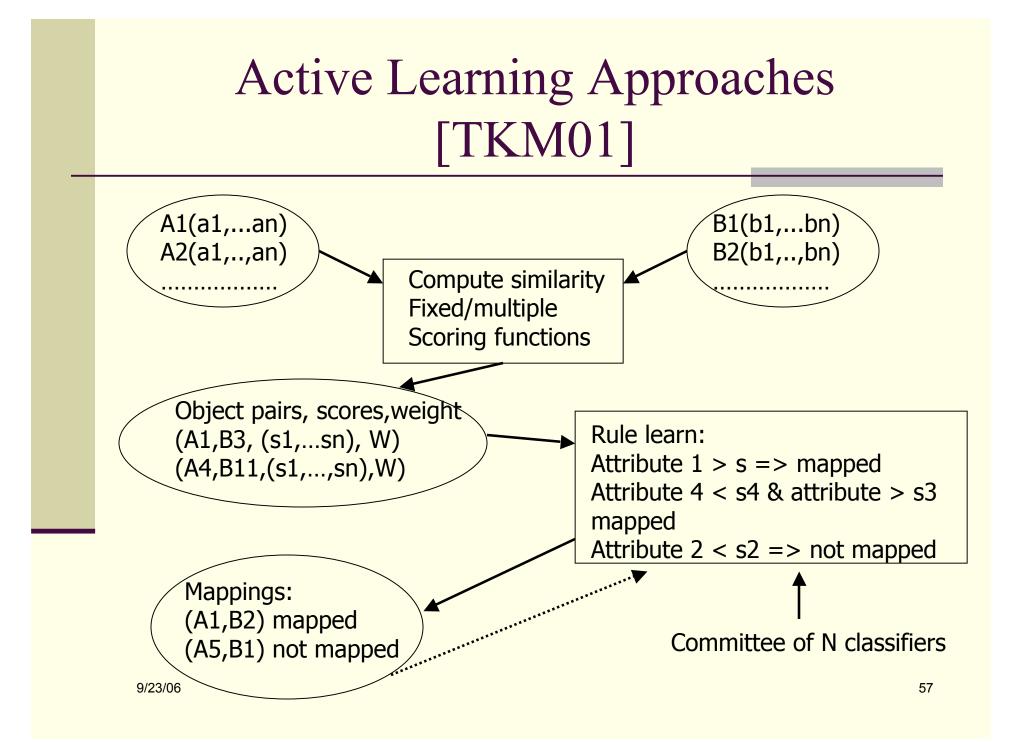
## The active learning approach



# Active Learning [SB02]

- Learn a 'similarity' function (classifier) from labeled data
- Small set of labeled data (pos,neg) and unlabeled data
- Seek instances that when labeled will strengthen the classification process
- Initial classifier sure about prediction on some unlabeled instances and unsure about others (confusion region)
  - Seek predictors on uncertain instances

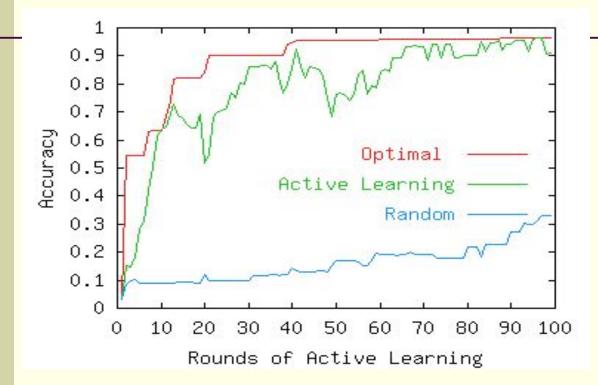




## Active learning algorithm

- Train k classifiers C1, C2,.. Ck on training data through
  - Data resampling,
  - Classifier perturbation
- For each unlabeled instance x
  - Find prediction y1,..., yk from the k classifiers
  - Compute uncertainty U(x) as entropy of above y-s
- Pick instance with highest uncertainty

#### Benefits of active learning



- Active learning much better than random
  - With only 100 active instances
    - 97% accuracy, Random only 30%
- Committee-based selection close to optimal

9/23/06

# Learning: beyond paired 0/1 classification

- Exploiting monotonicity between attribute similarity and class label to learn better
  - A Hierarchical Graphical Model for Record Linkage (Ravikumar, Cohen, UAI 2004)
- Exploiting transitivity to learn on groups
  - T. Finley and T. Joachims, Supervised Clustering with Support Vector Machines, Proceedings of the International Conference on Machine Learning (ICML), 2005.

#### Outline

- Part I: Motivation, similarity measures (90 min)
  - Data quality, applications
  - Linkage methodology, core measures
  - Learning core measures
  - Linkage based measures
- Part II: Efficient algorithms for approximate join (60 min)
- Part III: Clustering algorithms (30 min)

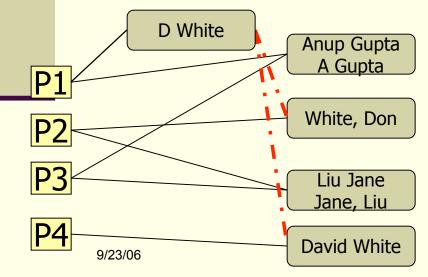
## Similarity based on linkage pattern

P1	D White, A Gupta	
P2	Liu, Jane & White, Don	
P3	Anup Gupta and Liu Jane	
P4	David White	

Relate D White and Don White through the third paper

Path in graph makes D White more similar to Don White than David White

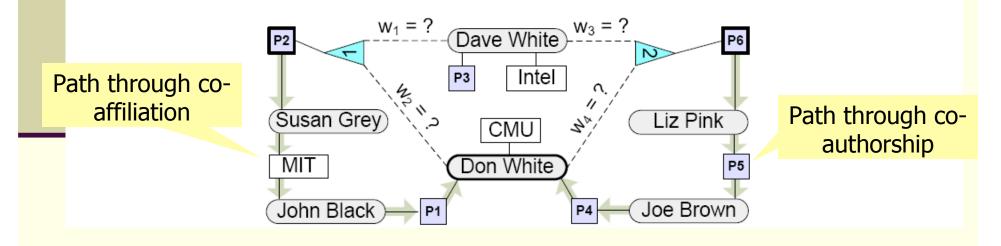
Lots of work on node similarities in graph
sim-rank, conductance models, etc ReIDC (Kalashnikov et al 2006)



#### RelDC: Example with multiple entity types

- $\langle A_1, \text{`Dave White', 'Intel'} \rangle$  $\langle A_2, \text{`Don White', 'CMU'} \rangle$  $\langle A_3, \text{`Susan Grey', 'MIT'} \rangle$  $\langle A_4, \text{`John Black', 'MIT'} \rangle$  $\langle A_5, \text{`Joe Brown', unknown} \rangle$  $\langle A_6, \text{`Liz Pink', unknown} \rangle$
- $\langle P_1, \text{`Databases ...', 'John Black', 'Don White'} \rangle$  $\langle P_2, \text{`Multimedia ...', 'Sue Grey', 'D. White'} \rangle$  $\langle P_3, \text{`Title3 ...', 'Dave White'} \rangle$  $\langle P_4, \text{`Title5 ...', 'Don White', 'Joe Brown'} \rangle$  $\langle P_5, \text{`Title6 ...', 'Joe Brown', 'Liz Pink'} \rangle$  $\langle P_6, \text{`Title7 ...', 'Liz Pink', 'D. White'} \rangle$

#### Task: resolve author references in papers to author table

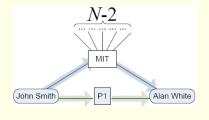


(From: Kalashninov et al 2006)

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# Quantifying strength of connection

- Given a graph G with edges denoting node similarity or some form of relationship, find connection strength between any two nodes u, v Methods
  - Simple methods: shortest path length or flow
    - Fails for high-degree nodes
  - Diffusion kernels
    - $\sum_k \lambda^k \sum_{\mathsf{path}_{Pk}\mathsf{-}\mathsf{long}} \prod_{edge(i,j)\in P} s(i,j)$
  - Electric circuit conductance model (Faloutsos et. al. 2004)
  - Walk-based model (WM)
    - Probabilistic
      - Treat edge weights as probability of transitioning out of node
      - Probability of reaching u from v via random walks
    - SimRank (Jeh&Widom 2002)
      - Expected distance to first meet of random walks from u and v
- Rel<sup>2/23/06</sup> extends (WM) to work for graphs with mutually exclusive choice node<sup>64</sup>



#### RelDC

- Resolve whatever is possible via textual similarity alone
- Create relationship graph with unresolved references connected via choice nodes to options
  - Weights of options related to similarity
- Find connection strength between each unresolved reference to options, resolve to strongest of these
- Results
  - Authors: Author names, affiliation (HP Search)
  - Papers: Titles and Author names (Citeseer)
  - 13% ambiguous references (cannot be resolved via text alone)
  - 100% accuracy on 50 random tests

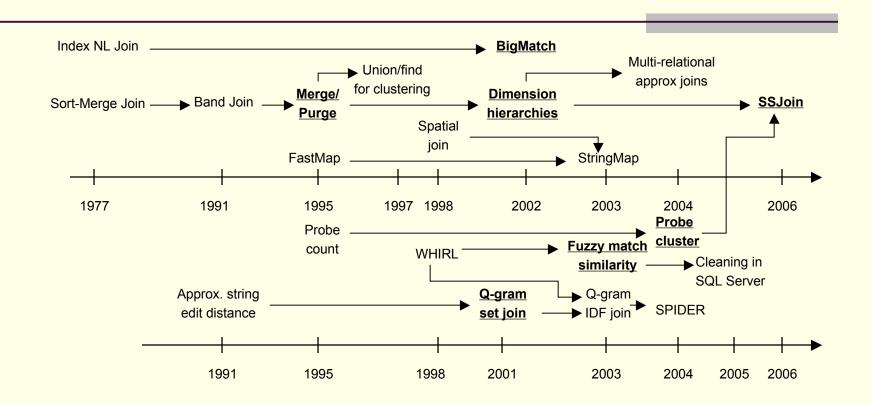
#### Outline

- Part I: Motivation, similarity measures (90 min)
- Part II: Efficient algorithms for approximate join (60 min)
  - Use traditional join methods
  - Extend traditional join methods
  - Commercial systems
- Part III: Clustering algorithms (30 min)

#### Approximate Joins: Baseline + Goal

- An **approximate join** of  $R_1(A_1, ..., A_n)$  and  $R_2(B_1, ..., B_m)$  is
  - A subset of the cartesian product of R<sub>1</sub> and R<sub>2</sub>
  - "Matching" specified attributes A<sub>i1</sub>, ..., A<sub>ik</sub> with B<sub>i1</sub>, ..., B<sub>ik</sub>
  - Labeled with a similarity score > t > 0
- Naïve method: for each record pair, compute similarity score
   I/O and CPU intensive, not scalable to millions of records
- Goal: reduce O(n<sup>2</sup>) cost to O(n\*w), where w << n</p>
  - Reduce number of pairs on which similarity is computed
  - Take advantage of efficient relational join methods

#### Historical Timelines



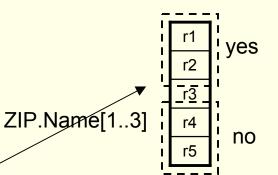
#### Sorted Neighborhood Method [HS95]

- Goal: bring matching records close to each other in linear list
- Background: duplicate elimination [BD83], band join [DNS91]
- Methodology: domain-specific, arbitrary similarity
  - Compute discriminating key per record, sort records
  - Slide fixed size window through sorted list, match in window
  - Use OPS5 rules (equational theory) to determine match
  - Multiple passes with small windows, based on distinct keys
- Lesson: multiple "cheap" passes faster than an "expensive" one

#### Sorted Neighborhood Method [HS95]

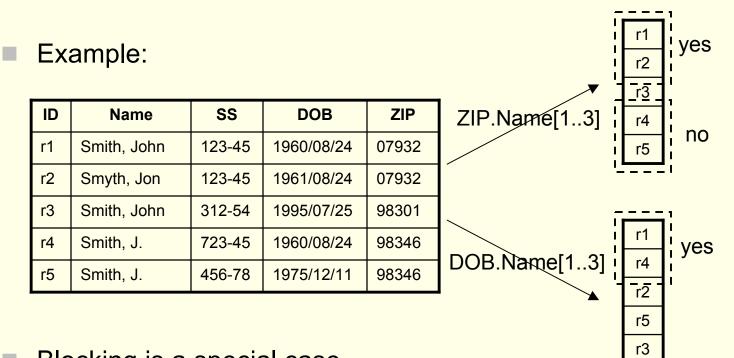
- Goal: bring matching records close to each other in linear list
- Example:

ID	Name	SS	DOB	ZIP
r1	Smith, John	123-45	1960/08/24	07932
r2	Smyth, Jon	123-45	1961/08/24	07932
r3	Smith, John	312-54	1995/07/25	98301
r4	Smith, J.	723-45	1960/08/24	98346
r5	Smith, J.	456-78	1975/12/11	98346



#### Sorted Neighborhood Method [HS95]

Goal: bring matching records close to each other in linear list



Blocking is a special case

# BigMatch [Y02]

- Goal: block/index matching records, based on multiple keys
- Background: indexed nested loop join [BE77]
- Methodology: domain-specific, Jaro-Winkler similarity
  - Store smaller table (100M) in main memory (4GB)
  - Create indexes for each set of grouping/blocking criteria
  - Scan larger table (4B), repeatedly probe smaller table
  - Avoids multiple matches of the same pair
- Lesson: traditional join technique can speed up approximate join

# BigMatch [Y02]

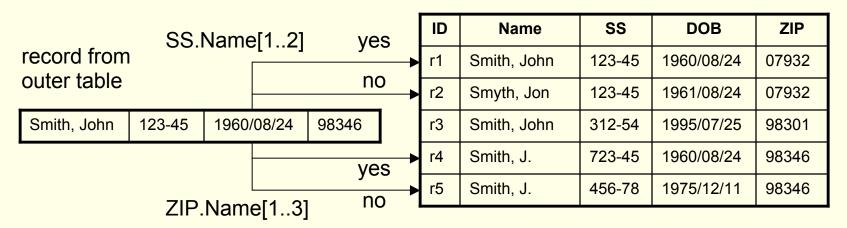
- Goal: block/index matching records, based on multiple keys
- Example:

	record from		[12] yes		ID	Name	SS	DOB	ZIP
					r1	Smith, John	123-45	1960/08/24	07932
outer table			nc	) •	r2	Smyth, Jon	123-45	1961/08/24	07932
Smith, John	123-45	1960/08/24	98346		r3	Smith, John	312-54	1995/07/25	98301
				-	r4	Smith, J.	723-45	1960/08/24	98346
					r5	Smith, J.	456-78	1975/12/11	98346

inner table

# BigMatch [Y02]

- Goal: block/index matching records, based on multiple keys
- Example:



inner table

Avoids multiple matches of the same pair

- Goal: exploit dimension hierarchies for duplicate elimination
- Background: clustering categorical data [GKR98]
- Methodology: domain-independent, structure+text similarity
  - Use hierarchical grouping, instead of sorting, to focus search
  - "Structural" similarity based on overlap of children sets
  - Textual similarity based on weighted token set containment
  - Top-down processing of dimension hierarchy for efficiency
- Lesson: useful to consider group structure in addition to content

Goal: exploit dimension hierarchies for duplicate elimination

ΑΙ	Address	CI	CI	City	SI	SI	State	YI	YI	Country
a1	10 Mountain Avenue	c1	c1	Summit	s1	s1	NJ	y1	y1	USA
a2	250 McCarter	c2	c2	Newark	s2	s2	New Jersey	y1	y2	United States
a3	250 McCarter Hwy	c3	c3	Newark	s3	s3	NJ	y2	у3	US
a4	10 Mountain	c4	c4	Summit	s4	s4	New Jersey	y2		
a5	10 Mountain Street	c5	c5	Summitt	s5	s5	NJ	у3		

Goal: exploit dimension hierarchies for duplicate elimination

#### Example:

ΑΙ	Address	CI	CI	City	SI	SI	State	YI	YI	Country
a1	10 Mountain Avenue	c1	c1	Summit	s1	s1	NJ	y1	y1	USA
a2	250 McCarter	c2	c2	Newark	s2	s2	New Jersey	y1	y2	United States
a3	250 McCarter Hwy	c3	c3	Newark	s3	s3	NJ	y2	<del>y3</del>	-US
a4	10 Mountain	c4	c4	Summit	s4	s4	New Jersey	y2		
а5	10 Mountain Street	c5	c5	Summitt	s5	s5	NJ	y1		

Textual similarity

Goal: exploit dimension hierarchies for duplicate elimination

#### Example:

AI	Address	CI	CI	City	SI	SI	State	YI	YI	Country
a1	10 Mountain Avenue	c1	c1	Summit	s1	s1	NJ	y1	y1	USA
a2	250 McCarter	c2	c2	Newark	s2	s2	New Jersey	y1	<del>y2</del>	United States
a3	250 McCarter Hwy	c3	c3	Newark	s3	s3	NJ	y1	<del>y3</del>	US
a4	10 Mountain	c4	c4	Summit	s4	s4	New Jersey	y1		
a5	10 Mountain Street	c5	c5	Summitt	s5	s5	NJ	y1		

Structural similarity

Goal: exploit dimension hierarchies for duplicate elimination

ΑΙ	Address	CI	CI	City	SI	SI	State	YI	YI	Country
a1	10 Mountain Avenue	c1	c1	Summit	s1	s1	NJ	y1	y1	USA
a2	250 McCarter	c2	c2	Newark	s2	s2	New Jersey	y1	<del>y2</del>	United States
a3	250 McCarter Hwy	c3	c3	Newark	s1	<del>-53</del>	NJ	-y1-	<del>y3</del>	US
a4	10 Mountain	c4	c4	Summit	s2	<del>-s</del> 4	New Jersey	-y1-		
а5	10 Mountain Street	c5	c5	Summitt	s1	<del>\$5</del>	NJ	-y1-		

Goal: exploit dimension hierarchies for duplicate elimination

ΑΙ	Address	CI	CI	City	SI	SI	State	YI	YI	Country
a1	10 Mountain Avenue	c1	c1	Summit	s1	s1	NJ	y1	y1	USA
a2	250 McCarter	c2	c2	Newark	s1	<del>s2</del>	New Jersey	-y1-	<del>y2</del>	United States
a3	250 McCarter Hwy	c3	c3	Newark	s1	<del>-\$3</del>	NJ	-y1-	<del>y3</del>	US
a4	10 Mountain	c4	c4	Summit	s1	<del>-s</del> 4	New Jersey	-y1-		
а5	10 Mountain Street	c5	c5	Summitt	s1	<del>-\$5</del>	NJ	-y1-		

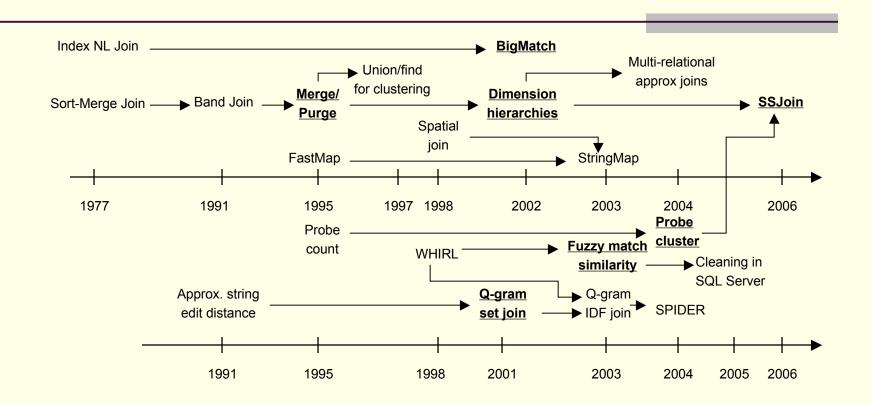
Goal: exploit dimension hierarchies for duplicate elimination

ΑΙ	Address	CI	CI	City	SI	SI	State	YI	YI	Country
a1	10 Mountain Avenue	c1	c1	Summit	s1	s1	NJ	y1	y1	USA
a2	250 McCarter	c2	c2	Newark	s1	<del>s2</del>	New Jersey	-y1-	<del>y2</del>	United States
a3	250 McCarter Hwy	c2	<del>-c3</del>	Newark	<del>-s1</del> -	<del>-\$3</del>	NJ	-y1-	<del>y3</del>	US
a4	10 Mountain	c1	<del>c</del> 4	Summit	<del>-s1</del> -	<del>-s4</del>	New Jersey	-y1-		
а5	10 Mountain Street	c1	<del>-c5</del>	Summitt	<u>-</u> \$1	<del>\$5</del>	NJ	-y1-		

Goal: exploit dimension hierarchies for duplicate elimination

ΑΙ	Address	CI	CI	City	SI	SI	State	YI	YI	Country
a1	10 Mountain Avenue	c1	c1	Summit	s1	s1	NJ	y1	y1	USA
<del>a2</del>	250 McCarter	c2	c2	Newark	s1	<del>s2</del>	New Jersey	y1	<del>y2</del>	United States
a3	250 McCarter Hwy	c2	<del>-c3</del>	Newark	<del>-s1</del> -	<del>-\$3</del>	NJ	y1	<del>y3</del>	US
a4	10 Mountain	c1	<del>c4</del>	Summit	<del>-s1</del> -	<del>s</del> 4	New Jersey	y1		
<del>-a5</del>	10 Mountain Street	c1	<del>-C5</del>	Summitt	<u>-</u> \$1	<del>-s5</del>	NJ	y1		

### Historical Timelines



- Goal: compute thresholded edit distance join on string attributes
- Background: combinatorial pattern matching [JU91]
- Methodology: domain-independent, edit distance similarity
  - Extract set of all overlapping q-grams Q(s) from string s
  - $ED(s_1,s_2) \le d \rightarrow |Q(s_1) \cap Q(s_2)| \ge max(|s_1|,|s_2|) (d-1)^*q 1$
  - Cheap filters (length, count, position) to prune non-matches
  - Pure SQL solution: cost-based join methods
- Lesson: reduce approximate join to aggregated set intersection

Goal: compute thresholded edit distance join on string attributes

ID	Name
r1	Srivastava
r2	Shrivastava
r3	Shrivastav

Goal: compute thresholded edit distance join on string attributes

ID	Name	3-grams
r1	Srivastava	##s, <mark>#sr</mark> , sri, riv, iva, vas, ast, sta, tav, ava, va\$, a\$\$
r2	Shrivastava	##s, <mark>#sh, shr, hr</mark> i, riv, iva, vas, ast, sta, tav, ava, va\$, a\$\$
r3	Shrivastav	

- $ED(s_1,s_2) \le d \Rightarrow |Q(s_1) \cap Q(s_2)| \ge max(|s_1|,|s_2|) (d-1)^*q 1$
- ED(r1, r2) = 1,  $|Q(r1) \cap Q(r2)| = 10$

Goal: compute thresholded edit distance join on string attributes

ID	Name	3-grams
r1	Srivastava	##s, <mark>#sr</mark> , sri, riv, iva, vas, ast, sta, tav, <mark>ava, va\$, a\$\$</mark>
r2	Shrivastava	
r3	Shrivastav	##s, <mark>#sh, shr, hri,</mark> riv, iva, vas, ast, sta, tav, <mark>av\$</mark> , v <mark>\$\$</mark>

- $ED(s_1, s_2) \le d \rightarrow |Q(s_1) \cap Q(s_2)| \ge max(|s_1|, |s_2|) (d-1)^*q 1$
- ED(r1, r2) = 2,  $|Q(r1) \cap Q(r2)| = 7$

Goal: compute thresholded edit distance join on string attributes

Q

ID	Name
r1	Srivastava
r2	Shrivastava
r3	Shrivastav

ID	Qg	
r1	##s	
r1	#sr	
r1	sri	
r1	riv	
r1	iva	
r1	vas	
r1	ast	
r1	sta	
r1	tav	
r1	ava	
r1	va\$	
r1	a\$\$	

ID	Qg
r3	##s
r3	#sh
r3	shr
r3	hri
r3	riv
r3	iva
r3	vas
r3	ast
r3	sta
r3	tav
r3	av\$
r3	v\$\$

Goal: compute thresholded edit distance join on string attributes

Q

#### Example:

ID	Name
r1	Srivastava
r2	Shrivastava
r3	Shrivastav

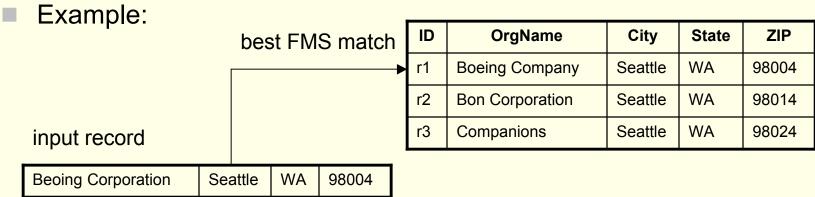
SELECT Q1.ID, Q2.ID FROM Q AS Q1, Q AS Q2 WHERE Q1.Qg = Q2.Qg GROUP BY Q1.ID, Q2.ID HAVING COUNT(\*) > T

ID	Qg	ID	Qg
r1	##s	r3	##s
r1	#sr	r3	#sh
r1	sri	r3	shr
r1	riv	r3	hri
r1	iva	r3	riv
r1	vas	r3	iva
r1	ast	r3	vas
r1	sta	r3	ast
r1	tav	r3	sta
r1	ava	r3	tav
r1	va\$	r3	av\$
r1	a\$\$	r3	v\$\$

- Goal: identify K "closest" reference records in on-line setting
- Background: IDF weighted cosine similarity, WHIRL [C98]
- Methodology: domain-independent, IDF+ED similarity
  - Similarity metric based on IDF weighted token edit distance
  - Approximate similarity metric using Jaccard on q-gram sets
  - Small error tolerant index table, sharing of minhash q-grams
  - Optimistic short circuiting exploits large token IDF weights
- Lesson: IDF weighting useful to capture erroneous tokens

- Goal: identify K "closest" reference records in on-line setting
  - ID OrgName City State ZIP r1 **Boeing Company** Seattle WA 98004 best ED match **Bon Corporation** 98014 r2 Seattle WA r3 Companions Seattle WA 98024 input record **Beoing Corporation** Seattle WA 98004
- Example:

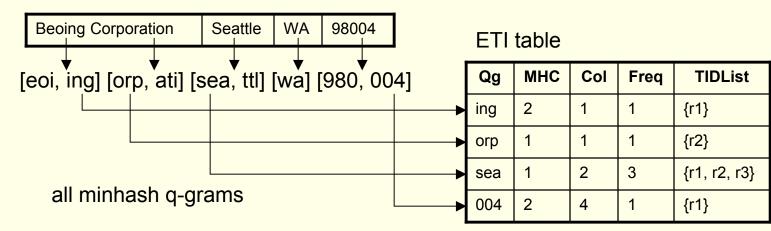
Goal: identify K "closest" reference records in on-line setting



- Goal: identify K "closest" reference records in on-line setting
- Example:

input record

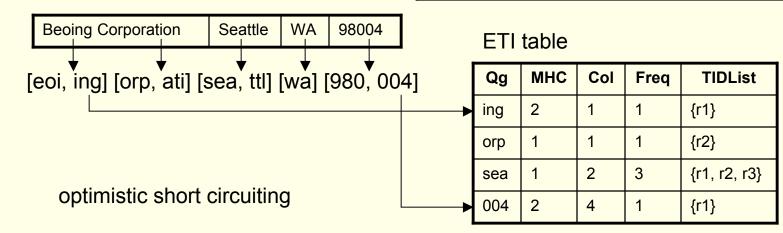
ID	OrgName	City	State	ZIP
r1	Boeing Company	Seattle	WA	98004
r2	Bon Corporation	Seattle	WA	98014
r3	Companions	Seattle	WA	98024



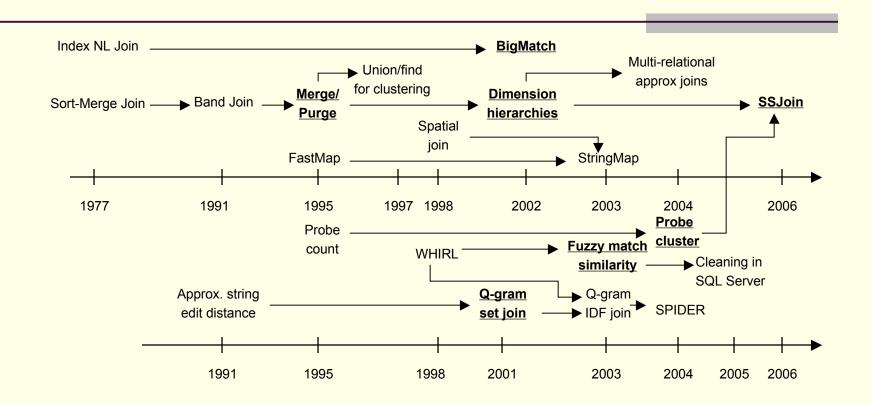
- Goal: identify K "closest" reference records in on-line setting
- Example:

input record

ID	OrgName	City	State	ZIP
r1	Boeing Company	Seattle	WA	98004
r2	Bon Corporation	Seattle	WA	98014
r3	Companions	Seattle	WA	98024



### Historical Timelines



- Goal: generic algorithm for set join based on similarity predicate
- Background: IR and probe count using inverted index [TF95]
- Methodology: domain-independent, weighted set similarity
  - Map a string to a set of elements (words, q-grams, etc.)
  - Build inverted lists on individual set elements
  - Optimization: process skewed lists in increasing size order
  - Optimization: sort lists in decreasing order of record sizes
- Lesson: IR query optimizations useful for approximate joins

Goal: generic algorithm for set join based on similarity predicate

Inverted index

ID	SVA
r1	{##s, #sr, sri, riv, iva, vas, ast, sta, tav, ava, va\$, a\$\$}
r2	{##s, #sh, shr, hri, riv, iva, vas, ast, sta, tav, ava, va\$, a\$\$}
r3	{##s, #sh, shr, hri, riv, iva, vas, ast, sta, tav, av\$, v\$\$}

-	
SE	IDs
##s	r1, r2, r3
#sr	r1
#sh	r2, r3
sri	r1
shr	r2, r3
hri	r2, r3
riv	r1, r2, r3
tav	r1, r2, r3
ava	r1, r2
v\$\$	r3

Goal: generic algorithm for set join based on similarity predicate

Inverted index

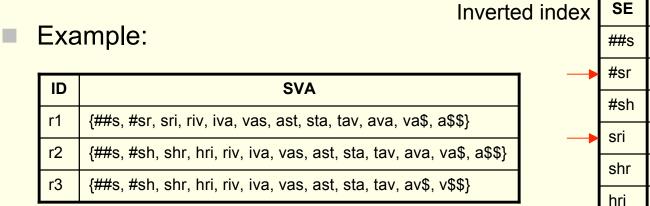
Example:

ID	SVA
r1	{##s, #sr, sri, riv, iva, vas, ast, sta, tav, ava, va\$, a\$\$}
r2	{##s, #sh, shr, hri, riv, iva, vas, ast, sta, tav, ava, va\$, a\$\$}
r3	{##s, #sh, shr, hri, riv, iva, vas, ast, sta, tav, av\$, v\$\$}

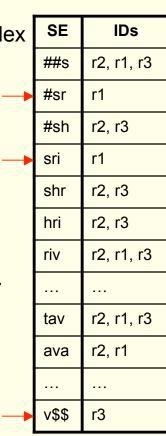
Sort lists in decreasing order of record sizes

-	
SE	IDs
##s	r2, r1, r3
#sr	r1
#sh	r2, r3
sri	r1
shr	r2, r3
hri	r2, r3
riv	r2, r1, r3
tav	r2, r1, r3
ava	r2, r1
v\$\$	r3

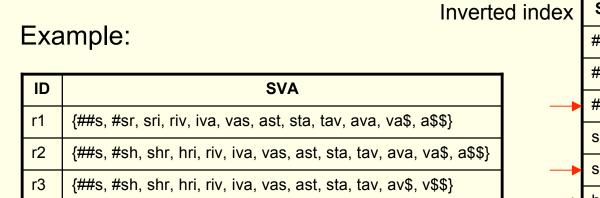
Goal: generic algorithm for set join based on similarity predicate



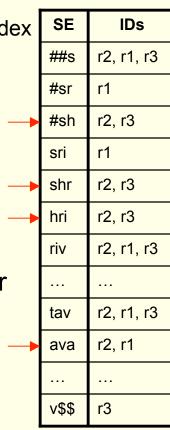
Process skewed lists in increasing size order



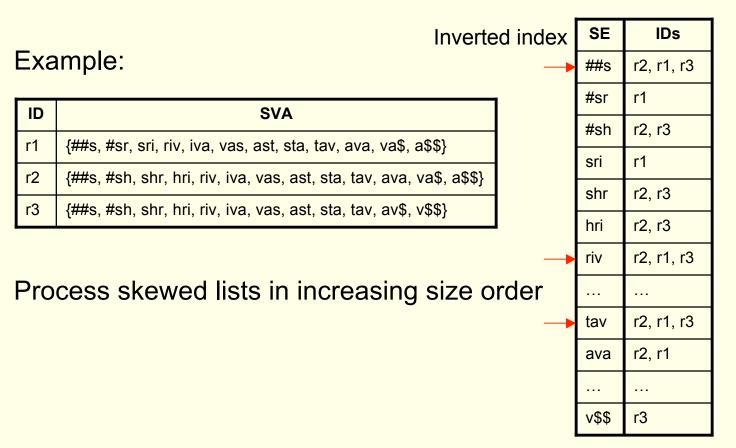
Goal: generic algorithm for set join based on similarity predicate



Process skewed lists in increasing size order



Goal: generic algorithm for set join based on similarity predicate



- Goal: generic algorithm for set join based on similarity predicate
- Background: Probe-Cluster, dimension hierarchies, q-gram join
- Methodology: domain-independent, weighted set similarity
  - Compare strings based on sets associated with each string
  - Problem:  $Overlap(s1, s2) \ge threshold$
  - Optimization: high set overlap  $\rightarrow$  overlap of ordered subsets
  - SQL implementation using equijoins, cost-based plans
- Lesson: Generic algorithms can be supported in DBMS

Goal: generic algorithm for set join based on similarity predicate

Q

#### Example:

ID	Name
r1	Srivastava
r4	Srivastav

SELECT Q1.ID, Q2.ID FROM Q AS Q1, Q AS Q2 WHERE Q1.Qg = Q2.Qg GROUP BY Q1.ID, Q2.ID HAVING COUNT(\*) > 8

		 _	
ID	Qg	ID	Qg
r1	##s	r4	##s
r1	#sr	r4	#sr
r1	sri	r4	sri
r1	riv	r4	riv
r1	iva	r4	iva
r1	vas	r4	vas
r1	ast	r4	ast
r1	sta	r4	sta
r1	tav	r4	tav
r1	ava	r4	av\$
r1	va\$	r4	v\$\$
r1	a\$\$		

- Goal: generic algorithm for set join based on similarity predicate
- Example:

ID	Name
r1	Srivastava
r4	Srivastav

SELECT Q1.ID, Q2.ID FROM Q AS Q1, Q AS Q2 WHERE Q1.Qg = Q2.Qg GROUP BY Q1.ID, Q2.ID HAVING COUNT(\*) > 8

ID	Qg	
r1	tav	
r1	ava	$\setminus$
r1	va\$	
r1	a\$\$	

Q

	ID	Qg
	r4	##s
	r4	#sr
	r4	sri
	r4	riv
	r4	iva
	r4	vas
$\setminus$	r4	ast
	r4	sta
	r4	tav
	r4	av\$
	r4	v\$\$

Optimization: use any 4 q-grams of r1 with all of r4

Goal: generic algorithm for set join based on similarity predicate

Q

#### Example:

ID	Name
r1	Srivastava
r4	Srivastav

SELECT Q1.ID, Q2.ID FROM Q AS Q1, Q AS Q2 WHERE Q1.Qg = Q2.Qg GROUP BY Q1.ID, Q2.ID HAVING COUNT(\*) > 8

Optimization: use any 3 q-grams of r4

ID	Qg	
r1	##s	
r1	#sr	
r1	sri	
r1	riv	
r1	iva	
r1	vas	
r1	ast	
r1	sta	
r1	tav	
r1	ava	
r1	va\$	
r1	a\$\$	

D	Qg
r4	sri
r4	av\$
r4	v\$\$

Goal: generic algorithm for set join based on similarity predicate

Q

Example:

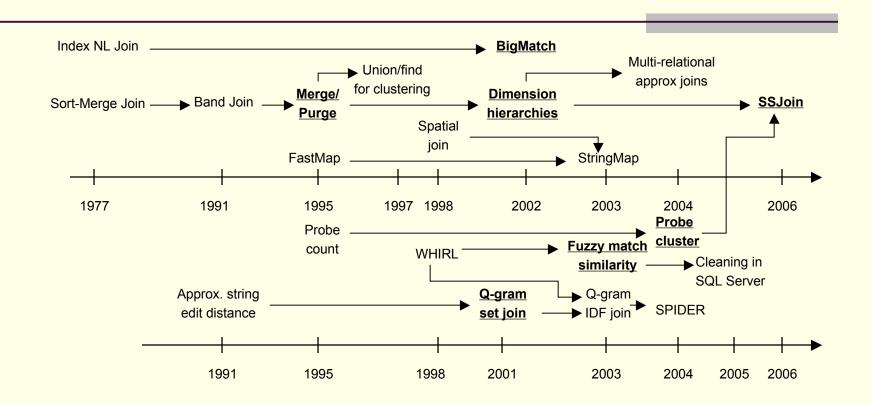
ID	Name	
r1	Srivastava	
r4	Srivastav	

SELECT Q1.ID, Q2.ID FROM Q AS Q1, Q AS Q2 WHERE Q1.Qg = Q2.Qg GROUP BY Q1.ID, Q2.ID HAVING COUNT(\*) > 8

ID	Qg	ID	Qg
r1	iva	r4	iva
r1	ast	r4	ast
r1	ava	r4	av\$
r1	a\$\$		

- Optimization: use ordered 4 q-grams of r1 and 3 q-grams of r4
- Suggested ordering: based on decreasing IDF weights

### Historical Timelines



# Commercial Systems: Comparisons

<u>Commercial</u>	<u>Record Linkage</u>	<u>Distance Metrics</u>	<u>Domain-Specific</u>	Additional Data
<u>System</u>	<u>Methodology</u>	<u>Supported</u>	<u>Matching</u>	Quality Support
SQL Server Integration Services 2005	Fuzzy Lookup; Fuzzy Grouping; uses Error Tolerant Index	customized, domain- independent: edit distance; number, order, freq. of tokens	unknown	unknown
OracleBl Warehouse Builder 10gR2 "Paris"	match-merge rules; deterministic and probabilistic matching	Jaro-Winkler; double metaphone	name & address parse; match; standardize: 3 <sup>rd</sup> party vendors	data profiling; data rules; data auditors
IBM's Entity	probabilistic matching	wide variety of fuzzy matching functions	name recognition;	data profiling;
Analytic	(information content);		identity resolution;	standardization;
Solutions,	multi-pass blocking;		relationship	trends and
QualityStage	rules-based merging		resolution: EAS	anomalies;

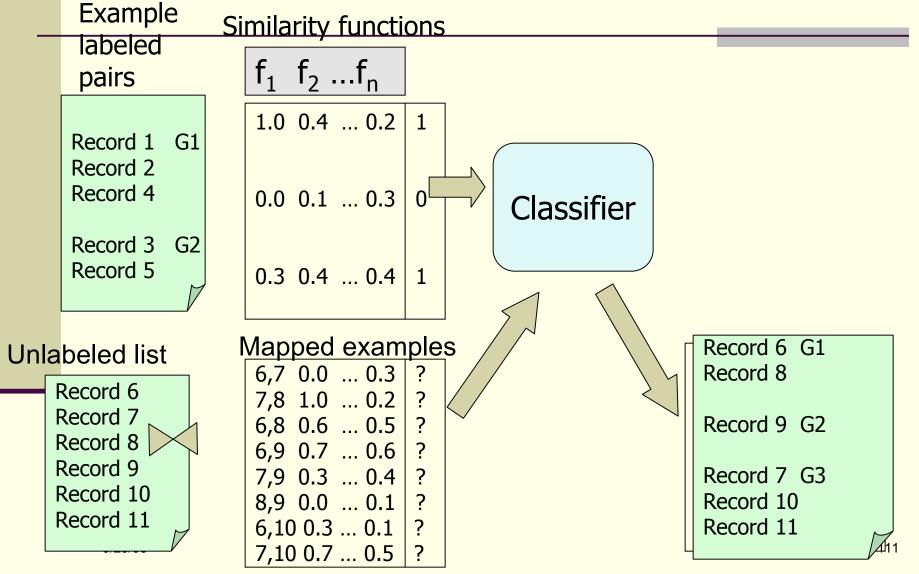
#### Outline

- Part I: Motivation, similarity measures (90 min)
- Part II: Efficient algorithms for approximate join (60 min)
- Part III: Clustering/partitioning algorithms (30 min)

# Partitioning/collective deduplication

- Single-entity types
  - A is same as B if both are same as C.
- Multiple linked entity types
  - If paper A is same as paper B then venue of A is the same as venue of B.

## Partitioning data records



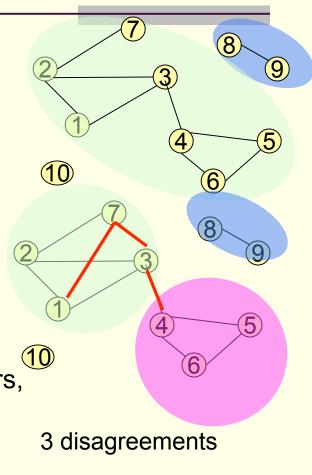
# Creating partitions

#### Transitive closure

 Dangers: unrelated records collapsed into a single cluster

Correlation clustering (Bansal et al 2002)

- Partition to minimize total disagreements
   Edges across partitions
   Missing edges within partition
- More appealing than clustering:
  - No magic constants: number of clusters, similarity thresholds, diameter, etc
- Extends to real-valued scores
- NP Hard: many approximate algorithms



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# Algorithms for correlation clustering

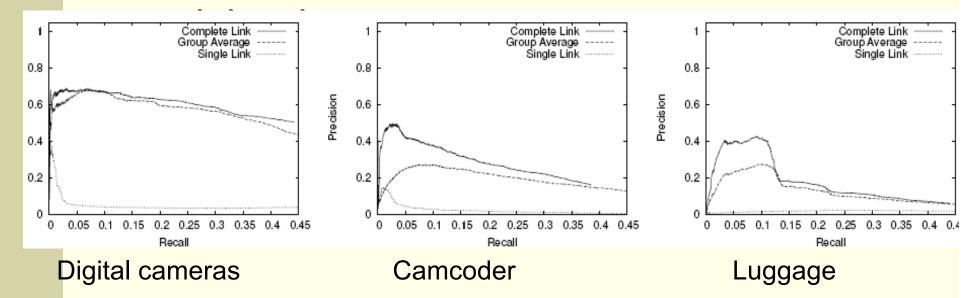
- Integer programming formulation (Charikar 03)
  - $X_{ij} = 1$  if i and j in same partition, 0 otherwise

$$\min \sum_{ij \in edges} (1 - x_{ij}) + \sum_{ij \notin edges} x_{ij}$$

such that  $x_{ij} \in \{0, 1\}$ 

- $x_{ij} + x_{jk} \le 1 + x_{ik}$  (partitioning constraint)
- Impractical: O(n<sup>3</sup>) constraints
- Practical substitutes (Heuristics, no guarantees)
  - Agglomerative clustering: repeatedly merge closest clusters
    - Efficient implementation possible via heaps (BG 2005)
    - Definition of closeness subject to tuning
      - Greatest reduction in error
      - Average/Max/Min similarity

# Empirical results on data



Setup: Online comparison shopping,

- Fields: name, model, description, price
- Learner: Online perceptron learner
- Complete-link clustering >> single-link clustering(transitive closure)
- An issue: when to stop merging clusters

(From: Bilenko et al, 2005)

# Other methods of partitioning

#### [Chaudhuri et al ICDE 2005]

- Partitions are compact and relatively far from other points
- A Partition has to satisfy a number of criteria
- Points within partition closer than any points outside
- #points within p-neighborhood of each partition < c</p>
- Either number of points in partition < K, or diameter <  $\theta$

# Algorithm

Consider case where partitions required to be of size <

- $K \rightarrow$  if partition P<sub>i</sub> of size m in output then
- m-nearest neighbors of all r in P<sub>i</sub> is P<sub>i</sub>
- Neighborhood of each point is sparse

For each record, do efficient index probes to get

Get K nearest neighbors

Count of number of points in p-neighborhood for each m nearest neighbors

Form pairs and perform grouping based on above insight to find groups

## Summary: partitioning

- Transitive closure is a bad idea
- No verdict yet on best alternative
- Difficult to design an objective and algorithms
- Correlation clustering
  - Reasonable objective with a skewed scoring function
  - Poor algorithms
- Greedy agglomerative clustering algorithms ok
  - Greatest minimum similarity (complete-link), benefit
  - Reasonable performance with heap-based implementation
- Dense/Sparse partitioning
  - Positives: Declarative objective, efficient algorithm
  - Parameter retuning across domains
- Need comparison between complete-link, Dense/Sparse, and Correlation clustering.

# Collective de-duplication: multiattribute

	Record	Title a <sup>1</sup>	Author a <sup>2</sup>	Venue a <sup>3</sup>	
	$b_1$	"Record Linkage using CRFs"	"Linda Stewart"	"KDD-2003"	
l	$b_2$	"Record Linkage using CRFs"	"Linda Stewart"	"9th SIGKDD"	
	$b_3$	"Learning Boolean Formulas"	"Bill Johnson"	"KDD-2003"	
	$b_4$	"Learning of Boolean Expressions"	"William Johnson"	"9th SIGKDD" $% \left( {{{\rm{SIGKDD}}}} \right)$	

 $R_{ii}$ 

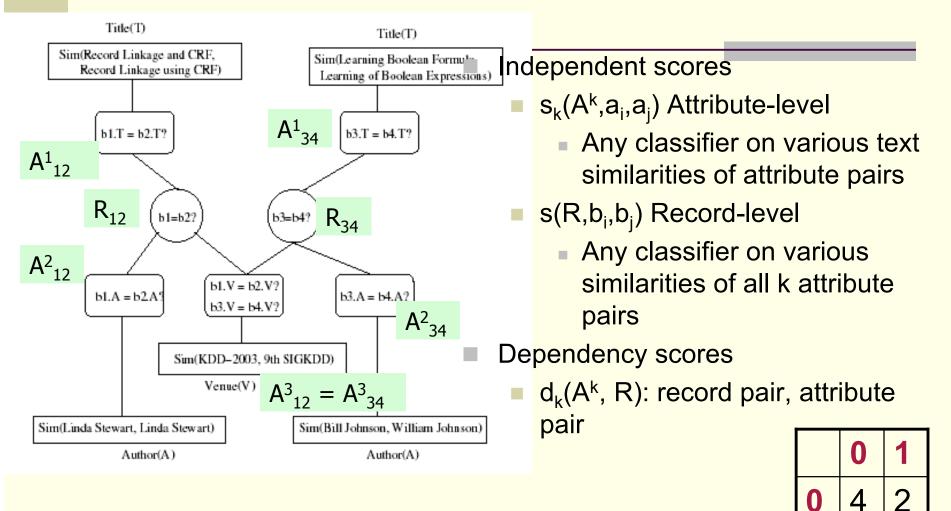
Associate variables for predictions for each attribute k each record pair (i,j)  $A^{k}_{ij}$ 

for each record pair

from Parag & Domingos 2005

## Dependency graph

#### Scoring functions





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### Joint de-duplication steps

 Jointly pick 0/1 labels for all record pairs R<sub>ij</sub> and all K attribute pairs A<sup>k</sup><sub>ij</sub> to maximize

$$\sum_{ij} [s(R_{ij}) + \sum_{k} s_{k}(A_{ij}^{k}) + d_{k}(R_{ij}, A_{ij}^{k})]$$

- When dependency scores associative
  - $d_{k}(1,1) + d_{k}(0,0) \ge d_{k}(1,0) + d_{k}(0,1)$
  - Can find optimal scores through graph MINCUT
- Assigning scores
  - Manually as in Levy et. al
  - Example-based training as in Domingos et al
    - Creates a weighted feature-based log-linear model
      - $s(R_{ij}) = w_1^* sim(a_i^1, a_j^1) + \dots + w_k^* sim(a_i^k, a_j^k)$
    - Very flexible and powerful.

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### Other issues and approaches

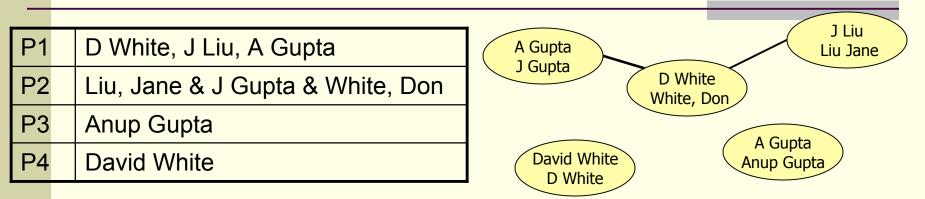
- Partitioning
  - Transitive-closure as a post processing
  - Results:

	Citation		Author		Venue	
	Ρ	Т	Ρ	Т	Ρ	Т
Independent	87	85	79	89	49	59
Collective	86	89	89	89	86	82

- Collective deduplication
  - does not help whole citations,
  - helps attributes
- Transitive closure can cause drop in accuracy
- Combined partitioning and linked dedup
  - Dong, HaLevy, Madhavan (SIGMOD 2005)
  - Bhattacharya and Getoor (2005)

# Collective linkage: set-oriented data

(Bhattacharya and Getoor, 2005)



#### **Scoring functions**

- S(A<sub>ij</sub>) Attribute-level
  - Text similarity
- $S(A_{ij}, N_{ij})$  Dependency with labels of co-author set
  - Fraction of co-author set assigned label 1.
- Final score:
  - $_{3/23/66}$  s(A<sub>ij</sub>) + (1-a) s(A<sub>ij</sub>, N<sub>ij</sub>)
  - a is the only parameter

Greedy agglomerative clustering

Merge author clusters with highest score

**Algorithm** 

- Redefine similarity between clusters of authors instead of single authors
  - Max of author-level similarity

#### Open Problem: Inside or Outside?

- Issue: optimizable processing in a relational database
- Background
  - Declarative data cleaning in AJAX [GFS+01]
  - Q-gram based metrics, SPIDER [GIJ+01,GIKS03,KMS04]
  - SSJoin [CGK06]
  - Compact sets, sparse neighborhood [CGM05]
- Goal: express arbitrary record linkage in SQL

#### **Open Problem: Multi-Table Joins**

- Issue: information in auxiliary tables can aid matching
- Background
  - Hierarchical models [ACG02]
  - Iterative matching [BG04]
  - Graphical models [KMC05]
- Goal: efficient multi-table approximate joins

## Open Problem: Benchmarking

- Issue: many algorithms and similarity measures, no benchmarks
- Background
  - Comparing quality of different similarity measures [CRF03]
- Goal: develop standard benchmarks (queries, data generation)

#### Conclusions

- Record linkage is critical when data quality is poor
  - Similarity metrics
  - Efficient sub-quadratic approximate join algorithms
  - Efficient clustering algorithms
- Wealth of challenging technical problems
  - Sophisticated similarity metrics, massive data sets
  - Important to work with real datasets

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