#### Social and Information Networks

#### University of Toronto CSC303 Winter/Spring 2024

Week 9: Mar 11-15

### This week's learning goals

- Define Mitochondrial Eve
  - ► Define the Wright-Fisher single-parent ancestry model (Ch 21.7)
  - Prove an estimation of time to convergence (Ch 21.8B)
- Define and reason about stability, with respect to bargaining in a Network Exchange Model
  - Explain power in the network exchange social experiment, from structural and non-structural factors (Ch 12.1-12.3)
  - Define stable outcomes (Ch 12.7), and determine if an outcome is stable
  - Define the Ultimatum Game (Ch 12.6), and explain how it compares to bargaining
  - Define balanced outcomes (Ch 12.5, 12.8), and determine if an outcome is balanced

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  - She lived sometime between 100,000 and 200,000 years ago
  - Probably living in Africa
- We'll ignores the issue of the location of Mitochondrial Eve and focuses on the basis (i.e. a model based on various assumptions) for this bold assertion of a common ancestry

**Note:** I suggest reading the text as to the caveats about the model (see Ch 21.7)

• To understand the assertion, we have to make some simplifying biological and mathematical assumptions (see section 21.8 B)

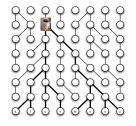
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  - The mathematical assumptions do not change any of the conclusions

#### **Mitochondrial Eve continued**

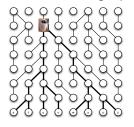
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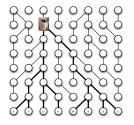


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- The model lets us conclude common mitochondrial DNA ancestry must have originated with a single female Mitochondrial Eve
- The model can also estimate for the time period in which she lived
- This does **not** say that Mitochondrial Eve was the only woman alive at this time, but that our mitochondrial DNA traces back to one woman
- Additionally, our genomic makeup does come from both parents

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  - ► assume a fixed population of *N* individuals throughout the entire period of time
- Problems? Obviously inconsistent with the fact that world population is growing
- Ultimately does not change the nature of the conclusions or even the nature of the analysis
  - In fact, once we accept that populations are growing, it is clear that certain individuals must be having multiple children which is also part of the model

#### Single parent ancestry model continued

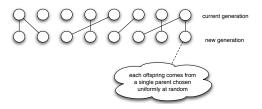
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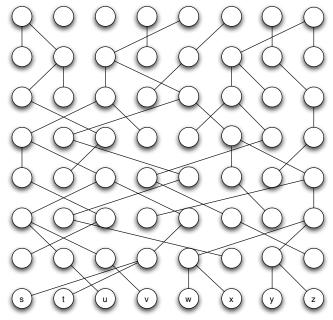
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  - A significant assumption given geography, ethnicity, etc...
  - To reconcile this (with respect to the assertion of a single Mitochondrial Eve), we need to understand the extent to which individual communities can be isolated
    - ★ Ultimately, the timing for when common ancestry would have taken place is not impacted by this assumption

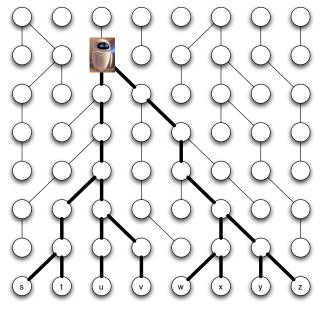


#### Figure: [Fig 21.11, E&K]

#### More generations of the model



## Ancestry depicted.



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Suppose we have a total population of N and at some point of time t + 1 that we are down to k candidates (lineages) for a common ancestor. We want to consider the probability that two lineages will collide so that there be (at most) k - 1 candidates.

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Case: k > 2. Lets consider the probability that none of the k nodes share a parent. There will be no collapsing if the second node doesn't collide with the first, the third doesn't collide with the first two, etc, so this means that the probability of no collapsing is :

$$(1-\frac{1}{N})(1-\frac{2}{N})\cdots(1-\frac{k-1}{N})$$

The previous product

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where g(k) depends only on k and not on N.

For any fixed k, the latter term is relatively negligible and we can say that the probability that none of the k share a parent is  $1 - \frac{k(k-1)}{2N}$ .

Fact: If we have a binary random variable  $Y_k$  (i.e., a heads coin flip) that is true with probability p, then the expected number of independent samples until  $Y_k$  is true (denoted  $E[X_k]$ ) is exactly 1/p

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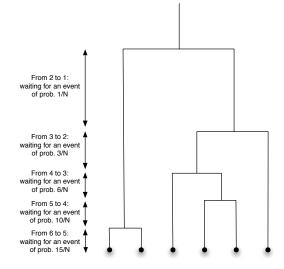
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**Note:** Initially when k is large, the decrease is expected every generation going back. But when k is a small constant, then the expected number of generations to show a decrease is proportional to N.

#### Depiction of the lineages colliding



**Figure:** Assuming no three lineages collide simultaneously. [Fig 21.1(a), E&K]

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Since 
$$\mathbb{E}[X_j] = \frac{2N}{j(j-1)}$$
 and  $\frac{1}{j(j-1)} = \frac{1}{j-1} - \frac{1}{j}$ , by linearity of expectations we have:  
 $\mathbb{E}[X^k] = \sum_{j=1}^k \frac{2N}{j(j-1)}$   
 $= 2N\left(\left[\frac{1}{1} - \frac{1}{2}\right] + \left[\frac{1}{2} - \frac{1}{3}\right] + \dots + \left[\frac{1}{k-1} - \frac{1}{k}\right]\right)$   
 $= 2N\left(1 - \frac{1}{k}\right)$ 

**Note:** Further more detailed analysis is consistent with the basic analysis that was presented in the text.

# Recap

- Mitochondrial Eve
  - Problem setup (Ch 21.7)
  - Wright-Fisher single-parent ancestry model (Ch 21.7)
  - Estimation of time to convergence (Ch 21.8B)

# **Chapter 12: Bargaining and Power in Networks**

- We begin a subtle and fascinating topic: how individuals in a network come to agreement on an outcome!
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  - We have a course (CSC304) which covers game theory; as opposed to necessary minimum we'll be covering
- To ensure we're all on the same page, we'll informally mention some basic concepts to keep in mind
  - We've seen these concepts, at least implicitly, in the course material already :)

 Individuals (agents) have strategies or actions and employ a (pure or mixed/randomized) strategy so as to act in self interest, always trying to maximize benefit or minimize cost

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- Agents are acting in self interest implies that their actions are decentralized
  - Mechanism design concerns how a central agent can introduce incentives to influence agents
  - Aside: An example of a result in Mechanism Design is Gibbard-Satterthwaite theorem, which states that any voting rule is either
    - \* Dictatorial
    - Only selecting the winner from a set of two candidates
    - ★ Susceptible to tactical voting

# Game theory concepts: Equilibrium

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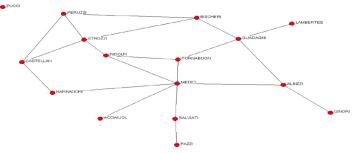


- Appeared in Schelling segregation model in Chapter 4, structural balance in Chapter 5, and will be important in Chapter 12 and the study of relative power
  - we will see them again in stable matchings and traffic equilibria

• Power between individuals can come from two distinct sources:

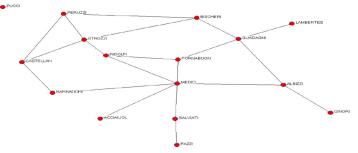
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- ★ In the second week of the course we discussed the *bridging capital* and the *bonding capital* of a node

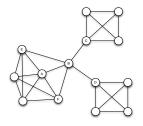
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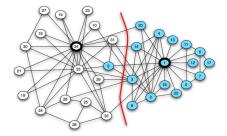


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- Second: The relative reputation, status, official position, exceptional attributes (intelligence, finances), etc...

# Power: Bridging and bonding capital of nodes

The early chapters of the text provided some insights about the importance of centrality and bonding capital and bridging capital with regard to the *flow of information* and *trust*.





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- The above is an informal definition of power, but the study of power in the context of imbalance is a well studied concept with precise definitions
- We will isolate power due to position in a network, and ignore the status aspects
- For motivation we begin with some illustrative network examples, we will follow this with a social experiment that will provide insight, and will in turn lead to precise definitions

## Some illustrative examples

- Assume \$1 is placed on each edge of the network
  - each node trying to reach an agreement (within a fixed amount of time) on how to split the dollar
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(a) 2-Node Path

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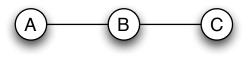


(a) 2-Node Path

#### Does either party have an advantage?

No; a  $\frac{1}{2} - \frac{1}{2}$  split is a reasonable predicted split that is observed in the experiments.

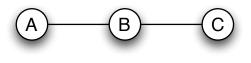
# A three node path



(b) 3-Node Path

#### What matching might occur and who each holds power?

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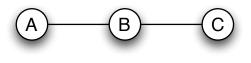
#### What matching might occur and who each holds power?

Clearly since we need a matching, either A and C will have to be left out. Intuitively then, node B holds much more power than A or C. The basic theory and experiments support this intuition. What fraction of the  $\$  would you expect B to obtain in negotiating between A and C?

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There is a difference between the basic theory and the social experiments. In the experiments , B gets a  $(\frac{5}{6})^{th}$  fraction of the \$. The basic theory would predict that B gets all almost all of the \$. Why the difference?

# A four node path



(c) 4-Node Path

What matching might occur and how might the money be split? Would *B* get more or less in this four node network than in the previous three node path?

# A four node path

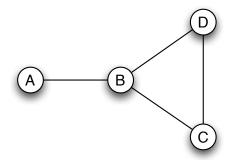


(c) 4-Node Path

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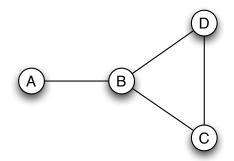
Here the experiments show that B gets a fraction of between  $\frac{7}{12}$ <sup>th</sup> and  $\frac{2}{3}$ <sup>rd</sup> of the \$, less than what we obtained in the three node network. Why?

# The stem graph in figure 12.3



What matching might occur and how might the money be split? Would *B* get more or less in this stem network than in the previous three and four node paths?

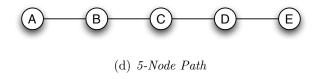
# The stem graph in figure 12.3



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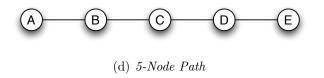
Experiments show that B in the stem graph makes slightly more money than B in the four node path (but less than in the three node path). Why?

## A five node path



Does C have any power (i.e. fraction of money obtained) compared to other nodes?

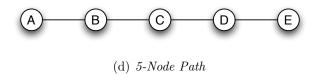
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Intuitively B and D have most of the power in the five node path network. The text states that in experiments, C has slightly more power than A or E.

# A five node path



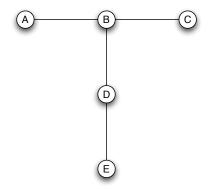
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Intuitively B and D have most of the power in the five node path network. The text states that in experiments, C has slightly more power than A or E.

Note that C is the most central node in terms of being on the most shortest paths. However, this has not translated into substantial bargaining power.

# Another graph to consider

The previous examples may help us reason about the following example from the text.



The following network exchange social experiment (and variants) is repeated a number of *rounds* so that some form of learning is taking place. There are many variants and the text presents one particular setting:

 Individuals (not knowing each other since we want to focus on the network aspects and not on the status, etc. of individuals) are placed at computer terminals and can interact with certain other individuals.

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- There is one more important condition on the experiment; namely in any given round, the outcome has to be a matching. i.e., you're only allowed to deal with one other person
  - This is called the 1-exchange rule

We would, of course, like to have results that are robust and do not differ that much in the exact "details".

- Results are reasonably robust with regard to how much network information is available
- Results are consistent across different countries and different cultures

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  - In certain networks, substantially different findings result if individuals can negotiate two or more exchanges in a round
  - In graph theory terms this is a *b*-matching; nodes can be adjacent to up to *b* edges in the matching
- Anonymity is important
  - How?

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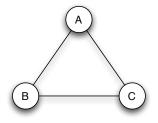
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  - How? Higher status individuals tend to inflate their "options", and those of lower status tend to underplay their options

#### **Demo time!**

I need two volunteers from the audience :)

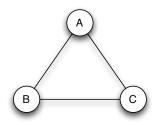


Question: For those who aren't volunteers – What solution will we converge to? As we go through the game, do you notice anything?

#### Do all experiments converge in a consistent manner?

In simple networks, each round tends to come to consistent outcomes within the specified time limits.

However, there are networks where this is not the case. Consider the following triangle graph:

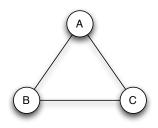


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Question: Notice anything?

Any two of the nodes can wind up excluding the other. Hence we would expect that the final outcome in any round will be determined by the two nodes who get to settle just before the time deadline.

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John Nash (the same Nash who showed that all finite games have mixed equilibria) introduced a specific stable outcome, the *Nash Bargaining Solution*.

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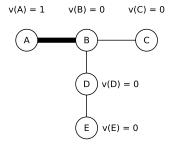
An outcome in a network exchange process on a graph G = (V, E) is a pair (M, v) where  $M \subseteq E$  is a matching and the value function  $v : V \rightarrow [0, 1]$  satisfies:

- For every edge  $e = (x, y) \in M$ ,  $v_x + v_y = 1$ .
- If a node x ∈ V is not part of the matching M (i.e. does not appear as a vertex in any edge (x, y) ∈ M), then v<sub>x</sub> = 0.

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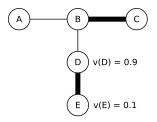


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$$v(A) = 0$$
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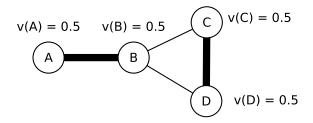
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Suppose  $v_{x'} + v_{y'} < 1$  for an edge  $(x', y') \notin M$ . Then the matching is unstable as there is a surplus of  $s = 1 - v_{x'} - v_{y'}$  that can be shared between x' and y' and there is no reason for them not to share this surplus and increase both their values.

#### Which stable outcome?

Stable solutions are necessary but there can be many stable solutions and some are more natural (in the sense of corresponding to real behaviour) than others.



### Mon. Mar 18th: Announcements & Corrections

- Assignment 2 is due Thu Mar 21
- Draft of critical review due Fri Mar 22 on PeerScholar
- You can access PeerScholar via the Quercus Assignment tab
  - Your individual peer reviews will be submitted via the same assignment! So make sure you're in the right group
- An oversight from last week: Naming the Nash Bargaining Solution!

- Suppose (x, y) ∈ M. What if x and y have other options other than to be in a given matching?
  - Suppose that x and y has the "outside options" of o<sub>x</sub> and o<sub>y</sub> respectively
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The Nash bargaining solution would be to keep (x, y) in the matching and equally divide up any surplus from the outside options:

• i.e., if 
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• Hence we get  $v_x + v_y = 1$ , with (x, y) in the matching.

#### Why extreme outcomes are not real outcomes

As stated earlier in this chapter, in the three node path example, the theory thus far would predict that *B* will obtain the entire \$. But we are told that in experiments, more typically *B* gets a fraction  $\frac{5}{6}$  and one other node gets a fraction  $\frac{1}{6}$ .

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In the following ultimatum game, we can perhaps better understand why participants tend to think beyond monetary rewards.

We again are considering how two individuals divide a \$. But now we have the following experiment:

• One person (say A) is given one \$ and is told to propose a division of it to person B.

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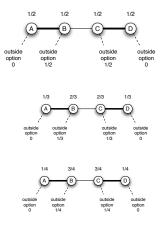
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Now in strictly monetary terms, person B should accept any offer (even a \$.01). But this is not what happens in experiments. In experiments, A tends to offer B about one third of the \$. Why? Aside: The Ultimatum Game is a little like the "I cut-you choose 2-person cake cutting algorithm" which ensures "fairness"

### Not all stable outcomes are "natural"

As we stated, there can be many stable outcomes for a given network. But some do not appear as natural as others and, in particular, stable outcomes can be "extreme solutions" that do not represent what we believe to be more realistic. Which of the following stable outcomes might be more expected "in practice"?



It turns out that the  $\frac{1}{3}$ ,  $\frac{2}{3}$  split between A and B and also between C and D is what happens more in experiments and can be considered "more natural" in the following way.

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It turns out that the  $\frac{1}{3}, \frac{2}{3}$  split is the Nash Bargaining solution which we argued seemed like a fair way to divide up surpluses.

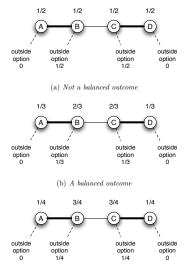
### What is a balanced outcome?

#### **Balanced outcomes**

An outcome (M, v) is balanced if for every edge in the matching M, the split of money  $\{v_x\}$  is the Nash bargaining solution for each node x, given the (best) outside options for each node.

**Fact:** For every exchange network with a stable outcome, there exists a balanced outcome.

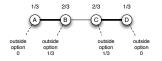
## Balanced and unbalanced outcomes for the four node path



(c) Not a balanced outcome

### Checking that the balanced outcome is the Nash Bargaining solution

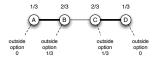
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Why is the best outside option for B (and similarly for C) equal to  $\frac{1}{3}$ ?

*B* has the option of offering  $\frac{2}{3}$  (or maybe  $\frac{2}{3} + \epsilon$  for some small  $\epsilon > 0$ ) to entice *C* to leave its current match with *D*. Therefore, *B* can receive at most  $\frac{1}{3} - \epsilon$ . Of course, *A* has no outside option so we we can calculate that surplus for the matched edge (A, B) is  $s = 1 - o_A - o_B = \frac{2}{3}$  and hence the Nash bargaining solution would be:

• 
$$v_A = o_A + \frac{s}{2} = 0 + \frac{1}{3} = \frac{1}{3}$$

• 
$$v_B = o_B + \frac{s}{2} = \frac{1}{3} + \frac{1}{3} = \frac{2}{3}$$

which is consistent with the balanced outcome.

Similarly, C and D follow the Nash Bargaining solution.

### Recap

With practice & review, you'll be able to:

- Define Mitochondrial Eve
  - ▶ Define the Wright-Fisher single-parent ancestry model (Ch 21.7)
  - Prove an estimation of time to convergence (Ch 21.8B)
- Define and reason about stability, with respect to bargaining in a Network Exchange Model
  - Explain power in the network exchange social experiment, from structural and non-structural factors (Ch 12.1-12.3)
  - Define stable outcomes (Ch 12.7), and determine if an outcome is stable
  - Define the Ultimatum Game (Ch 12.6), and explain how it compares to bargaining
  - Define balanced outcomes (Ch 12.5, 12.8), and determine if an outcome is balanced