Social and Information Networks

University of Toronto CSC303 Winter/Spring 2021

Week 12: April 9 (2021)

Today's agenda

- Having covered the course material, I've put much thought into what we should do on the last day of class
- We've all had a long term due to ongoing world events, so I've decided make a fun surprise for today
- Before we get into that, are there any questions about course material, the final, or anything else?
 - ► There are office hours today after class
 - I'm also happy to meet by appointment

Today's agenda

 Now without the least excuse or the slightest provocation, I'll fondly introduce (for your mental delectation), the first (pray it's not annual!) X-treme CSC303 review-and/or-additional-topics bingo! (and also Jeopardy)



 The rules are simple! Just download the list of bingo spots from Quercus, paste the list into https://jie-luo.github.io/BingoMaker/, and generate your card!

Exciting and dynamic demo here!

- As on Jeopardy, there is a grid of questions with various monetary values – a random person will be selected to chose & answer (or skip) a question, and the first person to answer correctly gets the points
- OR, for 50 points you can instead hear highlights from one of the papers that someone wrote their critical review about (if you correctly guess the number of your own paper then you get 500 points)

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- At the same time, if the answer is on your bingo card (or if you can find a connection with either the answer or the paper), then you can fill the space on your card
- At the end of the day, we'll have winners for Jeopardy, and Bingo
 - ▶ If you claim Bingo, you'll have to explain your spaces

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• Feel free to form teams via Zoom chat or external software!

EXCITING prizes!

 But what good is a game without prizes? So what can our lucky winners look forwards to?

EXCITING prizes!

 But what good is a game without prizes? So what can our lucky winners look forwards to?



Figure: Exciting prizes! Also, unnatural eyes, hands and arms!

 Nothing but the finest weird attribution free images, emailed directly to your home! (or not, I mean, do you really want images like that one on your hard drive? Ugh, it gives me the shivers)

That grid thing that always shows up in these (only now there's two of them, but not really because the second's a list)

The grid of questions:

Triangles,	Equilibria	Movin'	????
Triangles,		through a	
Triangles!		network	
100 (24)	100 (36)	100 (47)	100 (58)
200 (26)	200 (38)	200 (49)	200 (60)
300 (28)	300 (40)	300 (51)	300 (62)
400 (30)	400 (42)	400 (53)	400 (64)
500 (33)	500 (44)	500 (55)	500 (66)

For 50 points (or 500 if you get your own): Chose a paper, any paper!

Paper #1	Paper #2	Paper #3	Paper #4	Paper #5	Paper #6
(9)	(11)	(12)	(15)	(17)	(20)

PRIZES: B1 (69) B2 (70) B3 (71) J (72)

Paper # 1: The effects of balance on social fragmentation

- https://doi.org/10.1098/rsif.2020.0752
- Interesting paper, looks at finding a consistent explanation for increased connectivity between people (e.g., via the internet), and increased formation of isolated communities

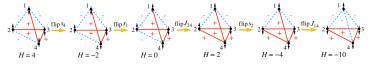


Figure 2. Example for how social stress, H, is gradually lowered in a sequence of changes of opinions and links between four individuals (assuming g = 1 for simplicity). Opinions of the nodes are given by \uparrow and \downarrow , positive links are red, negative are blue. In the course of this sequence, the number of balanced triangles changes from two to four. Note that in the second step social stress is temporarily increased. This is a consequence of the stochastic nature of the model, where also unfavourable events happen from time to time.

- Create a model with two opinions, edges can change sign or people can change opinion
- System changes over time to reduce total social tension
 - A combination of unstable triangles, and friends with opposite opinions
 - ► Therefore model creates stable triangles, and through social influence attempts to create homophily

Paper # 1: The effects of balance on social fragmentation

The authors ran their model on various graphs where all nodes have degree k.

They found that for a fixed amount of randomness T in the model, then there is a sharp threshold where on one side the graph is sharply fragmented, and on the other it is not.

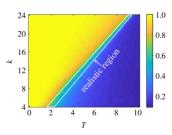


Figure 3. Phase diagram of the stochastic coevolutionary model with social balance. The balance level, f, is shown as a function of the average network degree, k, and social temperature, T. A phase separation line is visible. Below it, for low values of connectivity and high T, there exits a socially coherent phase (blue), above the line there is a phase of social fragmentation (yellow). Empirically reasonable values of f around 0.7 are indicated with a white box. Results were obtained for regular graphs ($\epsilon=0$), g=1, N=400 and are averaged over 500 realizations. Random initial conditions in links and opinions.

(8)

Paper # 2: Social network analysis: Characteristics of online social networks after a disaster

Table 6
Top 10 betweenness centrality and degree distribution.
Red-colored lines are all edges linked with top 10 betweenness centrality vertices excluding CBR.

and Marie	Users	Betweenness
	CBR	1,082,218
	Organization#1	166,500
	Organization#2	81,549
	Organization#3	49,982
	User#19	31,711
	User#20	29,937
	User#21	26,838
	Organization#4	22,838
	Organization#5	20,572
	User#22	13,791

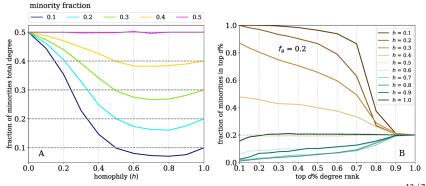
- https://doi.org/10.1016/j.ijinfomgt.2017.08.003
- Organizations have high betweenness centrality (total betweenness going through a node, where betweenness is as in Girvan-Newman algorithm), and function as gatekeeper nodes
- However, with the exception of the local government, the nodes of high out-degree tend to be people (8)

Paper # 3: Homophily influences ranking of minorities in social networks

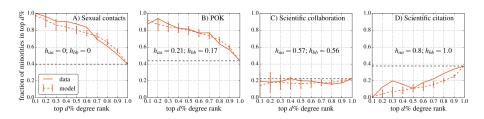
- https://doi.org/10.1038/s41598-018-29405-7
- Very nicely done paper, studies relationship between homophily (or it's opposite, heterophily) and the degree of nodes
- The authors modified the preferential-attachment model with a homophilily parameter, $h \in [0,1]$ (h < 0.5 is heterophily, and h > 0.5 is homophily)

Paper # 3: Homophily influences ranking of minorities in social networks

- Their simulations suggest that moderate homophily causes degree under-representation, both in total degree of nodes in the minority and and the proportion of the highest ranking nodes that are in the minority
- Strong homophily causes poor connectivity with the opposite group
- In contrast, heterophily benefits the minority group



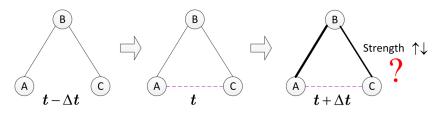
Paper # 3: Homophily influences ranking of minorities in social networks



 Model predictions compare favourably with real-life heterophilic networks (dating networks), and homophilic networks (scientific citations between research topics) (8)

Paper # 4: Will Triadic Closure Strengthen Ties in Social Networks?

- https://doi.org/10.1145/3154399
- We've previously talked about the interaction between tie-strength and triadic closure, but this paper seeks to go further by quantifying tie-strength



Paper # 4: Will Triadic Closure Strengthen Ties in Social Networks?

- The authors study two who-talks-to-whom networks
 - anonymized cellphone data (tie strength being number of calls)
 - Weibo data (tie strength being number of comments or retweets)
- Average tie-strength tends to increase over time
- ullet However, in pprox 80% of cases the original ties weaken

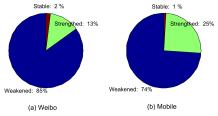


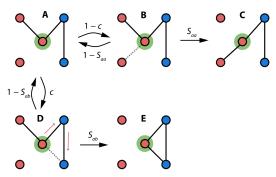
Fig. 3. Tie strength dynamics after triadic closure in social networks. In Weibo, ties in about 85% triads have been weakened, while in Mobile ties in 74% triads have been weakened.

- Demographics have a strong impact; all-male triads and celebrity triads are more strongly affected
- Men and celebrities are also more likely to maintain weaker triads (8)

Paper # 5: Cumulative effects of triadic closure and homophily in social networks

- This paper looks at the causes of homophily
 - ▶ Homophily due to selection (which they call *choice homophily*)
 - Homophily due to restrictions in who you have contact with (which they call induced homophily)
- The authors study the question of whether triadic closure can drive induced homophily

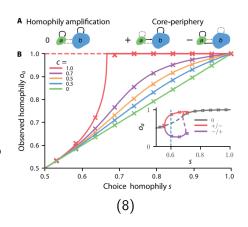
Paper # 5



- The model selects a focal node uniformly at random
 - ▶ with probability *c* it attempts to close a triangle
 - lacktriangle with probability 1-c it attempts to form a friendship with a random node
 - ► The proposal succeeds or fails with a probability determined by the endpoints' group
 - On success, the edge forms and a pre-existing edge connected to the focus is deleted

Paper # 5

- The model tends to one of two types of stable configuration (illustrated in A), homophily amplification or core-periphery
- We can see that as the probability of triadic closure, c, increases (line colour), the observed homophily is amplified from the choice homophily
- Observed homophily: proportion of neighbours of the same group
- Choice homophily: s is the probability of accepting a member of the same group, and 1 s is the probability of accepting an opposite member
- Results consistent with datasets

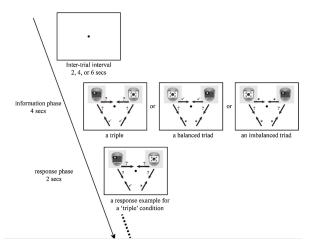


Paper # 6: Triadic balance in the brain: Seeking brain evidence for Heiders structural balance theory

- https://doi.org/10.1016/j.socnet.2020.05.003
- This paper attempted to find physical evidence of mental stress arising from unstable triangles
- Participants where placed in an fMRI machine, and played a 3-player variant of the prisoner's dilema for money
 - ► There are 3 players, they can defect or cooperate, the best outcome is when all players cooperate
 - Player A can chose to give player B 10 points, if B reciprocates then A gets a net gain of 10 points

Paper # 6:

 Participants where shown a triangle, and given the choice to change their decisions



Paper # 6:

 Participants have a weak, but significant, preference to change to balanced triads

 Table 1

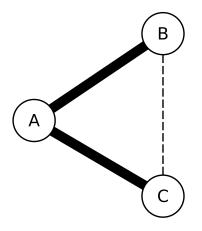
 Participants' decisions of triads given a specific initial triadic scenario.

Decision Stimuli	Balanced Triads	Unbalanced Triads
Balanced Triads ($N = 2083$)	1470 (71 %)	613 (29 %)
Unbalanced Triads ($N = 2042$)	721 (35 %)	1321 (65 %)

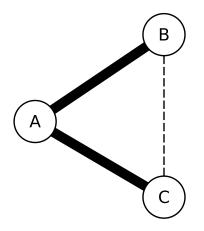
"Among the active brain areas ... two of them were significantly correlated with participants probability of settling on unbalanced triads: the left superior frontal gyrus ... and the right supramarginal gyrus ... intensity of activation of the two brain areas in unbalanced triads is even stronger for people who have stronger preferences for structural imbalance. The result suggests not only that cognitive dissonance occurs when people are confronted with structural imbalance, but also that it intensifies when people choose to settle on structural imbalance." (8)



• Does this graph satisfy strong triadic closure?



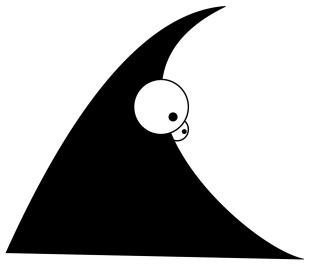
• Does this graph satisfy strong triadic closure?



• Yes, it only has the strong edges (A,B) and (A,C). STC implies that we should have the edge (B,C), and indeed it is present

• From the bounty of "triangle" search results, your prize is...

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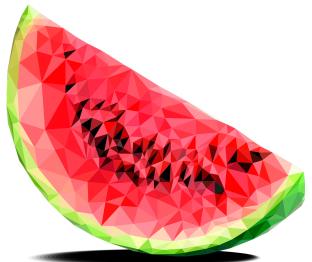
absolutely adorable! Can we keep it?;) (8)

• What are the three types of closure in a social-affiliation network?

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 - Triadic Closure: Two people become friends, because of a mutual friend
 - * May correspond to either selection, or influence
 - ② Focal Closure: Two people become friends, because they are both members of the same focus
 - ★ Corresponds to selection
 - Membership Closure: A person joins a focus, because they are friends with a member
 - ★ Corresponds to influence

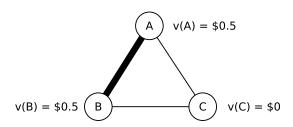
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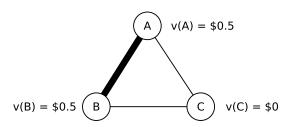


Polygonal watermelon! (Caution: sharp seeds) (8)

- In class we saw bargaining networks, in which adjacent nodes can enter a deal to split \$1 between themselves, and each node can only make a deal with one of its neighbours
- Is the following outcome stable?



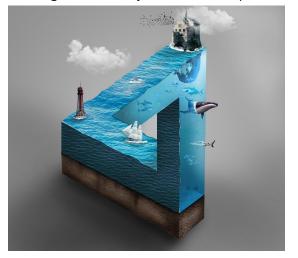
- In class we saw bargaining networks, in which adjacent nodes can enter a deal to split \$1 between themselves, and each node can only make a deal with one of its neighbours
- Is the following outcome stable?



- No, a matching is stable iff there is no pair of nodes which can enter a deal that strictly improves both of their payouts
 - ▶ Node C can offer node A \$0.60
- More generally, we saw in class that this bargaining network has no stable solution

• How have the triangles rewarded you? You're the proud owner of...

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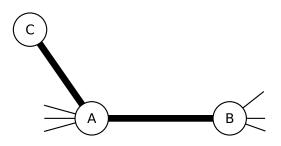


 A maritime-themed optimal illusion! May contain small parts unsuitable for children, spare whales sold separately (8)

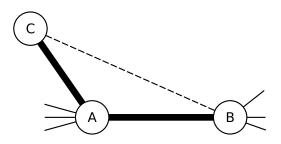
- Recall from class, that Strong Triadic Closure (STC) requires that for any strong edges (A, B) and (B, C), then there must exist the edge (A, C)
- What is the connection between STC and inter-community edges?

- Recall from class, that Strong Triadic Closure (STC) requires that for any strong edges (A, B) and (B, C), then there must exist the edge (A, C)
- What is the connection between STC and inter-community edges?
- At a high level, STC tells us that inter-community edges tend to be weak

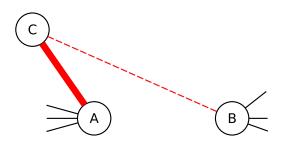
- More rigorously, we proved that for any node A with at least 2 strong edges, then any local bridge adjacent to A must be weak
 - ▶ Recall that an edge is a local bridge if, upon the removal of the edge, the shortest path between the endpoints is of length at least 3 (or no such path exists)



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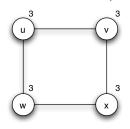
• Congratulations, you've just won a brand new...

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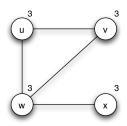


 Industrialized ancient Egypt? Remodeled Louvre? Well sue me, I think it looks pretty anyhow (8)

- In class we saw a revolt model in which nodes have a threshold t and will revolt if they know that t nodes (including themselves) will revolt
- Why is it impossible for a revolt to occur in b), yet a revolt will occur if we create the triangle in c)?



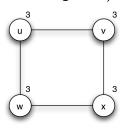
(b) An uprising will not occur



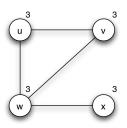
(c) An uprising can occur

[Modified from 19.14 of E&K]

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(b) An uprising will not occur



(c) An uprising can occur

[Modified from 19.14 of E&K]

• The triangle in (c) makes willingness to revolt shared knowledge, whereas in (b) no node knows that it's neighbours know that there enough willing people to meet their thresholds

• With a search term of "triangles", you are the lucky winner of...

• With a search term of "triangles", you are the lucky winner of...



• An antique, highly concerned (and concerning) cow-penguin! May you cherish it always. (8)

Equilibria



- Consider the Schelling model with a threshold of t = 3
- What does an equilibrium in this model look like?

- Consider the Schelling model with a threshold of t = 3
- What does an equilibrium in this model look like?
- Each node (either of class A or B) has at least 3 neighbours of the same class
- The Schelling model tends to converge to highly segregated solutions, this homophily is driven purely by (relatively mild) local choices

• And your prize is...

• And your prize is...



• "airship"! Methinks having sails perpendicular to propellers might not be the best arrangement, but at least the piece looks nice:) (8)

- Recall that for a game, the price of anarchy (POA) is defined as $\frac{cost(S)}{cost(OPT)}$ where cost is the social cost of a solution to the game, S is the worst Nash equilbrium for the game, and OPT is the optimal solution to the game
- In a congestion network, cost is the total ammount of time people spend commuting
- From class we know that the price of anarchy in a congestion network with linear cost functions is at most $\frac{4}{3}$
- Given that such a congestion network has an optimal social cost of 900 minutes, then what is the range of values that the real social cost would be in? (Hint: assume real traffic converges to the Nash Equilibrium)

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- In a congestion network, cost is the total ammount of time people spend commuting
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- Given that such a congestion network has an optimal social cost of 900 minutes, then what is the range of values that the real social cost would be in? (Hint: assume real traffic converges to the Nash Equilibrium)
- The cost would be at least 900, and at most $900 \times \frac{4}{3} = 1200$

• With a search term of "equilibria", you are the lucky winner of...

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• Disco Justice! (8)

- Recall the stable marriage problem from lecture
- What is the male-pessimal solution, and how can it be found?

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- What is the male-pessimal solution, and how can it be found?
- Recall that an outcome to the stable marriage problem is a matching between women W and men M
 - An outcome is stable iff it contains no blocking pair (m, w) where both prefer the other member of the blocking pair over their current match
- The male-pessimal solution is the outcome where each man is matched with their least-preferred woman, out of all woman that the man can be matched to in a stable matching
- The male-pessimal solution is stable, and can be found using the Gale-Shapely algorithm
 - Specifically, FPDA (Female Proposing, Deferred Acceptance)

• Your prize for getting this question right is...

• Your prize for getting this question right is...



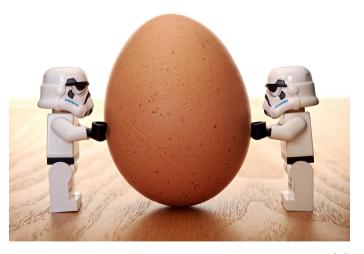
• The strangest looking picture of a peacock I've ever seen! (8)

- Recall that the hubs and authorities algorithm converges to an equilibrium, in which further updates no longer affect the value
- Given the vector of authority values at time t, \mathbf{a}_t , what is the matrix M such that we can compute our updated hub values as $\mathbf{h}_{t+1} = M\mathbf{a}_t$

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- Given the vector of authority values at time t, \mathbf{a}_t , what is the matrix M such that we can compute our updated hub values as $\mathbf{h}_{t+1} = M\mathbf{a}_t$
- *M* is the adjacency matrix of the graph
 - ▶ i.e., M_{ij} is 1 iff we have a hyperlink on page i pointing to page j

• With a search term of "equilibria", you can now possess...

• With a search term of "equilibria", you can now possess...



• a pair of Lego stormtroopers making a very large breakfast (8)

 What information can we extract from the eigenvalues and eigenvectors of the signed Laplacian matrix of a graph, G?

- What information can we extract from the eigenvalues and eigenvectors of the signed Laplacian matrix of a graph, G?
- If we let λ_1 to be the smallest eigenvalue, and x be a corresponding eigenvector, then:
 - If $\lambda_1=0$ then the graph is completable to a strongly balanced network
 - ★ If none of x's elements are zero (which will be the case if G is connected), then G is completable to two communities comprised of the positive and negative elements of x respectively
 - Otherwise, smaller values of λ_0 indicate that the graph is closer to being completable to a strongly balanced network

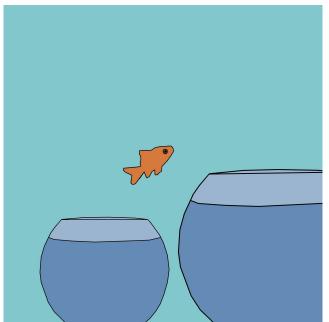
• For your marvelous memory of the matrix, your prize is...

• For your marvelous memory of the matrix, your prize is...



• the Lego stormtrooper's circus act! (8)

Movin' through a network



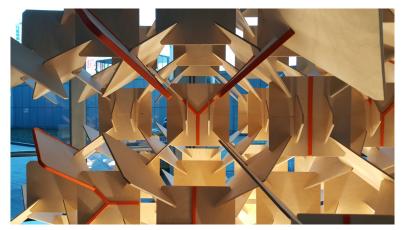
Movin' through a network for 100

• What is the basic reproductive number, R_0 ? (Hint: we covered this in the section on disease spread)

- What is the basic reproductive number, R_0 ? (Hint: we covered this in the section on disease spread)
- R₀ is the expected number of new cases created by one case of the disease
- In a branching process where each person is connected to k new individuals and the disease has a probability p of spreading, then $R_0 = pk$

• With that, you now have...

With that, you now have...



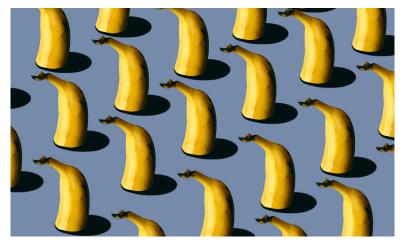
• Your own... uh... hmmm. Well, it's pretty in a headache inducing kind of way anyways. (8)

• Describe the unscaled PageRank update

- Describe the unscaled PageRank update
- Each node's rank is evenly divided among it's outgoing edges
 - Nodes with no outgoing edges keep all of their rank
- Recall that in the limit, this update is equivalent to looking at proportion of time spend in each node by a random walk on the graph

• With that, you have the dubious privilege of taking home with you...

• With that, you have the dubious privilege of taking home with you...



• The march of the bananas! (8)

- Recall the modified Watts-Strogatz model from class:
 - where every node u is on a grid
 - u is connected to its neighbours
 - u has a random long-distance edge to some node v
 - ▶ v is randomly chosen with probability proportional to $d(u, v)^{-2}$, where d is the grid distance
- Recall also decentralized search, where a node tries to pass a message to a target node t by passing it to neighbour with the lowest grid distance to t
- How does the length of the average path length grow with *n*?

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- Recall also decentralized search, where a node tries to pass a message to a target node t by passing it to neighbour with the lowest grid distance to t
- How does the length of the average path length grow with *n*?
- $O(\log n)$ from lecture, and we covered a proof of the 1D case in tutorial

• With a search term of "decentralized", you're the proud owner of...

• With a search term of "decentralized", you're the proud owner of...



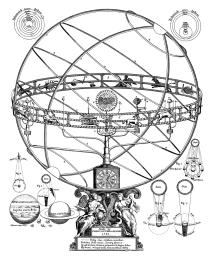
An evaporating physical Bitcoin – Bitcoin smoke, don't breath this!

- In class we saw that we can have efficient decentralized search in a modified Watts-Strogatz model by selecting v, the endpoint of the long-distance edges of a node u with probability proportional $d(u,v)^{-2}$
- In practice however, this doesn't work as human population is not uniformly dense
- What where the two alternatives to distance that we saw in class?

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- What where the two alternatives to distance that we saw in class?
- The alternatives we saw were:
 - Rank
 - ★ How many people are closer
 - ★ We have efficient decentralized search with $r_u(v)^{-1}$
 - Social Distance
 - ★ The size of the smallest shared social focus
 - ***** We have efficient decentralized search with $s(u, v)^{-1}$

• For getting this question right, your prize is...

• For getting this question right, your prize is...



• A diagram of an armillary sphere! (8)

- In class we covered various models for the diffusion of an idea or product through a network
- This led to the problem of finding an optimal initial set of adopters
- Given f(S) (the expected final number of adopters given the initial set of adopters, S), we can find a candidate solution by a greedy marginal improvement method:
 - ▶ We start with $S = \emptyset$, and repeatedly add the node that causes the greatest increase in f
- What were the properties of f that we saw are sufficient to ensure a constant approximation ratio of 1-1/e?

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- What were the properties of f that we saw are sufficient to ensure a constant approximation ratio of 1-1/e?
- f must be:
 - **1** Normalized: $f(\emptyset) = 0$
 - **2** Monotone: $\forall S, T \subseteq V : S \subseteq T \implies f(S) \leq f(T)$
 - Submodular (i.e. diminishing returns):

$$\forall S, T \subseteq V : \forall v \in V : S \subseteq T \implies f(S \cup \{v\}) - f(S) \ge f(T \cup \{v\}) - f(T)$$

• For your exemplary memory and dedication, you shall receive...

• For your exemplary memory and dedication, you shall receive...



• a picture of a portrait of the Right Honourable Pug-face himself! (β)/74

????



• In class, we demonstrated that processes driven by _____, often result in power laws

- In class, we demonstrated that processes driven by _____, often result in power laws
- Popularity
- More specifically:
 - We demonstrated that power laws have to emerge from correlated individual decisions
 - ► We saw examples where power laws emerge, that could be partially explained by popularity (e.g., sizes of cities, book sales, website in-links)
 - We analyzed the preferential attachment model, which is a model driven by popularity which produces power laws

• What "????" of your own have you won? Well...

• What "????" of your own have you won? Well...



 You've won your very own train hypnotist! (don't make eye contact and keep away from hypnotoads) (8)

 What is the clustering coefficient of a node? You can provide a mathematical definition, or a description

- What is the clustering coefficient of a node? You can provide a mathematical definition, or a description
- For a node A, the clustering coefficient is the probability that any randomly chosen pair of friends of A are themselves friends.
 Equivalently, it is the extent of triadic closure, measured as the proportion of possible triangles that are closed.
 For undirected edges E then:

$$cc(A) := \frac{|\{(u, v) \in E | u, v \in N(A), u \neq v\}|}{|\{\{u, v\} | u, v \in N(A) : u \neq v\}|} \in [0, 1]$$

 Recall that a low clustering coefficient can be suggestive of mental stress, as show in the Bearman & Moody study from week 1

• For getting this question right, you'll get to take home with you nothing but the finest result from searching "????"...

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 your very own bundle of raw asparagus! (now with strawberries and flowers). Enjoy! (8)

- Recall the Wright-Fisher model that we analyzed when studying at mitochondrial Eve
- Suppose that we have 10 single-cellular organisms in a petri-dish, each of which reproduces asexually (more specifically, by mitosis)
- A classmate proposes applying the Wright-Fisher model, and using it to argue that eventually all organisms in the dish will be the same species. What are some possible concerns with applying the Wright-Fisher model?

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- Some possible concerns are that the population will not be static (it will probably follow some sort of geometric progression), and some species may be better suited to the environment than others
- To apply the model, we would have to address these points

• Congratulations! You've won...

Congratulations! You've won...



• What appears to be a stained glass window of a rooster surrounded by some *rather* grouchy teddy-bears! (8)

• Briefly describe the Rozenstein algorithm for minimizing weak edges, while maintaining strong connectivity of provided communities

- Briefly describe the Rozenstein algorithm for minimizing weak edges, while maintaining strong connectivity of provided communities
- The algorithm considers the edge causing the most STC violations, and makes it weak if it can do so without disconnecting one of the provided communities – once considered an edge will not be considered again

• For getting this question right, your prize is...

• For getting this question right, your prize is...



• a house of cards made of iphones! (8)

????? for 500

- Recall that we consider an arbitrary graph to be balanced iff it can be completed into a complete balanced graph
 - ► A complete graph is balanced if all triangles are balanced (either 3 or 2 positive edges)
- What is the compact characterization of whether an arbitrary graph is balanced? (Hint: it's closely related to the characterization of whether or not a graph is bipartite)

???? for 500

- Recall that we consider an arbitrary graph to be balanced iff it can be completed into a complete balanced graph
 - ► A complete graph is balanced if all triangles are balanced (either 3 or 2 positive edges)
- What is the compact characterization of whether an arbitrary graph is balanced? (Hint: it's closely related to the characterization of whether or not a graph is bipartite)
- An arbitrary graph is strongly balanced if it contains no cycles with an odd number of negative edges

???? for 500

- Brief proof:
 - ▶ By Harary Balance theorem, a balanced complete graph has no negative edges, or we can be divided into two camps (X and Y) which are internally positive, and dislike eachother
 - ▶ If the graph is completable then the completed graph contain no cycles with an odd number of negative edges. Therefore, the original cannot



Figure 5.10: If a signed graph contains a cycle with an odd number of negative edges, then it is not balanced. Indeed, if we pick one of the nodes and try to place it in X, then following the set of friend/enemy relations around the cycle will produce a conflict by the time we get to the starting node.

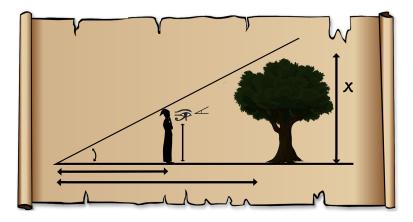
If the graph doesn't contain a cycle with an odd number of negative edges then we can find the communities connected by positive edges (they can't contain a negative edges as that would make an odd cycle). We can collapse these communities into supernodes connected by negative edges. This graph is bipartite as there are no cycles with an odd number of negative edges. We now do a 2-colouring using BFS.

????? for 500

• Are you ready for your prize? You are the proud winner of...

????? for 500

• Are you ready for your prize? You are the proud winner of...



• The eye of Horus compelling a man in a night cap to do trigonometry! (8)

Bingo Prize: 1rst Place

• Our lucky, first place Bingo winner gets Gold Medal...

Bingo Prize: 1rst Place

• Our lucky, first place Bingo winner gets Gold Medal...



• flour. I for one, blame the search results;) (8)

Bingo Prize: 2nd Place

• Our Bingo runner-up gets...

Bingo Prize: 2nd Place

• Our Bingo runner-up gets...



Concerned fennel! (8)

Bingo Prize: 3rd Place

• Our third place Bingo winner gets...

Bingo Prize: 3rd Place

• Our third place Bingo winner gets...



• A small toy bulldozer pushing caviar off of (buttered?) toast! (8)

Jeopardy Prize: 3rd Place

• Our third place Jeopardy winner gets...

Jeopardy Prize: 3rd Place

• Our third place Jeopardy winner gets...



Parent-child bonding on a frosty morning! (8)

Jeopardy Prize: 2nd Place

• Our Jeopardy runner-up gets...

Jeopardy Prize: 2nd Place

• Our Jeopardy runner-up gets...



• Scenic hiking! (8)

Jeopardy Prize: 1rst Place

• Our lucky, first place Jeopardy winner gets...

Jeopardy Prize: 1rst Place

• Our lucky, first place Jeopardy winner gets...



• Lego Star Wars bohemian rhapsody! What you've never known you've always wanted! (8)