

424 Prove the following definitions implement simple data-stack theory (Subsection 7.0.2).

$$stack = [nil], [stack; X]$$

$$push = \langle s: stack \cdot \langle x: X \cdot [s; x] \rangle \rangle$$

$$pop = \langle s: stack \cdot s 0 \rangle$$

$$top = \langle s: stack \cdot s 1 \rangle$$

After trying the question, scroll down to the solution.

§ Consider the implementation to be four axioms, named by their left sides. Now I prove each of the axioms of simple data-stack theory. First,  $stack \neq null$  by contradiction.

$$\begin{aligned}
 & stack = null && \text{conjoin } stack \text{ axiom} \\
 = & stack = null \wedge stack = [nil], [stack; X] && \text{context, then specialize} \\
 \Rightarrow & null = [nil], [null; X] && \text{both ; and [ ] distribute over ,} \\
 = & null = [nil], null && null \text{ is identity for ,} \\
 = & null = [nil] && \text{transparency} \\
 \Rightarrow & \phi null = \phi [nil] && \text{size axioms; note that [nil] is an element} \\
 & && \text{because all 0 of its items are elements} \\
 = & 0 = 1 && \text{arithmetic axiom} \\
 = & \perp
 \end{aligned}$$

Let  $s: stack$  and  $x: X$ . Then

$$\begin{aligned}
 & push\ s\ x : stack && \text{use } push \text{ and } stack \text{ axioms} \\
 = & \langle s: stack \cdot \langle x: X \cdot [s; x] \rangle \rangle s\ x : [nil], [stack; X] && \text{apply} \\
 = & [s; x]: [nil], [stack; X] && \text{generalization} \\
 \Leftarrow & [s; x]: [stack; X] \\
 = & \top \\
 & pop\ (push\ s\ x) = s && \text{use } pop \text{ and } push \text{ axioms} \\
 = & \langle s: stack \cdot s\ 0 \rangle \langle s: stack \cdot \langle x: X \cdot [s; x] \rangle \rangle s\ x = s && \text{apply} \\
 = & \langle s: stack \cdot s\ 0 \rangle [s; x] = s && \text{apply} \\
 = & [s; x]\ 0 = s && \text{index} \\
 = & \top \\
 & top\ (push\ s\ x) = x && \text{use } top \text{ and } push \text{ axioms} \\
 = & \langle s: stack \cdot s\ 1 \rangle \langle s: stack \cdot \langle x: X \cdot [s; x] \rangle \rangle s\ x = x && \text{apply} \\
 = & \langle s: stack \cdot s\ 1 \rangle [s; x] = x && \text{apply} \\
 = & [s; x]\ 1 = x && \text{index} \\
 = & \top
 \end{aligned}$$

The last step, indexing, requires  $x$  to be an item, so this implementation requires  $X$  to be a bunch of items.