

411 Let x be an integer variable.

(a) Using the recursive time measure, add time and then find the strongest implementable specification S that you can find for which

$S \Leftarrow$ **if** $x=0$ **then** *ok*
else if $x>0$ **then** $x:=x-1$. S
else $x' \geq 0$. S **fi fi**

Assume that $x' \geq 0$ takes no time.

(b) What do we get from recursive construction starting with $t' \geq t$?

After trying the question, scroll down to the solution.

- (a) Using the recursive time measure, add time and then find the strongest implementable specification S that you can find for which

$$S \Leftarrow \begin{array}{l} \mathbf{if } x=0 \mathbf{ then } ok \\ \mathbf{else if } x>0 \mathbf{ then } x:=x-1. S \\ \mathbf{else } x' \geq 0. S \mathbf{ fi fi} \end{array}$$

Assume that $x' \geq 0$ takes no time.

§ Adding time,

$$S \Leftarrow \begin{array}{l} \mathbf{if } x=0 \mathbf{ then } ok \\ \mathbf{else if } x>0 \mathbf{ then } x:=x-1. t:=t+1. S \\ \mathbf{else } x' \geq 0 \wedge t'=t. t:=t+1. S \mathbf{ fi fi} \end{array}$$

the strongest implementable solution for S is

$$x'=0 \wedge \mathbf{if } x \geq 0 \mathbf{ then } t' = t+x \mathbf{ else } t' \geq t+1 \mathbf{ fi}$$

If we replace $x' \geq 0 \wedge t'=t$ by $x:=c$ where c is an arbitrary natural number, then we can prove

$$x'=0 \wedge \mathbf{if } x \geq 0 \mathbf{ then } t' = t+x \mathbf{ else } t' = t+1+c \mathbf{ fi}$$

- (b) What do we get from recursive construction starting with $t' \geq t$?

$$\begin{array}{l} \S S_n = \begin{array}{l} 0 \leq x < n \wedge x'=0 \wedge t'=t+x \\ \vee \neg 0 \leq x < n \wedge t' \geq t+n \\ \vee x < 0 \wedge x'=0 \wedge t+1 \leq t' < t+n \end{array} \\ S_\infty = \begin{array}{l} 0 \leq x \wedge x'=0 \wedge t'=t+x \\ \vee x < 0 \wedge t' = \infty \\ \vee x < 0 \wedge x'=0 \wedge t+1 \leq t' < \infty \end{array} \end{array}$$

S_∞ is a solution to the given implication, but not as strong as the solution shown in part (a). It is interesting to note that if the given implication were an equation, then S_∞ would not be a solution (fixed-point), but the solution of part (a) would still be a solution. $\Downarrow n \cdot S_n$ is the same as S_∞ .