

277 Given three sorted lists having at least one item common to all three, write a program to find the smallest item occurring in all three lists.

After trying the question, scroll down to the solution.

§ Let the three lists be A , B , and C and let a , b , and c be natural variables to act as indexes of the three lists. The specification is R , defined as

$$\begin{aligned} R = & \quad A(0,..a') \wedge B(0,..b') \wedge C(0,..c') = null \\ & \wedge A a' = B b' = C c' \\ & \wedge t' = t + a' + b' + c' \end{aligned}$$

which says there isn't a common item before indexes a' , b' , and c' , and there is one at those indexes, and the sum of those indexes is the execution time. We also need the specification Q , defined as

$$\begin{aligned} Q = & \quad A(a,..a') \wedge B(b,..b') \wedge C(c,..c') = null \\ & \wedge A a' = B b' = C c' \\ & \wedge t' = t + a' - a + b' - b + c' - c \end{aligned}$$

which says the same thing, but starting at indexes a , b , and c . The refinements are

$$\begin{aligned} R & \Leftarrow a := 0. b := 0. c := 0. Q \\ Q & \Leftarrow \text{if } A a < B b \text{ then } a := a + 1. t := t + 1. Q \\ & \quad \text{else if } B b < C c \text{ then } b := b + 1. t := t + 1. Q \\ & \quad \text{else if } C c < A a \text{ then } c := c + 1. t := t + 1. Q \\ & \quad \text{else ok fi fi fi} \end{aligned}$$

The proof of the first refinement is immediate after using the Substitution Law 3 times. The proof of the last refinement breaks into 12 pieces (3 conjuncts \times 4 cases). Let's start with the first conjunct of Q and the first case of the **if**.

$$\begin{aligned} & A a < B b \wedge (a := a + 1. t := t + 1. A(a,..a') \wedge B(b,..b') \wedge C(c,..c') = null) \\ & \hspace{15em} \text{substitution law twice} \\ = & \quad A a < B b \wedge A(a+1,..a') \wedge B(b,..b') \wedge C(c,..c') = null \\ & \hspace{10em} \text{Because } B \text{ is sorted, } B b \text{ is the minimum item of } B(b,..b'). \\ & \hspace{10em} \text{And since } A a < B b, \text{ therefore } A a \text{ is unequal to any item in } B(b,..b'). \\ \Rightarrow & \quad A(a,..a') \wedge B(b,..b') \wedge C(c,..c') = null \end{aligned}$$

Now we prove the middle conjunct of Q with the same first case of the **if**.

$$\begin{aligned} & A a < B b \wedge (a := a + 1. t := t + 1. A a' = B b' = C c') \hspace{5em} \text{substitution law twice} \\ = & \quad A a < B b \wedge A a' = B b' = C c' \hspace{10em} \text{specialize} \\ \Rightarrow & \quad A a' = B b' = C c' \end{aligned}$$

Now we prove the last conjunct of Q with the same first case of the **if**.

$$\begin{aligned} & A a < B b \wedge (a := a + 1. t := t + 1. t' = t + a' - a + b' - b + c' - c) \hspace{5em} \text{substitution law twice} \\ = & \quad A a < B b \wedge t' = t + 1 + a' - a - 1 + b' - b + c' - c \hspace{5em} \text{simplify and specialize} \\ \Rightarrow & \quad t' = t + a' - a + b' - b + c' - c \end{aligned}$$

The proofs of the next two cases are exactly the same. That leaves the last case.

$$\begin{aligned} & A a \geq B b \wedge B b \geq C c \wedge C c \geq A a \wedge ok \hspace{5em} \text{antisymmetry, and expand } ok \\ & A a = B b = C c \wedge a' = a \wedge b' = b \wedge c' = c \wedge t' = t \\ \Rightarrow & \quad Q \end{aligned}$$