

222 (natural square root) The natural square root of a natural number n is the natural number s satisfying

$$s^2 \leq n < (s+1)^2$$

- (a) Write a program to find the natural square root of a given natural number n in $\log n$ time.
- (b) Write a program to find the natural square root of a given natural number n in $\log n$ time using only addition, subtraction, doubling, halving, and comparisons (no multiplication or division).

After trying the question, scroll down to the solution.

(a) Write a program to find the natural square root of a given natural number n in $\log n$ time.

§ The solution is a binary search for s in the interval $0..n+1$.

$$s'^2 \leq n < (s'+1)^2 \iff s:=0. r:=n+1. s^2 \leq n < r^2 \implies s'^2 \leq n < (s'+1)^2$$

$$s^2 \leq n < r^2 \implies s'^2 \leq n < (s'+1)^2 \iff$$

if $s+1 = r$ **then** *ok*

else $m := \text{div}(s+r) 2$.

if $m^2 \leq n$ **then** $s := m$ **else** $r := m$ **fi**.

$s^2 \leq n < r^2 \implies s'^2 \leq n < (s'+1)^2$ **fi**

Execution time is identical to binary search.

(b) Write a program to find the natural square root of a given natural number n in $\log n$ time using only addition, subtraction, doubling, halving, and comparisons (no multiplication or division).

§ The problem with the program in part (a) is the occurrence of m^2 . We'll have to work in the squared space to avoid the need for squaring. Our three variables will be:

s is the left end of the remaining search space times the size of the remaining search space. The left end is initially 0, so s is initially 0. The size of the search space is finally 1, so s' is the answer.

b is the square of the size of the remaining search space. To cut the search space in half, we must divide b by 4. To repeatedly divide b by 4, it must be a power of 4.

n is the number whose square root is sought minus the square of the left end of the search space. Initially the left end is 0, so n starts as the number whose square root is sought. n measures, in squared space, the distance from s to the answer.

Define

$$P = n' = n < b' : b \times 4^{nat} \wedge s' = s \wedge \text{if } b \leq 4 \times n \text{ then } b' \leq 4 \times n \text{ else } b' = b \text{ fi}$$

$$\wedge t' = t + (\log(b'/b))/2$$

$$Q = n < b + 2 \times s \wedge (n \leq 2 \times s \implies s = 0) \wedge (\exists m, l. b = 4^m \wedge s = l \times 2^m)$$

$$\implies b' = 1 \wedge n' = n - s'^2 \leq 2 \times s' \wedge t' = t + (\log b)/2$$

Then

$$s'^2 \leq n < (s'+1)^2 \wedge t' \leq t + \log(n+1) + 2 \iff s:=0. b:=1. P. Q$$

$P \iff$ **if** $b > n$ **then** *ok* **else** $b := 2 \times b. b := 2 \times b. t := t+1. P$ **fi**

$Q \iff$ **if** $b = 1$ **then** *ok*

else $b := b/2. b := b/2. c := s+b. s := s/2.$

if $c > n$ **then** *ok* **else** $(n := n-c. s := s+b)$ **fi**.

$t := t+1. Q$ **fi**