

124 Let n be a natural state variable. Is the following specification implementable?

(a) $n := n - 1$

(b) $n > 0 \Rightarrow (n := n - 1)$

(c) **if** $n > 0$ **then** $n := n - 1$ **else** *ok* **fi**

After trying the question, scroll down to the solution.

$$\begin{aligned}
& \text{(a)} && n := n-1 \\
& \S && \forall n: \text{nat} \exists n': \text{nat} \cdot n := n-1 && \text{expand assignment} \\
& = && \forall n: \text{nat} \exists n': \text{nat} \cdot n' = n-1 && \text{specialization} \\
& \Rightarrow && \exists n': \text{nat} \cdot n' = 0-1 && \text{arithmetic} \\
& = && \exists n': \text{nat} \cdot n' = -1 \\
& = && \perp
\end{aligned}$$

So no, $n := n-1$ is not implementable. From the line

$$\begin{aligned}
& \forall n: \text{nat} \exists n': \text{nat} \cdot n' = n-1 && \text{we can use an identity law} \\
= & \forall n: \text{nat} \exists n': \text{nat} \cdot n' = n-1 \wedge \top && \text{but now we cannot use the one-point law to get} \\
& \forall n: \text{nat} \exists n': \text{nat} \cdot \top && \text{because the one-point law requires } n-1: \text{nat}
\end{aligned}$$

$$\begin{aligned}
& \text{(b)} && n > 0 \Rightarrow (n := n-1) \\
& \S && \forall n: \text{nat} \exists n': \text{nat} \cdot n > 0 \Rightarrow (n := n-1) && \text{expand assignment} \\
& = && \forall n: \text{nat} \exists n': \text{nat} \cdot n > 0 \Rightarrow n' = n-1 && \text{distributive and identity} \\
& = && \forall n: \text{nat} \cdot n > 0 \Rightarrow \exists n': \text{nat} \cdot n' = n-1 \wedge \top && \text{In the context } n > 0, n-1: \text{nat}. \\
& && && \text{So we can use one-point.} \\
& = && \forall n: \text{nat} \cdot n > 0 \Rightarrow \top && \text{base and identity} \\
& = && \top
\end{aligned}$$

So yes, $n > 0 \Rightarrow (n := n-1)$ is implementable.

$$\begin{aligned}
& \text{(c)} && \mathbf{\text{if } n > 0 \text{ then } n := n-1 \text{ else } ok \text{ fi}} \\
& \S && \forall n: \text{nat} \exists n': \text{nat} \cdot \mathbf{\text{if } n > 0 \text{ then } n := n-1 \text{ else } ok \text{ fi}} && \text{expand assignment and } ok \\
& = && \forall n: \text{nat} \exists n': \text{nat} \cdot \mathbf{\text{if } n > 0 \text{ then } n' = n-1 \text{ else } n' = n \text{ fi}} && \text{case analysis} \\
& = && \forall n: \text{nat} \exists n': \text{nat} \cdot (n > 0 \wedge n' = n-1) \vee (n = 0 \wedge n' = n) && \text{splitting} \\
& = && \forall n: \text{nat} \cdot (\exists n': \text{nat} \cdot n > 0 \wedge n' = n-1) \vee (\exists n': \text{nat} \cdot n = 0 \wedge n' = n) && \text{In the context } n > 0, \\
& && && n-1: \text{nat}. \text{ And in the context } n = 0, n: \text{nat}. \text{ So we can apply one-point twice.} \\
& = && \forall n: \text{nat} \cdot n > 0 \vee n = 0 \\
& = && \top
\end{aligned}$$

So yes, $\mathbf{\text{if } n > 0 \text{ then } n := n-1 \text{ else } ok \text{ fi}}$ is implementable.