

# ERGO on a Page (v1.5)

## The ERGO File

An *ERGO* file for an application typically has three parts: a basic action theory, the definitions of some *ERGO* procedures (using `define` as usual), and an optional robotic interface. These three parts are described further below, and may be loaded from other files by using `include`. The *ERGO* file usually ends with a main function that calls `ergo-do` or one of the planning functions:

```
(define (main) (ergo-do [#:mode mode] pgm))
```

The *pgm* here is an expression that evaluates to an *ERGO* program. If the execution *mode* is specified, it should be one of 'offline', 'first' (the default), or 'online'. (The robotic interface part is needed only when the mode is 'online'.)

## Using Scheme

Within an *ERGO* file, *ERGO* can be intermingled with Scheme variables and function definitions that appear in the usual way. In what follows, we use “*fexpr*” to mean any Scheme expression where the fluents of the basic action theory (see `define-fluents` below) may appear as global variables.

## Running ERGO

Once *ERGO* has been properly installed, it is possible to call the main function in an *ERGO* file called `my-ergo-app.scm` as follows:

```
> racket -l ergo -f my-ergo-app.scm -m
```

## Basic Action Theory

A basic action theory has definitions for the fluents and actions.

### Fluents

The fluents of a basic action theory are defined using one or more expressions in the file of the form

```
(define-fluents
  fluent ini
  ...
  fluent ini)
```

where each *fluent* is a symbol and each *ini* is a Scheme form that provides the value of the fluent in the initial state. This has the effect of defining the fluents as global variables that can then be used later in *fexprs* for actions and programs.

Any valid Scheme datum can be used as the value of a fluent, including lists, vectors, hash-tables, and functions.

### Actions

Each action is defined by an expression of the following form:

```
(define-action action
  fluent fexpr
  ...
  fluent fexpr)
```

The fluents listed are those that are considered to be changed by the action. The value of the *fluent* after the action will be the value of the corresponding *fexpr* before the action. (All changes are considered to be done in parallel.)

The *action* in the definition can be a symbol or a list of symbols (*name var ... var*) for an action with parameters..

In addition, the *fluent* can be the special symbol `#:prereq` or `#:sensing`, in which case, the corresponding *fexpr* defines the prerequisite or sensing result of the action.

The `define-action` expression has the effect of defining the action itself as a global variable whose value is the action symbol (or a list of the action symbol and its arguments).

## ERGO Programs

Each of the following expressions evaluates to an *ERGO* program:

```
:nil
  The program that always succeeds.

:fail
  The program that always fails.

(:test fexpr)
  Succeed or fail according to whether or not the current
  value of fexpr is true.

(:act action)
  Fail if the action has a prerequisite that evaluates to false,
  but succeed otherwise, and move to a new state where the
  fluents are updated as per its define-action (see above).

(:begin pgm ... pgm)
  Sequentially perform all of the programs.

(:choose pgm ... pgm)
  Nondeterministically perform one of the programs.

(:if fexpr pgm1 pgm2)
  Behave like pgm1 if the current value of fexpr is true, but
  like pgm2 otherwise.

(:when fexpr pgm ... pgm)
  Behave like (:if fexpr (:begin pgm...pgm) :nil)

(:unless fexpr pgm ... pgm)
  Behave like (:when (not fexpr) pgm...pgm).

(:until fexpr pgm ... pgm)
  Perform (:begin pgm...pgm) repeatedly until the value
  of fexpr becomes true.

(:while fexpr pgm ... pgm)
  Behave like (:until (not fexpr) pgm...pgm).

(:star pgm ... pgm)
  Perform (:begin pgm...pgm) repeatedly for some
  nondeterministically chosen number of times.

(:for-all var list pgm ... pgm)
  Perform (:begin pgm...pgm) repeatedly for all values of
  the variable var taken from the list list.
```

```
(:for-some var list pgm ... pgm)
  Perform (:begin pgm...pgm) for some value of var from
  the list, chosen nondeterministically.

(:conc pgm ... pgm)
  Concurrently perform all of the programs,
  nondeterministically interleaving any single steps.

(:atomic pgm ... pgm)
  Perform (:begin pgm...pgm) but with no interleaving
  from concurrent programs.

(:monitor pgm1 pgm2 ... pgmn)
  Perform pgm1 before every step of pgm2, which is
  performed before every step of pgm3, and so on.

(>> fexpr ... fexpr)
  Succeed after evaluating the expressions (for effect).

(<< fexpr ... fexpr)
  Like >> except that evaluation only happens on failure.

(:let ((var fexpr) ... (var fexpr)) pgm ... pgm)
  Perform (:begin pgm...pgm) in an environment where
  each variable var has the value of fexpr.

(:wait)
  Succeed after the next exogenous action happens.

(:search pgm ... pgm)
  Perform (:begin pgm...pgm) online, but using lookahead
  for any nondeterminism to guard against failure.
```

Note that obtaining the value of these expressions is not the same as executing them. Execution is what is done by `ergo-do` alone.

## Robotic Interface

A robotic interface is defined by a set of expressions of the form:

```
(define-interface 'out body)
  The body should evaluate to a function of one argument
  (like displayln) that will send an endogenous action to
  an outside target, blocking until the target is ready.

(define-interface 'in body)
  The body should evaluate to a function of no arguments
  (like read) that will return the next exogenous action
  received from an outside source, blocking when none.
```

The bodies can perform whatever initialization is needed for the functions they return to work properly (such as opening files, or making TCP connections). More than one in and out interface can be defined. The functions `write-endogenous` and `read-exogenous` can be used as default bodies.