

1. Consider the following statement:

If m and n are odd integers, then mn is an odd integer.

(a) Express the statement using logical notation.

$\forall m \in \mathbb{Z}, \forall n \in \mathbb{Z}, [(m \text{ is odd} \wedge n \text{ is odd}) \Rightarrow (mn \text{ is odd})]$

Alternate: $\forall m \in \mathbb{Z}, \forall n \in \mathbb{Z}, [(\exists k \in \mathbb{Z}, m = 2k + 1) \wedge (\exists k \in \mathbb{Z}, n = 2k + 1) \Rightarrow (\exists k \in \mathbb{Z}, mn = 2k + 1)]$

(b) This statement can be proven using a direct proof. Write a detailed proof *structure* for the statement. **Don't write a complete proof** — for now, focus on the proof structure only and leave out *all* of the “middle” of the argument.

Assume $m, n \in \mathbb{Z}$. # m and n are arbitrary elements of \mathbb{Z}

Assume $(m \text{ is odd} \wedge n \text{ is odd})$. # the antecedent

⋮

Then mn is odd. # definition of odd

Then $(m \text{ is odd} \wedge n \text{ is odd}) \Rightarrow (mn \text{ is odd})$. # introduce \Rightarrow

Then $\forall m \in \mathbb{Z}, \forall n \in \mathbb{Z}, [(m \text{ is odd} \wedge n \text{ is odd}) \Rightarrow (mn \text{ is odd})]$ # introduce \forall

(c) Now, complete the proof of the statement.

Assume $m, n \in \mathbb{Z}$. # m and n are arbitrary elements of \mathbb{Z}

Assume $(m \text{ is odd} \wedge n \text{ is odd})$. # the antecedent

Then $(\exists k \in \mathbb{Z}, m = 2k + 1)$ and $(\exists k \in \mathbb{Z}, n = 2k + 1)$. # definition of odd

Let $i \in \mathbb{Z}$ be such that $m = 2i + 1$. # label the quotient $m/2$ by i

Let $j \in \mathbb{Z}$ be such that $n = 2j + 1$. # label the quotient $n/2$ by j

Then $mn = (2i + 1)(2j + 1)$ # substitution

$= 4ij + 2i + 2j + 1$ # algebraic manipulation

$= 2(2ij + i + j) + 1$

Let $p = 2ij + i + j$.

Then $p \in \mathbb{Z}$. # since \mathbb{Z} closed under $+$, \times .

Then $mn = 2p + 1$. # substitution

Then $\exists k \in \mathbb{Z}, mn = 2k + 1$.

Then mn is odd. # definition of odd

Then $(m \text{ is odd} \wedge n \text{ is odd}) \Rightarrow (mn \text{ is odd})$. # introduce \Rightarrow

Then $\forall m \in \mathbb{Z}, \forall n \in \mathbb{Z}, [(m \text{ is odd} \wedge n \text{ is odd}) \Rightarrow (mn \text{ is odd})]$ # introduce \forall

2. Consider the following statement:

If m and n are integers with mn odd, then m and n are odd.

(a) Express the statement using logical notation.

$\forall m \in \mathbb{Z}, \forall n \in \mathbb{Z}, [(mn \text{ is odd}) \Rightarrow (m \text{ is odd} \wedge n \text{ is odd})]$

Alternate: $\forall m \in \mathbb{Z}, \forall n \in \mathbb{Z}, [(\exists k \in \mathbb{Z}, mn = 2k + 1) \Rightarrow (\exists k \in \mathbb{Z}, m = 2k + 1) \wedge (\exists k \in \mathbb{Z}, n = 2k + 1)]$

(b) This statement can be proven using an **indirect** proof. Write a detailed proof *structure* for the statement. **Don't write a complete proof** — for now, focus on the proof structure only and leave out *all* of the “middle” of the argument.

Assume $m, n \in \mathbb{Z}$. # m and n are arbitrary elements of \mathbb{Z}

Assume $(m \text{ is even} \vee n \text{ is even})$. # the negation of the consequent

[Since at least one of m or n is even, let us label one of the even numbers as m and make no assumption about n . The number n could be odd or even. (This argument is often labelled “Without loss of generality, assume m is even.” or “WLOG assume m is even.”)]

WLOG, assume m is even.

⋮

Then mn is even.

Then $(m \text{ is even}) \Rightarrow (mn \text{ is even})$. # introduce \Rightarrow

Then $(m \text{ is even} \vee n \text{ is even}) \Rightarrow (mn \text{ is even})$. # introduce disjunction in antecedent

Then $(mn \text{ is odd}) \Rightarrow (m \text{ is odd} \wedge n \text{ is odd})$. # apply contrapositive

Then $\forall m \in \mathbb{Z}, \forall n \in \mathbb{Z}, [(mn \text{ is odd}) \Rightarrow (m \text{ is odd} \wedge n \text{ is odd})]$. # introduce \forall

(c) Now, complete the proof of the statement.

Assume $m, n \in \mathbb{Z}$. # m and n are arbitrary elements of \mathbb{Z}

Assume $(m \text{ is even} \vee n \text{ is even})$. # the negation of the consequent

WLOG, assume m is even.

Then $\exists k \in \mathbb{Z}, m = 2k$. # definition of even

Let $i \in \mathbb{Z}$ be such that $m = 2i$. # label the quotient $m/2$ by i

Then $mn = 2in$ # substitution

$= 2(in)$ # associativity

Let $p = in$.

Then $p \in \mathbb{Z}$. # since \mathbb{Z} closed under \times .

Then $mn = 2p$. # substitution

Then $\exists k \in \mathbb{Z}, mn = 2k$.

Then mn is even. # definition of even

Then $(m \text{ is even}) \Rightarrow (mn \text{ is even})$. # introduce \Rightarrow

Then $(m \text{ is even} \vee n \text{ is even}) \Rightarrow (mn \text{ is even})$. # introduce disjunction in antecedent

Then $(mn \text{ is odd}) \Rightarrow (m \text{ is odd} \wedge n \text{ is odd})$. # apply contrapositive

Then $\forall m \in \mathbb{Z}, \forall n \in \mathbb{Z}, [(mn \text{ is odd}) \Rightarrow (m \text{ is odd} \wedge n \text{ is odd})]$. # introduce \forall