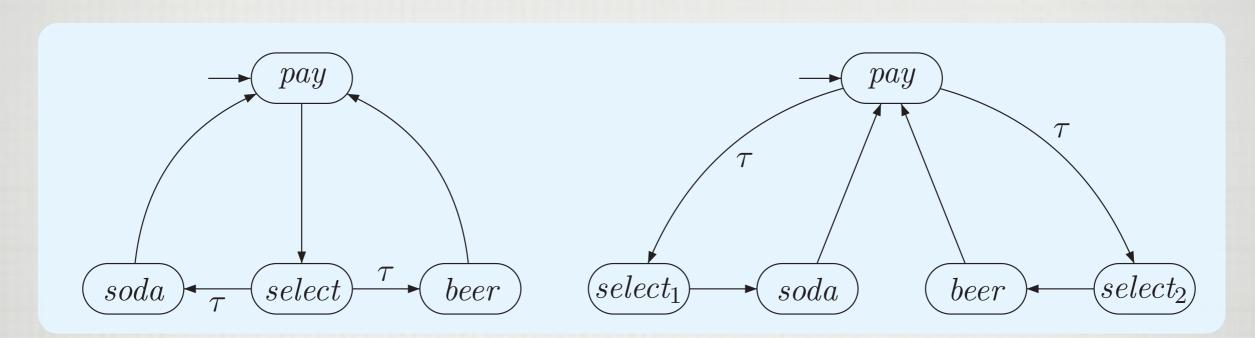
# CSC410

AZADEH FARZAN

FALL 2020

# Computational Tree Logic (CTL)

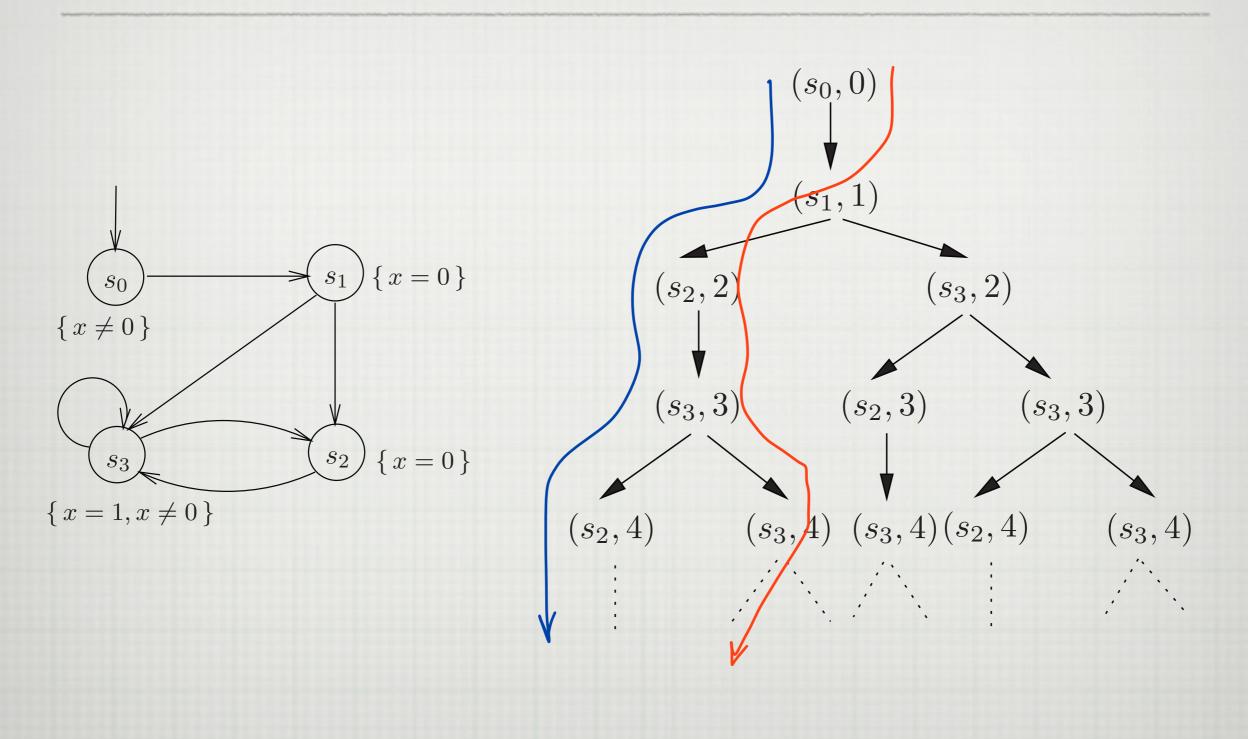
#### Limitations of LTL



They are trace equivalent.

These two transition systems satisfy the same set of LTL formulas. But they function in different ways.

## Computational Tree Logic (CTL)



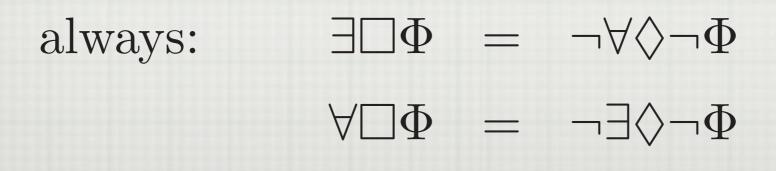
State Formula
 
$$\Phi ::= true$$
 $a$ 
 $\Phi_1 \land \Phi_2$ 
 $\neg \Phi$ 
 $\exists \varphi$ 
 $\forall \varphi$ 

 Path Formula
  $\varphi ::= \bigcirc \Phi$ 
 $\Phi_1 \cup \Phi_2$ 

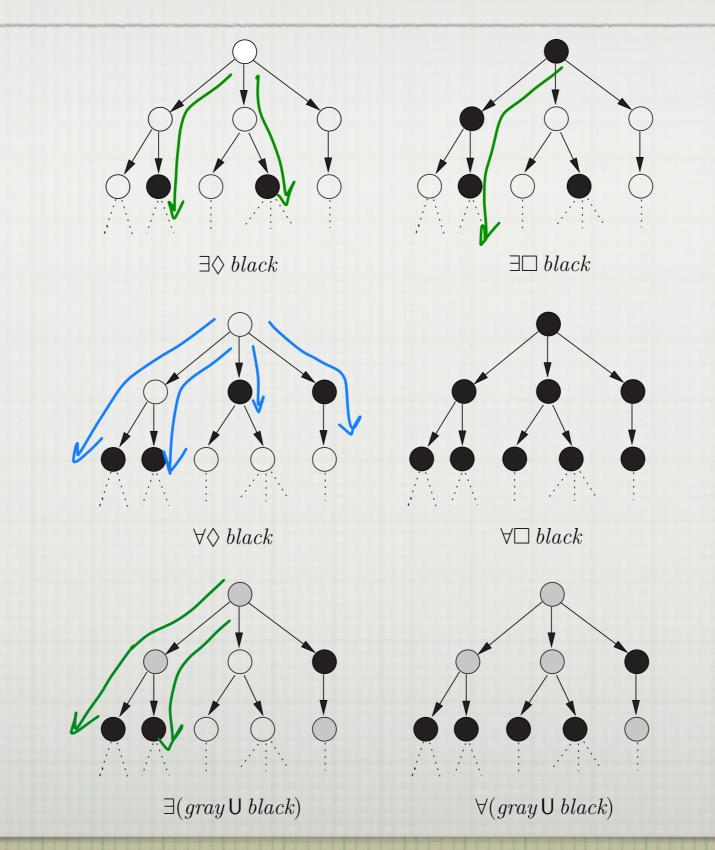
Examples:  $\exists \bigcirc (x = 1)$   $\forall \bigcirc (x = 1)$  needs a path quantification But not:  $\exists (x = 1 \land \forall \bigcirc (x \ge 3))$   $\exists \bigcirc (true \bigcup (x = 1))$ needs a femporal operator!

#### **Eventually and Always**

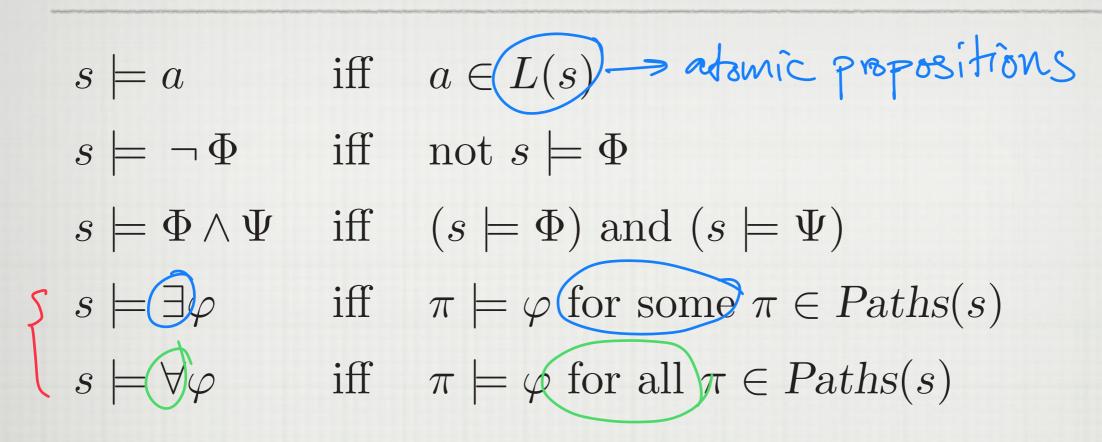
eventually:  $\exists \Diamond \Phi = \exists (\operatorname{true} \bigcup \Phi)$  $\forall \Diamond \Phi = \forall (\operatorname{true} \bigcup \Phi)$ 



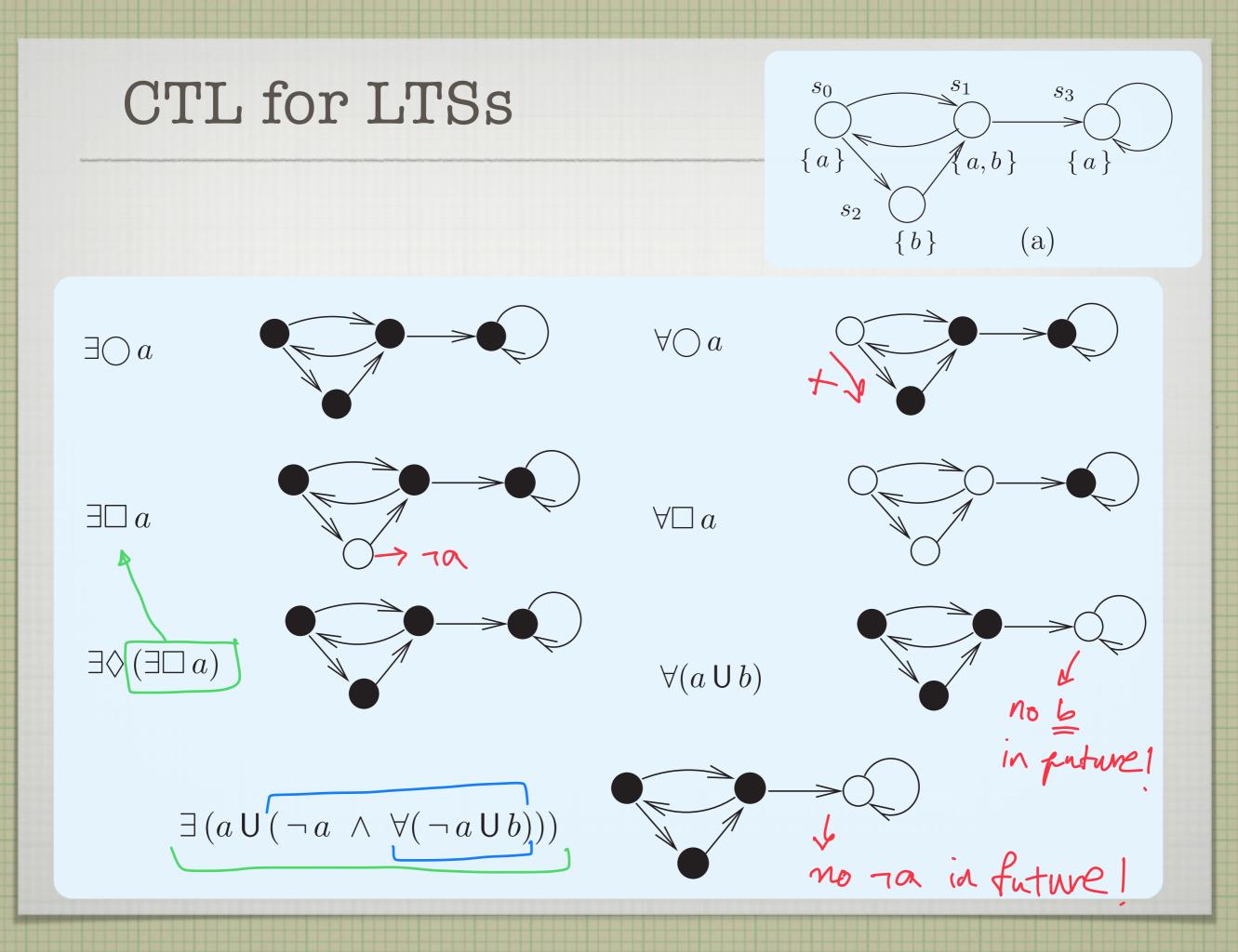
## Meaning of CTL: Examples



#### **CTL** Semantics



Same as LTL  $\begin{cases} \pi \models \bigcirc \Phi & \text{iff} \quad \pi[1] \models \Phi \\ \pi \models \Phi \cup \Psi & \text{iff} \quad \exists j \ge 0. \ (\pi[j] \models \Psi \land (\forall 0 \le k < j. \pi[k] \models \Phi)) \end{cases}$ 



# End of CTL