

# Designing Samplers is Easy: The Boon of Testers

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(Relevant Publications: AAAI-19, FMCAD-21, CP-22)

**Input** A CNF Formula  $F$  and tolerance parameter  $\varepsilon$

**Output**  $\sigma \in \text{Sol}(F)$  such that

$$\frac{1}{(1 + \varepsilon)|\text{Sol}(F)|} \leq \Pr[\mathcal{A}(F) = \sigma] \leq \frac{1 + \varepsilon}{|\text{Sol}(F)|}$$

**Motivation:** Fundamental problem in CS (theory) and applications in hardware and software testing (practice)

Snapshot from early 2010's

**Scalability** WES04, NRJK+06, KK07

**Guarantees** JVV86, BGP00, YAPA04

- Core Idea: Use 3-wise independence (random XORs) to partition the solution space
- Makes  $\mathcal{O}(\log n)$  calls to SAT oracle
- Theoretical guarantees

$$\frac{1}{(1 + \varepsilon)|\text{Sol}(F)|} \leq \Pr[\mathcal{A}(\mathcal{F}) = y] \leq \frac{1 + \varepsilon}{|\text{Sol}(F)|}$$

- Scalability: CryptoMiniSat (A specialized solver for CNF+XOR)

## How do you test a sampler is uniform?

**Input:** A reference sampler  $\mathcal{U}$ , a test sampler  $\mathcal{A}$ , and a formula  $F$

**Approach:** Run both samplers and plot their distributions

- Eyeball the distributions
- Run statistical tests (KL divergence, chi-square)

**Caveat** Requires number of samples  $\gg$  number of solutions

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What if you try to draw conclusions based on fewer samples?

**DLBS18: Efficient Sampling of SAT Solutions for Testing**

“We can see that SearchTreeSampler and UniGen2 are more uniform, but **QuickSampler** is still close to uniform on most benchmarks. However, this result should be taken with care, since the uniformity test is not very reliable on benchmarks where QuickSampler completed a small number of epochs or when the number of produced samples is too low.”

# In search of principled approach

**Input:** A reference sampler  $\mathcal{U}$ , a test sampler  $\mathcal{A}$ , and a formula  $F$

**Problem:** Return **Yes** if the distribution of  $\mathcal{U}(F)$  (known to be uniform) and  $\mathcal{A}(F)$  are close, else return **No**

**Approach II:** Just keep sampling and stop the first time you see a collision

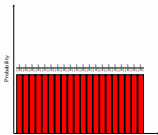


Figure:  $\mathcal{U}$ : Reference Distribution

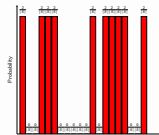


Figure:  $\mathcal{A}$ : far from uniform

No collisions until you have generated at least  $\sqrt{|Sol(F)|}$  solutions!

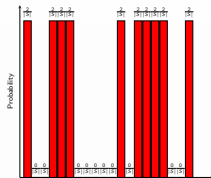
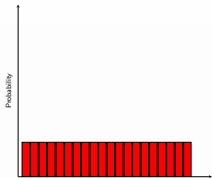
**BFRSW98**  $\implies$  The above technique is *optimal* (i.e., if we are only allowed to look at samples)

## Definition (Conditional Sampling)

Given a distribution  $\mathcal{D}$  on  $S$ ; allow one to specify a set  $T \subseteq S$  and draw samples from  $\mathcal{A}$  conditioned on  $T$

$$\Pr[\sigma \text{ is generated}] = \begin{cases} 0 & \text{if } \sigma \notin T \\ \frac{\mathcal{D}(\sigma)}{\sum_{\sigma \in T} \mathcal{D}(\sigma)} & \text{otherwise} \end{cases}$$

Conditional sampling is at least as powerful as drawing normal samples but is it more powerful?



- Draw  $\sigma_1$  uniformly at random from the domain and draw  $\sigma_2$  according to the distribution  $\mathcal{A}$ . Let  $\mathcal{T} = \{\sigma_1, \sigma_2\}$ .
- In the case of the “far” distribution, with constant probability,  $\sigma_1$  will have “low” probability and  $\sigma_2$  will have “high” probability.
- We will be able to distinguish the far distribution from the uniform distribution using constant number of samples from  $\mathcal{A}|\mathcal{T}$ .
- The constant depend on the fairness parameter.

The above algorithm works for all cases



- Input formula:  $F$  over variables  $X$
- **Challenge:** Conditional Sampling over  $T = \{\sigma_1, \sigma_2\}$ .
- Construct  $G = F \wedge (X = \sigma_1 \vee X = \sigma_2)$
- Most of the samplers will just enumerate all the solutions when the number of solutions is very small
- Need way to construct formulas whose solution space is large but every solution can be mapped to either  $\sigma_1$  or  $\sigma_2$ .

Input: A Boolean formula  $\varphi$ , two assignments  $\sigma_1$  and  $\sigma_2$ , and desired number of solutions  $\tau$

Output: Formula  $\hat{\varphi}$

- $\tau = |\text{Sol}(\hat{\varphi})|$
- $z \in \text{Sol}(\hat{\varphi}) \implies z_{\downarrow \text{Supp}(\varphi)} \in \{\sigma_1, \sigma_2\}$
- $|\{z \in \text{Sol}(\hat{\varphi}) \mid z_{\downarrow \text{Supp}(\varphi)} = \sigma_1\}| = |\{z \in \text{Sol}(\hat{\varphi}) \mid z_{\downarrow \text{Supp}(\varphi)} = \sigma_2\}|$
- $\varphi$  and  $\hat{\varphi}$  has "*similar*" structure

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### Definition

The **non-adversarial sampler assumption** states that the distribution of the projection of samples obtained from  $\mathcal{A}(\hat{\varphi})$  to variables of  $\varphi$  is same as the conditional distribution of  $\mathcal{A}(\varphi)$  restricted to either  $\sigma_1$  or  $\sigma_2$

- If  $\mathcal{A}$  is a uniform sampler for all the input formulas, it satisfies non-adversarial sampler assumption
- If  $\mathcal{A}$  is not a uniform sampler for all the input formulas, it may not necessarily satisfy non-adversarial sampler assumption

**Input:** A sampler under test  $\mathcal{A}$ , a reference uniform sampler  $\mathcal{U}$ , a tolerance parameter  $\varepsilon > 0$ , an intolerance parameter  $\eta > \varepsilon$ , a guarantee parameter  $\delta$  and a CNF formula  $\varphi$

**Output:** ACCEPT or REJECT with the following guarantees:

- if the generator  $\mathcal{A}$  is an  $\varepsilon$ -additive almost-uniform generator then Barbarik ACCEPTS with probability at least  $(1 - \delta)$ .
- if  $\mathcal{A}(\varphi, \cdot)$  is  $\eta$ -far from a uniform generator and if non-adversarial sampler assumption holds then Barbarik REJECTS with probability at least  $1 - \delta$ .
- Barbarik needs at most  $K = \tilde{O}\left(\frac{1}{(\eta - \varepsilon)^4}\right)$  samples.

- Samplers without guarantees (Uniform-like Samplers):
  - STS (Ermon, Gomes, Sabharwal, Selman,2012)
  - QuickSampler (Dutra, Laeuffer, Bachrach, Sen, 2018)
- Sampler with guarantees:
  - UniGen3

	QuickSampler	STS	UniGen3
ACCEPT <sub>s</sub>	0	14	50
REJECT <sub>s</sub>	50	36	0

To ACCEPT, we needed  $10^6$  samples but we could reject with just 250 samples

How can we use the availability of Barbarik to design a good sampler? Is it even possible ?

### Wishlist

- Sampler should be at least as fast as STS and QuickSampler.
- Sampler should pass the Barbarik test.
- Sampler should perform well on real world applications.

- Exploits the flexibility of CryptoMiniSat.
- Pick polarities and branch on variables at random.
  - To explore the search space as evenly as possible.
  - To have samples over all the solution space.
- Turn off all pre and preprocessing.
  - Processing techniques: bounded variable elimination, local search, or symmetry breaking, and many more.
  - Can change solution space of instances.
- Restart at static intervals.
  - Helps to generate samples which are very hard to find.

```
./cryptominisat5 -maxsol $1 -nobansol -restart fixed -maple 0 --verb 0 -scc 1 -n 1  
-presimp 0 -polar rnd -freq 0.9999 -fixedconfl $2 -random $3 -dumpresult $4 [CNFFILE]
```

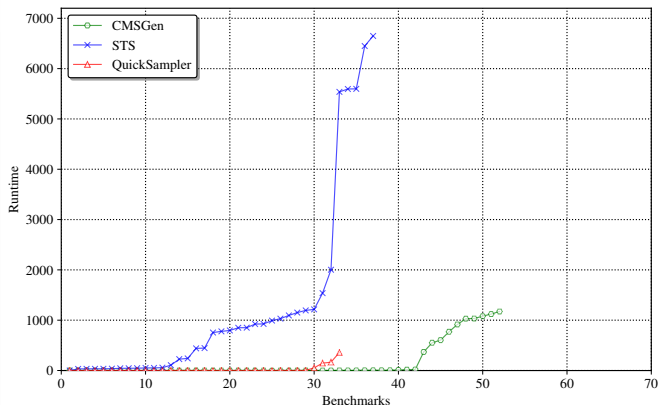
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- Parameters of CMSGen are decided iteratively with the help of Barbarik
- Lack of theoretical analysis.



# CMSGen vs. Other State-of-the-Art Samplers (I)



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QuickSampler	STS	CMSGen
33	37	52

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  - **CMSTGen**
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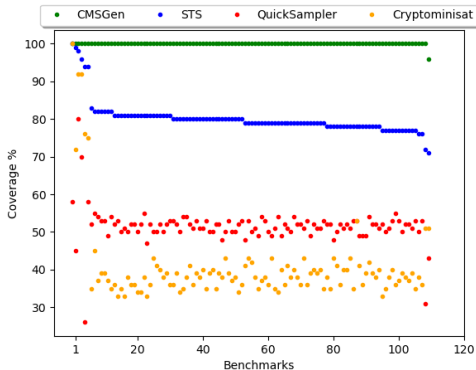
- Sampler should be at least as fast as STS and QuickSampler. ✓
- Sampler should pass the Barbarik test. ✓
- Sampler should perform well on real world applications.

- A powerful paradigm for testing configurable system.
- Challenge: To generate test suites that maximizes  $t$ -wise coverage.

$$\text{t-wise coverage} = \frac{\# \text{ of t-sized combinations in test suite}}{\text{all possible valid t-sized combinations}}$$

- To generate the test suites use constraint samplers.
- Uniform sampling to have high  $t$ -wise coverage (Plazar, Acher, Perrouin et al., 2019).
- Experimental Evaluations:
  - Generate 1000 samples (test cases).
  - 110 Benchmarks, Timeout: 3600 seconds
  - 2-wise coverage  $t = 2$ .

Higher is better



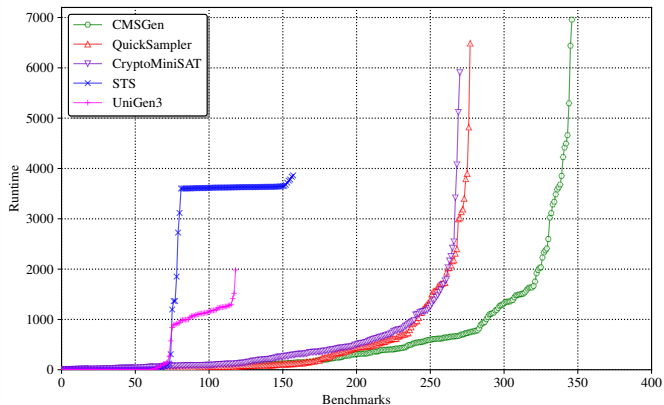
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	QuickSampler	STS	CMSGen
Avg. Coverage	51.5%	80.15%	~ 100%

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Remark: UniGen3 could sample for only 6 benchmarks

State of the art approach (Manthan): Sampling + Machine Learning + Counter-example guided repair



**Summary** Design of a practically efficient sampler via test-driven development that works well in real-world applications

**Practice** A Virtuous cycle: Improve Barbarik so that it can reject CMSGen and then improve CMSGen

- Trade-off between runtime performance and quality
- Frequent restarts degrade solution quality

**Theory** Explain why CMSGen works well

- Perhaps CDCL with randomization is all you need in practice?
- Perhaps, you don't really need uniformity in most cases. What do we really need?

**Theory and Practice** And a testing methodology independent of non-adversarial assumption