Midterm Test

Friday February 25, 2000

**Duration:** 50 minutes

**Aids allowed:** None

**Family Name:** ____________________________  **Given names:** ____________________________

**Student #:** ____________________________  **Tutor:** ____________________________

- There are 7 pages, including this one. The test is out of 35 marks and the value of each question is provided; please use this information to manage your time effectively.

- For questions that involve writing code, comments are not necessary. If you need to call a standard method but can’t remember the correct order of arguments, just indicate the meaning of each argument.

<table>
<thead>
<tr>
<th>Part</th>
<th>Score</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>____ / 5</td>
</tr>
<tr>
<td>B</td>
<td>____ / 6</td>
</tr>
<tr>
<td>C</td>
<td>____ / 4</td>
</tr>
<tr>
<td>D</td>
<td>____ / 6</td>
</tr>
<tr>
<td>E</td>
<td>____ / 5</td>
</tr>
<tr>
<td>F</td>
<td>____ / 4</td>
</tr>
</tbody>
</table>

**Total** ____ / 30
Part A [5 marks in total; no marks will be given without the “Tree / Explanation”]
Here is a table of characters and their corresponding frequency in a file we would like to compress.

<table>
<thead>
<tr>
<th>character</th>
<th>frequency</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>60%</td>
</tr>
<tr>
<td>B</td>
<td>30%</td>
</tr>
<tr>
<td>C</td>
<td>5%</td>
</tr>
<tr>
<td>D</td>
<td>2%</td>
</tr>
<tr>
<td>E</td>
<td>3%</td>
</tr>
</tbody>
</table>

Below are five different possible encodings. For each one, circle whether or not it could be a Huffman encoding. If you circle “yes”, draw the corresponding tree. If you circle “no”, explain why it is not possible for the encoding to be a Huffman encoding.

### Example 1

<table>
<thead>
<tr>
<th>character</th>
<th>code</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>1</td>
</tr>
<tr>
<td>B</td>
<td>11</td>
</tr>
<tr>
<td>C</td>
<td>110</td>
</tr>
<tr>
<td>D</td>
<td>1110</td>
</tr>
<tr>
<td>E</td>
<td>1111</td>
</tr>
</tbody>
</table>

**Yes**   **No**

Tree / Explanation: The code for A is also a prefix for several other codes. So there would be no way to decode a message like 1110 — one can’t tell whether it represents AC or just D.

### Example 2

<table>
<thead>
<tr>
<th>character</th>
<th>code</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>0</td>
</tr>
<tr>
<td>B</td>
<td>11</td>
</tr>
<tr>
<td>C</td>
<td>101</td>
</tr>
<tr>
<td>D</td>
<td>1001</td>
</tr>
<tr>
<td>E</td>
<td>1000</td>
</tr>
</tbody>
</table>

**Yes**   **No**

Tree / Explanation: (I have chosen to put 0 branches to the left and 1 branches to the right. It could just as well be drawn the other way.)
### Character Code Table

<table>
<thead>
<tr>
<th>Character</th>
<th>Code</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>110</td>
</tr>
<tr>
<td>B</td>
<td>10</td>
</tr>
<tr>
<td>C</td>
<td>1110</td>
</tr>
<tr>
<td>D</td>
<td>0</td>
</tr>
<tr>
<td>E</td>
<td>1111</td>
</tr>
</tbody>
</table>

**Yes**  **No**

**Tree / Explanation:** Huffman coding would make sure that A’s code is no longer than D’s, since is is more frequent than D.

### Character Code Table

<table>
<thead>
<tr>
<th>Character</th>
<th>Code</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>00</td>
</tr>
<tr>
<td>B</td>
<td>01</td>
</tr>
<tr>
<td>C</td>
<td>10</td>
</tr>
<tr>
<td>D</td>
<td>111</td>
</tr>
<tr>
<td>E</td>
<td>110</td>
</tr>
</tbody>
</table>

**Yes**  **No**

**Tree / Explanation:** Because A’s frequency is more than 50%, Huffman coding is guaranteed to give it a code of length 1. (Think about it: the total of the other characters’ frequencies is less than A’s frequency, so A is guaranteed to be hooked into the tree at the last step.)

### Character Code Table

<table>
<thead>
<tr>
<th>Character</th>
<th>Code</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>1</td>
</tr>
<tr>
<td>B</td>
<td>01</td>
</tr>
<tr>
<td>C</td>
<td>001</td>
</tr>
<tr>
<td>D</td>
<td>0001</td>
</tr>
<tr>
<td>E</td>
<td>0000</td>
</tr>
</tbody>
</table>

**Yes**  **No**

**Tree / Explanation:**

![Huffman Tree Diagram]

**Tree / Explanation:**
Part B [6 marks in total]

Consider 3 files of records with very different layouts.

File A Contains records about students, with the following fixed length fields, in this order:
- StudentName 30 characters, left justified, padded with blanks
- ID 9 digits (represented as text)
- GPA 4 characters xx.x where x is a digit 0..9
There are no delimiters between fields or between records.

File B Contains records about companies, with the following fields, in this order:
- CompanyName variable length, character data
- Address variable length, character data
- PostalCode 6 characters
- ContactName variable length, character data
There is a delimiter after each field and no additional delimiter between records.

File C Contains records about images, where each record has the following fields, in this order:
- ImageNameLength 5 bits indicates the length of the ImageName field
- ImageName variable character data
- Xoffset 4 bits (value 0..15)
- Yoffset 4 bits (value 0..15)
- Scale 4 byte floating point value, in binary format
There are no delimiters between fields or between records.

1. In File B, the actual character for the delimiter is not stated. The designer could choose to use a readable or non-readable character for this. State one advantage for each alternative.
   - Advantage of non-readable delimiter:
     That would leave all the readable characters as valid characters for within a field’s contents.
   - Advantage of readable delimiter:
     That would make the file easier for a human to look at.

2. Which field in any of the files might contribute to internal fragmentation?
   - StudentName

3. In File B, which delimiters could be removed without making other changes to the file format? Circle the one correct answer.
   - (a) those between CompanyName and Address
   - (b) those between Address and PostalCode
   - (c) **those between PostalCode and ContactName**
   - (d) those between ContactName and CompanyName
   - (e) all the delimiters are needed

4. State the limit on the number of characters in the ImageName field in File C. 31 or 32

5. Which file could have records referenced by RRN? Circle one.
   - A
   - B
   - C
Part C [4 marks in total]
Suppose an operating system keeps track of the blocks of a file using a multi-level index with 3 levels. Each block of the index holds $2^{10}$ bytes and a pointer to a file block takes 4 bytes. What is the size of the biggest file possible in this file system?

Answer: $2^{24}$ blocks

Rough work:

- There are $2^{10}$ bytes per block and $2^2 = 4$ bytes per pointer, so a block can hold $2^{10}/2^2 = 2^8$ pointers.
- So there are $2^8$ blocks (nodes) in the 2nd level of the index, and $2^8 \times 2^8 = 2^{16}$ in the 3rd.
- Those $2^{16}$ leaf blocks can each hold $2^8$ pointers, so in total $2^{16} \times 2^8 = 2^{24}$ file blocks can be referenced.

Part D [6 marks in total]

Suppose we have a file of integers, in binary format. Write a fragment of C++ code that goes to the integer that is 6 integers away from the beginning of the file, reads that integer, and then prints it to the standard output. The code to open the file is already written.

```cpp
fstream file;
file.open("fileOfInts", ios::in);
```

Answer:

```cpp
file.seekg( 6 * sizeof(int), ios::beg );
int num;
file.read( &num, sizeof(int) );
cout << num;
```
Part E [5 marks in total]

Below is a small program, and the statement used to compile it. It won't compile. Make any changes necessary below so that it will compile and run.

File Node.h

```cpp
template <class X>
class Node {
public:
  Node<X> () { data=0; next=0; };
  // Set my data field to "value".
  void set(X * value);
  // Return my data field.
  X * get();
private:
  X * data;
  Node<X> * next;
};
```

File Node.cpp

```cpp
#include "Node.h"

template<class X> void Node<X>::set (X * value){
  data = value;
};

template<class X> X * Node<X>::get (){ return data;
};
```

File driver.cpp

```cpp
#include <fstream.h>
#include "Client.h"
#include "Node.h" <-- #include "Node.cpp"

class Client {
public:
  int accountNumber;
  float balance;
};

int main(void) {
  Client c1;
  c1.accountNumber = 91524;
  c1.balance = 0.0;
  Node<Client> n;
  n.set(c1);
  cout << n.get();
  return 0;
}
```

**Compiling command: g++ driver.cpp**

There are 2 problems that have been fixed above:

1. `driver.cpp` must `#include` the node file that contains the bodies of the functions. (See my web posting if you're still unclear on this issue.) Stylistically, it would be better to call that node file "NodeImp.h", but stylistic improvements were not required for this question.

2. Node's set method expects to receive a pointer to an object. This necessitates the 3 changes shown inside main.
Part F [4 marks in total]

Suppose we have a B-tree of order 5 (i.e., \( M = 5 \)), and that the tree has nodes on three levels — a root node and two levels below it.

(a) What is the minimum number of leaf nodes that this tree might have? **6**

(b) What is the maximum number of leaf nodes that this tree might have? **25**

(c) What is the minimum number of keys that this tree might have at the leaf level? **18**

(d) What is the maximum number of keys that this tree might have at the leaf level? **125**

Explanation: (No explanation was required for full marks)

Smallest possible tree with 3 levels has 6 leaves, and between 18 and 30 keys at the leaves.

Biggest possible tree with 3 levels has 25 leaves, and between 75 and 125 keys at the leaves.