

Near Optimal Dimensionality Reductions that Preserve Volumes

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Dimension Reduction

$P \subseteq \mathbb{R}^t$: set of n points

Goal: Find $f : P \rightarrow \mathbb{R}^d$ ($d \ll n, t$) s.t. some property is preserved.

Measure of quality (Distance)

f has *distortion* $1 + \varepsilon$ if

$$\forall p, q \in P \quad \|p - q\| \leq \|f(p) - f(q)\| \leq (1 + \varepsilon) \|p - q\|.$$

Measure of quality (Volume)

f has *volume distortion* $1 + \varepsilon$ if

$$\forall S \subset P, |S| \leq k \quad 1 \leq \left(\frac{\text{vol}(f(S))}{\text{vol}(S)} \right)^{\frac{1}{|S|-1}} \leq 1 + \varepsilon.$$

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Measure of quality (Volume) (This talk)

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Johnson Lindenstrauss Lemma

Lemma (Distances)

Let P an n -point subset of Euclidean space. There exists a mapping f from P into \mathbb{R}^d , $d = O(\varepsilon^{-2} \log n)$ such that

$$\forall x, y \in P \quad (1 - \varepsilon)\|x - y\| \leq \|f(x) - f(y)\| \leq (1 + \varepsilon)\|x - y\|$$

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Almost tight Lower bound $\Omega(\varepsilon^{-2} \log n / \log(1/\varepsilon))$ [Alon, 2003].

Random Projections

Many ways to generate such a (linear) mapping (encoded by $X \in \mathbb{R}^{n \times d}$):

- $X_{i,j} \sim N(0, 1)$
- $X_{i,j} \sim \pm 1$ w.p. $1/2$.
- Sparse Gaussian matrix (with preprocessing)
- Entries with Subgaussian tails
- ECC and Rademacher r.v.
- Lean Walsh Transform. (Next talk, [Liberty et al., 2008])

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Related Work: Extensions of JL to other cases

[Magen, 2002] Preserve volumes of subsets of size up to k and affine distances using $O(k\epsilon^{-2} \log n)$ dimensions.

[Sarlos, 2006] Preserve distances of *all* points lying in any k dim. linear subspace by projecting into $O(k\epsilon^{-2} \log(k/\epsilon))$ dimensions.

[Wakin and Baraniuk, 2006, Agarwal et al., 2007, Clarkson, 2008] Moving points, curves, surfaces and manifolds etc.

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Our Contribution

- Improve Magen's result for volumes, by showing that $O(\max\{k/\epsilon, \epsilon^{-2} \log n\})$ dimensions are enough.
- JL Lemma preserves more than distances. It preserves volumes of subsets of size up to $\log n/\epsilon$.

Our Result

Theorem

Let $P \subset \mathbb{R}^n$. There $\exists f : P \rightarrow \mathbb{R}^d$, $d = O(\max\{\frac{k}{\varepsilon}, \varepsilon^{-2} \log n\})$, s.t. \forall subset S of P , $1 < |S| < k$,

$$1 - \varepsilon \leq \left(\frac{\text{vol}(f(S))}{\text{vol}(S)} \right)^{\frac{1}{|S|-1}} \leq 1 + \varepsilon.$$

Overview of proof:

- There are roughly $O(n^s)$ sets of size s .
- It suffices to prove the failure probability for a subset of size s is roughly $e^{-\Omega(sd\varepsilon^2)}$.
- Union bound implies that a volume-preserving mapping exists.

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Overview of proof:

- There are roughly $O(n^s)$ sets of size s .
- It suffices to prove the failure probability for a subset of size s is roughly $e^{-\Omega(sd\varepsilon^2)}$. (**Core of the talk.**)
- Union bound implies that a volume-preserving mapping exists.

Proof

Two steps:

- 1 Prove it for the regular n -simplex.
- 2 Reduce the general case to the above case.

The n -simplex

- Assume input points are $\{e_1, \dots, e_n\}$.

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- Random Projection (**without normalization**)

$$\begin{bmatrix} 1 & 0 & 0 & \dots & 0 \\ 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & \ddots & & \vdots \\ \vdots & \vdots & & 1 & 0 \\ 0 & 0 & \dots & 0 & 1 \end{bmatrix}_{n \times n} \begin{bmatrix} X_{ij} \sim \mathcal{N}(0, 1) \\ \vdots \\ \vdots \\ \vdots \\ \vdots \end{bmatrix}_{n \times d}$$

The n -simplex

- Assume input points are $\{e_1, \dots, e_n\}$.
- Projected points are **Random Gaussian Vectors** in \mathbb{R}^d .

$$\left[\begin{array}{c} X_{ij} \sim \mathcal{N}(0, 1) \\ \vdots \\ \vdots \end{array} \right]_{n \times d}$$

The n -simplex

- Assume input points are $\{e_1, \dots, e_n\}$.
- Projected points are **Random Gaussian Vectors** in \mathbb{R}^d .
- Pick any subset S , $|S| = s$ of rows of X

$$X_S := \left[\begin{array}{c} X_{ij} \sim \mathcal{N}(0, 1) \\ \phantom{X_{ij} \sim \mathcal{N}(0, 1)} \\ \phantom{X_{ij} \sim \mathcal{N}(0, 1)} \\ \phantom{X_{ij} \sim \mathcal{N}(0, 1)} \end{array} \right]_{s \times d}$$

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- $\text{vol}(S \cup \{\mathbf{0}\}) = \sqrt{\det(X_S X_S^T)}/s!$.

The n -simplex

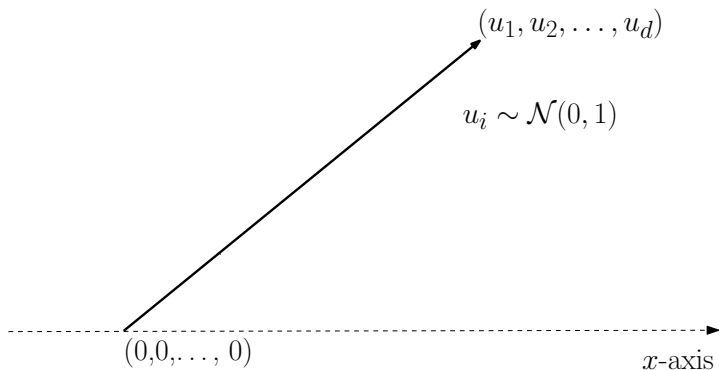
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- $\text{vol}(S \cup \{\mathbf{0}\}) = \sqrt{\det(X_S X_S^T)} / s!$.
- What's the distribution of $\sqrt{\det(X_S X_S^T)}$?

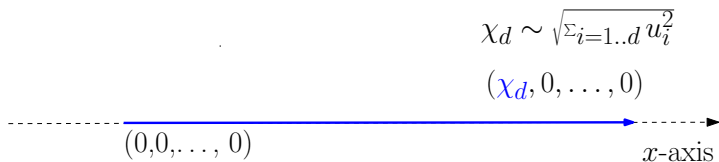
Distribution of $\sqrt{\det(\mathbf{X}_S \mathbf{X}_S^\top)}$ for $s = 2$

1) Pick a random vector



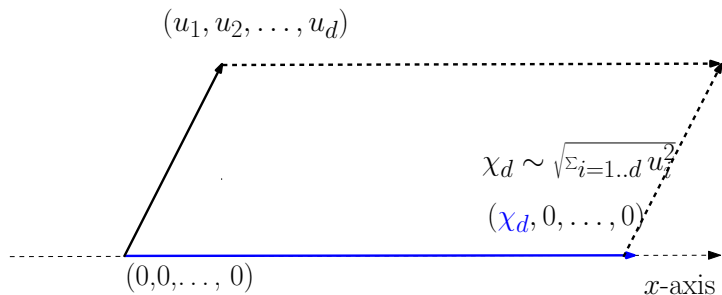
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- 2) Rotate to x -axis (Rotational invariance).



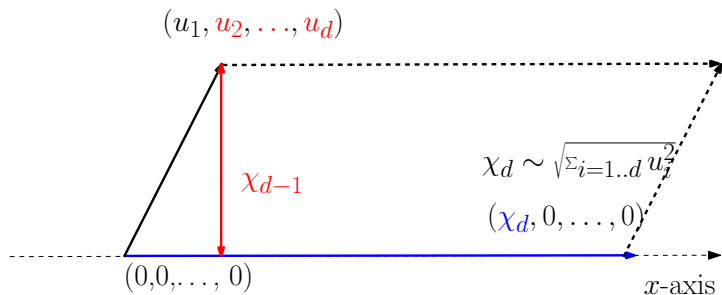
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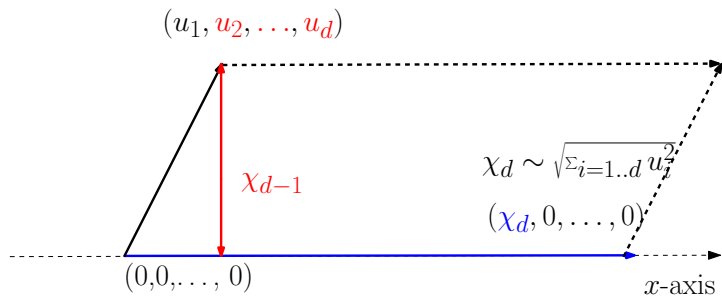
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$$\text{vol}(\diamond) \sim X_d X_{d-1}$$

Random Determinants

- For $s = 2$, $\sqrt{\det(\mathbf{X}_S \mathbf{X}_S^\top)} = \chi_d \chi_{d-1}$
- Using induction, we can show that:

Let \mathbf{X}_S s -by- d ($s \leq d$) Gaussian Random Matrix,

$$\det(\mathbf{X}_S \mathbf{X}_S^\top) \sim \prod_{i=1}^s \chi_{d-i+1}^2$$

where the Chi square r.v. are independent.

Facts about χ^2 distribution:

- $\chi_t^2 = \sum_{i=1}^t u_i^2$, $u_i \sim \mathcal{N}(0, 1)$.
- χ_t^2 is sharply concentrated around its expected value ($e^{-\Omega(t\varepsilon^2)}$).

$$\Pr\left[|\chi_t^2 - E[\chi_t^2]| > \varepsilon E[\chi_t^2]\right] < e^{-\Omega(t\varepsilon^2)}$$

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What about the concentration of the (normalized) product of χ^2 ?

Normalized Product of Independent χ^2

Theorem ([Gordon, 1989])

$$\left(\prod_{i=1}^s \chi_{d-i+1}^2 \right)^{\frac{1}{s}} \sim \frac{1}{s} \chi_{sd \pm O(s^2)}^2$$

Reminder: χ_t^2 concentrated with $e^{-\Omega(t\varepsilon^2)}$.

Implication:

- Take any subset S of the vertices of the n -simplex, call it I_S .
- $\left(\frac{\text{vol}(f(I_S))}{\text{vol}(I_S)} \right)^{\frac{1}{s}} = \left(\frac{\sqrt{\det(X_S X_S^T)}/s!}{\sqrt{\det(I_S I_S^T)}/s!} \right)^{\frac{1}{s}} = \left(\det(X_S X_S^T) \right)^{\frac{1}{2s}}$

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- The normalized volume of any subset I_S of the n -simplex is concentrated with $\exp(-\Omega(sd\varepsilon^2))$

Points in General Position

- Let P a n -point set in general position.

$$\begin{bmatrix} * & * & * & \dots & * \\ * & * & * & \dots & * \\ * & * & P_{ij} & & \vdots \\ \vdots & \vdots & & * & * \\ * & * & \dots & * & * \end{bmatrix}_{n \times n} \begin{bmatrix} \\ \\ \\ \\ \\ \end{bmatrix}_{n \times d} \quad X_{ij} \sim \mathcal{N}(0, 1)$$

Points in General Position

- Pick any subset P_S of P

$$\begin{bmatrix} * & * & * & * & * \\ * & * & P_S & * & * \\ * & * & * & * & * \\ * & * & * & * & * \end{bmatrix}_{s \times n} \quad \left[\begin{array}{c} X_{ij} \sim \mathcal{N}(0, 1) \\ \\ \\ \end{array} \right]_{n \times d}$$

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- Let $Y_S = P_S X$.
- Observation 1: Y_S is a correlated Gaussian matrix.
- Using the *stability* property of Gaussian Distribution, we can show that

$$\frac{\det(Y_S Y_S^T)}{\det(P_S P_S^T)} \sim \det(X_S X_S^T)$$

Points in General Position

- $\left(\frac{\text{vol}(Y_S)}{\text{vol}(P_S)}\right)^{\frac{1}{s}} = \left(\frac{\sqrt{\det(Y_S Y_S^T)}/s!}{\sqrt{\det(P_S P_S^T)}/s!}\right)^{\frac{1}{s}} = \left(\det(X_S X_S^T)\right)^{\frac{1}{2s}} \approx \chi_{sd}^2$
- Same distribution as n -simplex case.
- Observation 2: (Normalized) volume distribution **independent** of P_S .

Conclusion:

- For points in general position we have the desired concentration.
- Therefore, by union bound there exists a volume preserving mapping to \mathbb{R}^d .

Why k/ε on $d = O(\max\{k/\varepsilon, \varepsilon^{-2} \log n\})$?

- Let μ_s be the expected normalized volume of subsets of size s .
- μ_s is decreasing in s , $\mu_s \leq \sqrt{d-s/2}$.
- Also, $\mu_2 \geq \sqrt{d-1}$.
- We must satisfy:

$$\sqrt{\frac{d-k/2}{d-1}} \geq 1 - O(\varepsilon)$$

- Therefore, $d = \Omega(k/\varepsilon)$.

Digestion slide

	Distance (JL)	Volume (This work)
	$ S = s$	
Parameter	$s = 1$	$s < k$
Quantity (Squared & normalized)	Length	Volume
	$\frac{\ SX\ ^2}{\ S\ ^2}$	$\left(\frac{\det(SXX^T S^T)}{\det(SS^T)}\right)^{\frac{1}{s}}$
Random Variable	χ_d^2	$\left(\prod_{i=1}^s \chi_{d-i+1}^2\right)^{\frac{1}{s}}$
Concentration	$e^{-\Omega(d\varepsilon^2)}$	$e^{-\Omega(sd\varepsilon^2)}$
Lower Bound	[Alon, 2003]	?

Open Question - Lower Bound

Open Question

Let A be an $n \times n$ positive semidefinite matrix such that the determinant of every $s \times s$ principal minor ($s \leq k$) is between $(1 - \varepsilon)^{s-1}$ to 1. Then is the rank of A at least $\min\{\Omega(k/\varepsilon), n\}$?

Summary

We discussed:

- A $O(\max\{\varepsilon^{-2} \log n, k/\varepsilon\})$ upper bound for volume-preserving (to subset of size up to k) dimension reduction.
- JL Lemma preserves more than distances. It preserves volumes of subsets of size up to $\log n/\varepsilon$.
- An open question for the lower bound.

Future work:

- Extension to \mathbb{R}^d , $d = o(\log n)$.

Thank You

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Extra slides

Concentration bounds for χ^2

Lemma

Let $\chi_t^2 = \sum_{i=1}^t X_i^2$, where $X_i \sim \mathcal{N}(0, 1)$. Then for every ε , with $0 < \varepsilon \leq 1/2$, we have that

$$\Pr\left[\chi_t^2 \leq (1 - \varepsilon)E[\chi_t^2]\right] \leq \exp\left(-t\frac{\varepsilon^2}{6}\right)$$

and

$$\Pr\left[\chi_t^2 \geq (1 + \varepsilon)E[\chi_t^2]\right] \leq \exp\left(-t\frac{\varepsilon^2}{6}\right).$$

Extra slides

Theorem ([Gordon, 1989])

Let $u_i := \chi_{d-i+1}^2$ be independent Chi-square random variables for $i = 1, 2, \dots, s$. If u_i are independent, then the following holds for every $s \geq 1$,

$$\chi_{s(d-s+1) + \frac{(s-1)(s-2)}{2}}^2 \geq s \left(\prod_{i=1}^s u_i \right)^{1/s} \geq \chi_{s(d-s+1)}^2. \quad (1)$$

- By $X \geq Y$ we denote that the r.v. X is stochastically greater than Y .