

# Temporally-Expressive Planning as Constraint Satisfaction Problems

Yuxiao Hu

Department of Computer Science  
University of Toronto

yuxiao (a) cs toronto edu

September 25, 2007

# Motivation

As shown by Cushing *et al.* (2007), there are “temporally-expressive” planning problems that

- can be represented by PDDL 2.x
- cannot be solved by many *state-of-the-art* planners

This is due to their strong assumptions on

- 1 temporal annotation  
(Over-all preconditions, at end effects)
- 2 decision epochs  
(An action can happen only when another event is happening)

# Motivation

As shown by Cushing *et al.* (2007), there are “temporally-expressive” planning problems that

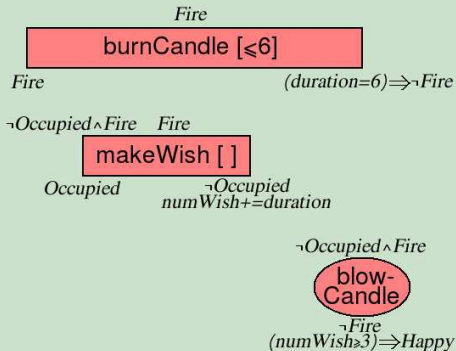
- can be represented by PDDL 2.x
- cannot be solved by many *state-of-the-art* planners

This is due to their strong assumptions on

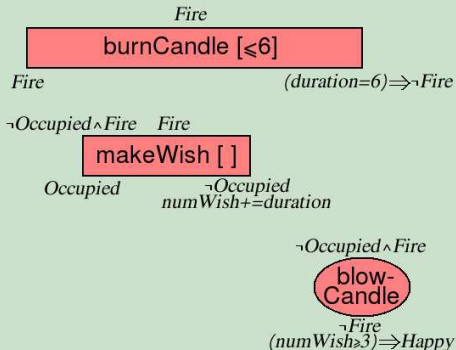
- 1 temporal annotation  
(Over-all preconditions, at end effects)
- 2 decision epochs  
(An action can happen only when another event is happening)

Goal: Solve general PDDL planning problems with a unified approach

# A Running Example



# A Running Example



- Required concurrency (“makeWish” be contained within “burnCandle”)
- Duration inequality and duration-related effects

# Approach

Our approach consists of two steps:

① PDDL  $\implies$  BAT (Basic Action Theory)

Based on a concurrent extension to the situation-calculus semantics of PDDL (Claßen *et al.* 2007)

② BAT  $\implies$  CSP

Encode the BAT into a CSP problem, and solve the CSP to obtain the plan.

# Approach

Our approach consists of two steps:

① PDDL  $\implies$  BAT (Basic Action Theory)

Based on a concurrent extension to the situation-calculus semantics of PDDL (Claßen *et al.* 2007)

② BAT  $\implies$  CSP

Encode the BAT into a CSP problem, and solve the CSP to obtain the plan.

The intuition behind it is to

- model durative actions with simple (instantaneous) actions
- treat time as a numerical property, and advance it with constraints.

# Logical Foundations

Concurrent temporal situation calculus (Reiter 2001) with the following syntax

- $A(\vec{x}, t)$  denotes the happening of simple action  $A(\vec{x})$  at time  $t$
- A durative action  $A(\vec{x})$  is represented by
  - Action  $start(A(\vec{x}), t)$ : the start event of  $A(\vec{x})$  at time  $t$
  - Action  $end(A(\vec{x}), t)$ : the end event of  $A(\vec{x})$  at time  $t$
  - Predicate  $Performing(A(\vec{x}))$ : whether  $A(\vec{x})$  is in progress
  - Function  $since(A(\vec{x}))$ : the last starting time of  $A(\vec{x})$
- $\{a_1, \dots, a_n\}$  means the concurrent happening of  $a_i$
- $[c]\alpha$  means  $\alpha$  holds after a list  $c$  of concurrent actions
- $\Box\alpha$  means  $\alpha$  holds in any situation
- $now$  is a special functional fluent representing the current time



# The Basic Action Theory

The basic action theory  $\Sigma$  consists of

- The initial database  $\Sigma_0$  e.g.:  
 $\neg \text{Fire}, \text{numWish} = 0, \neg \text{Performing}(\text{burnCandle}), \text{now} = 0;$
- The precondition axiom  $\Sigma_{pre}$ , obtained from, e.g.:  
 $\Box \text{Poss}(\text{end}(\text{burnCandle}, t)) \supset (t - \text{since}(\text{burnCandle})) \leq 6;$
- The successor state axioms  $\Sigma_{post}$ , e.g.:  
 $\Box [c] \text{Fire} \equiv \exists t. \text{start}(\text{burnCandle}, t) \in c \vee$   
 $\text{Fire} \wedge \neg (\exists t. \text{end}(\text{burnCandle}, t) \in c \wedge (t - \text{since}(\text{burnCandle})) = 6) \vee$   
 $\exists t. \text{blowCandle}(t) \in c;$
- The unique names axioms  $\Sigma_{una}$ ;
- The foundational axioms  $\mathcal{FA}$ , e.g.:  
 $\Box \text{Poss}(c) \supset \text{now} < \text{time}(c)$

# The Variable Structure

Planning by encoding the basic action theory into a CSP

- Search for increasing plan length  $n = 1, 2, 3, \dots$ , where “length” means the number of concurrent happenings.
- When searching for a plan of length  $n$ , create
  - $n$  boolean variables for each ground action term  $A(\vec{o})$ :  
 $A_{\vec{o}}^{(0)}, \dots, A_{\vec{o}}^{(n-1)}$
  - $n + 1$  boolean variables for each ground predicate  $P(\vec{o})$ :  
 $P_{\vec{o}}^{(0)}, \dots, P_{\vec{o}}^{(n)}$
  - $n + 1$  numerical variables for each ground function  $f(\vec{o})$ :  
 $f_{\vec{o}}^{(0)}, \dots, f_{\vec{o}}^{(n)}$

# The Variable Structure

An example of searching for a plan of length 2:

*Fire*<sup>(0)</sup>  
*Occupied*<sup>(0)</sup>  
*Happy*<sup>(0)</sup>  
*numWish*<sup>(0)</sup>  
*Performingburn*<sup>(0)</sup>  
*Performingwish*<sup>(0)</sup>  
*sinceburn*<sup>(0)</sup>  
*sincewish*<sup>(0)</sup>  
*now*<sup>(0)</sup>

**Fact layer 0**

*startburn*<sup>(0)</sup>  
*endburn*<sup>(0)</sup>  
*startwish*<sup>(0)</sup>  
*endwish*<sup>(0)</sup>  
*blowC*<sup>(0)</sup>

**Action layer 0**

*Fire*<sup>(1)</sup>  
*Occupied*<sup>(1)</sup>  
*Happy*<sup>(1)</sup>  
*numWish*<sup>(1)</sup>  
*Performingburn*<sup>(1)</sup>  
*Performingwish*<sup>(1)</sup>  
*sinceburn*<sup>(1)</sup>  
*sincewish*<sup>(1)</sup>  
*now*<sup>(1)</sup>

**Fact layer 1**

*startburn*<sup>(1)</sup>  
*endburn*<sup>(1)</sup>  
*startwish*<sup>(1)</sup>  
*endwish*<sup>(1)</sup>  
*blowC*<sup>(1)</sup>

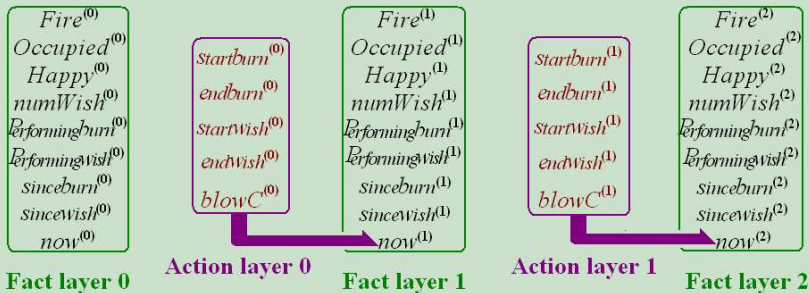
**Action layer 1**

*Fire*<sup>(2)</sup>  
*Occupied*<sup>(2)</sup>  
*Happy*<sup>(2)</sup>  
*numWish*<sup>(2)</sup>  
*Performingburn*<sup>(2)</sup>  
*Performingwish*<sup>(2)</sup>  
*sinceburn*<sup>(2)</sup>  
*sincewish*<sup>(2)</sup>  
*now*<sup>(2)</sup>

**Fact layer 2**

# The Variable Structure

An example of searching for a plan of length 2:



## Constraints: Initial and Goal States

The 0th fact layer encodes the initial state

- Problem-specific fluents set according to the initial description
- Auxiliary fluents set to 0 (or `FALSE`)

The last fact layer encodes the goal condition

- Problem-specific goals
- All “Performing” variables must be false

# Constraints: Initial and Goal States

The 0th fact layer encodes the initial state

- Problem-specific fluents set according to the initial description
- Auxiliary fluents set to 0 (or FALSE)

The last fact layer encodes the goal condition

- Problem-specific goals
- All “Performing” variables must be false

*Fire*<sup>(0)</sup>=0  
*Occupied*<sup>(0)</sup>=0  
*Happy*<sup>(0)</sup>=0  
*numWish*<sup>(0)</sup>=0  
*Performingburn*<sup>(0)</sup>=0  
*Performingwish*<sup>(0)</sup>=0  
*sinceburn*<sup>(0)</sup>=0  
*sincewish*<sup>(0)</sup>=0  
*now*<sup>(0)</sup>=0

*startburn*<sup>(0)</sup>  
*endburn*<sup>(0)</sup>  
*startwish*<sup>(0)</sup>  
*endwish*<sup>(0)</sup>  
*blowC*<sup>(0)</sup>

*Fire*<sup>(1)</sup>  
*Occupied*<sup>(1)</sup>  
*Happy*<sup>(1)</sup>  
*numWish*<sup>(1)</sup>  
*Performingburn*<sup>(1)</sup>  
*Performingwish*<sup>(1)</sup>  
*sinceburn*<sup>(1)</sup>  
*sincewish*<sup>(1)</sup>  
*now*<sup>(1)</sup>

*startburn*<sup>(1)</sup>  
*endburn*<sup>(1)</sup>  
*startwish*<sup>(1)</sup>  
*endwish*<sup>(1)</sup>  
*blowC*<sup>(1)</sup>

*Fire*<sup>(2)</sup>  
*Occupied*<sup>(2)</sup>  
*Happy*<sup>(2)</sup>=1  
*numWish*<sup>(2)</sup>  
*Performingburn*<sup>(2)</sup>=0  
*Performingwish*<sup>(2)</sup>=0  
*sinceburn*<sup>(2)</sup>  
*sincewish*<sup>(2)</sup>  
*now*<sup>(2)</sup>

# Constraints: Action Preconditions

For each formula of the form  $\Box Poss(A) \supset \pi_A$ , construct an action precondition constraint, e.g.:

From (part of) precondition axiom in the BAT

$$\Box Poss(end(burnCandle, t)) \supset (t - since(burnCandle)) \leq 6$$

we obtain the action precondition constraint

$$e_{nd}burn^{(i)} \supset (now^{(i+1)} - sinceburn^{(i)}) \leq 6$$

# Constraints: Action Preconditions

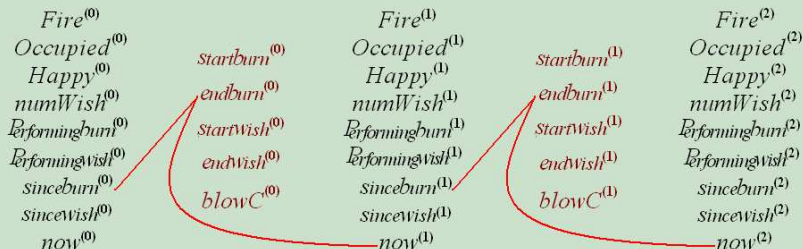
For each formula of the form  $\Box Poss(A) \supset \pi_A$ , construct an action precondition constraint, e.g.:

From (part of) precondition axiom in the BAT

$$\Box Poss(end(burnCandle, t)) \supset (t - since(burnCandle)) \leq 6$$

we obtain the action precondition constraint

$$e_{nd}burn^{(i)} \supset (now^{(i+1)} - sinceburn^{(i)}) \leq 6$$





# Constraints: Successor States

For each formula of the form  $\Box[c]F \equiv \Phi_F$ , construct a successor state constraint, e.g.:

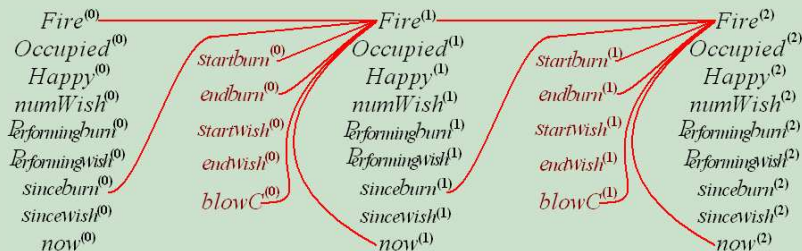
From the successor state axiom in the BAT

$$\begin{aligned}\Box[c]Fire &\equiv \exists t.start(burnCandle, t) \in c \vee \\ &Fire \wedge \neg(\exists t.end(burnCandle, t) \in c \wedge (t - since(burnCandle) = 6) \vee \\ &\exists t.blowCandle(t) \in c)\end{aligned}$$

we obtain the successor state constraint

$$\begin{aligned}Fire^{(i+1)} &\equiv s_{start}burn^{(i)} \vee \\ &Fire^{(i)} \wedge \neg(e_{nd}burn^{(i)} \wedge (now^{(i+1)} - sinceburn^{(i)} = 6) \vee blowC^{(i)})\end{aligned}$$

# Constraints: Successor States



$$Fire^{(i+1)} \equiv startburn^{(i)} \vee$$

$$Fire^{(i)} \wedge \neg(e_{nd}burn^{(i)} \wedge (now^{(i+1)} - sinceburn^{(i)} = 6) \vee blowC^{(i)})$$

# Constraints: Action Happening Times

To ensure chronological order of happenings, the foundational axiom in BAT

$$\square Poss(c) \supset now < time(c)$$

is encoded as the action happening time constraint

$$now^{(i)} < now^{(i+1)}$$

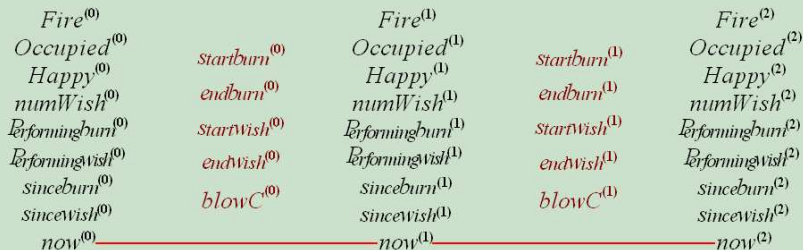
# Constraints: Action Happening Times

To ensure chronological order of happenings, the foundational axiom in BAT

$$\Box Poss(c) \supset now < time(c)$$

is encoded as the action happening time constraint

$$now^{(i)} < now^{(i+1)}$$



## Other Constraints

In addition, we need the following constraints (details in paper)

- Invariant constraints;
- Non-null step constraints;
- Mutex constraints;
- Timed-initial literal enforcement constraints.

# Implementation and Test Result

We hand-encoded several problem with constraint programming language *Choco* ([choco.sourceforge.net](http://choco.sourceforge.net)), which supports

- 1 Higher-order constraints, e.g. `implies(A, and(B,C))`;
- 2 Numerical constraints.

Can solve problems with required concurrency, duration inequalities and duration-related effects.

# Example Result

$Fire^{(0)}=0$   
 $Occupied^{(0)}=0$   
 $Happy^{(0)}=0$   
 $numWish^{(0)}=0$   
 $Performingburn^{(0)}=0$   
 $Performingwish^{(0)}=0$   
 $sinceburn^{(0)}=0$   
 $sincewish^{(0)}=0$   
 $now^{(0)}=0$

$startburn^{(0)}=1$   
 $endburn^{(0)}=0$   
 $startwish^{(0)}=0$   
 $endwish^{(0)}=0$   
 $blowC^{(0)}=0$

$Fire^{(1)}=1$   
 $Occupied^{(1)}=0$   
 $Happy^{(1)}=0$   
 $numWish^{(1)}=0$   
 $Performingburn^{(1)}=1$   
 $Performingwish^{(1)}=0$   
 $sinceburn^{(1)}=1$   
 $sincewish^{(1)}=0$   
 $now^{(1)}=1$

$startburn^{(1)}=0$   
 $endburn^{(1)}=0$   
 $startwish^{(1)}=1$   
 $endwish^{(1)}=0$   
 $blowC^{(1)}=0$

$Fire^{(2)}=1$   
 $Occupied^{(2)}=1$   
 $Happy^{(2)}=0$   
 $numWish^{(2)}=0$   
 $Performingburn^{(2)}=1$   
 $Performingwish^{(2)}=1$   
 $sinceburn^{(2)}=1$   
 $sincewish^{(2)}=2$   
 $now^{(2)}=2$

$startburn^{(2)}=0$   
 $endburn^{(2)}=0$   
 $startwish^{(2)}=0$   
 $endwish^{(2)}=1$   
 $blowC^{(2)}=0$

$Fire^{(3)}=1$   
 $Occupied^{(3)}=0$   
 $Happy^{(3)}=0$   
 $numWish^{(3)}=3$   
 $Performingburn^{(3)}=1$   
 $Performingwish^{(3)}=0$   
 $sinceburn^{(3)}=1$   
 $sincewish^{(3)}=2$   
 $now^{(3)}=5$

$startburn^{(3)}=0$   
 $endburn^{(3)}=1$   
 $startwish^{(3)}=0$   
 $endwish^{(3)}=0$   
 $blowC^{(3)}=1$

$Fire^{(4)}=0$   
 $Occupied^{(4)}=0$   
 $Happy^{(4)}=1$   
 $numWish^{(4)}=3$   
 $Performingburn^{(4)}=0$   
 $Performingwish^{(4)}=0$   
 $sinceburn^{(4)}=1$   
 $sincewish^{(4)}=2$   
 $now^{(4)}=6$

# Example Result

$Fire^{(0)}=0$   
 $Occupied^{(0)}=0$   
 $Happy^{(0)}=0$   
 $numWish^{(0)}=0$   
 $Performingburn^{(0)}=0$   
 $Performingwish^{(0)}=0$   
 $sinceburn^{(0)}=0$   
 $sincewish^{(0)}=0$   
 $now^{(0)}=0$

$startburn^{(0)}=1$   
 $endburn^{(0)}=0$   
 $startwish^{(0)}=0$   
 $endwish^{(0)}=0$   
 $blowC^{(0)}=0$

$Fire^{(1)}=1$   
 $Occupied^{(1)}=0$   
 $Happy^{(1)}=0$   
 $numWish^{(1)}=0$   
 $Performingburn^{(1)}=1$   
 $Performingwish^{(1)}=0$   
 $sinceburn^{(1)}=1$   
 $sincewish^{(1)}=0$   
 $now^{(1)}=1$

$startburn^{(1)}=0$   
 $endburn^{(1)}=0$   
 $startwish^{(1)}=1$   
 $endwish^{(1)}=0$   
 $blowC^{(1)}=0$

$Fire^{(2)}=1$   
 $Occupied^{(2)}=1$   
 $Happy^{(2)}=0$   
 $numWish^{(2)}=0$   
 $Performingburn^{(2)}=1$   
 $Performingwish^{(2)}=1$   
 $sinceburn^{(2)}=1$   
 $sincewish^{(2)}=2$   
 $now^{(2)}=2$

$startburn^{(2)}=0$   
 $endburn^{(2)}=0$   
 $startwish^{(2)}=0$   
 $endwish^{(2)}=1$   
 $blowC^{(2)}=0$

$Fire^{(3)}=1$   
 $Occupied^{(3)}=0$   
 $Happy^{(3)}=0$   
 $numWish^{(3)}=3$   
 $Performingburn^{(3)}=1$   
 $Performingwish^{(3)}=0$   
 $sinceburn^{(3)}=1$   
 $sincewish^{(3)}=2$   
 $now^{(3)}=5$

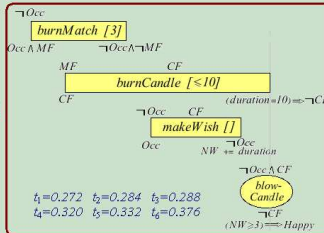
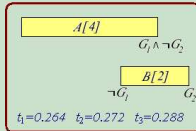
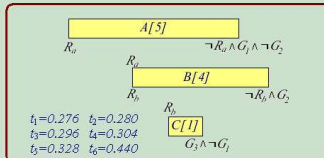
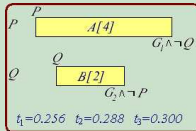
$startburn^{(3)}=0$   
 $endburn^{(3)}=1$   
 $startwish^{(3)}=0$   
 $endwish^{(3)}=0$   
 $blowC^{(3)}=1$

$Fire^{(4)}=0$   
 $Occupied^{(4)}=0$   
 $Happy^{(4)}=1$   
 $numWish^{(4)}=3$   
 $Performingburn^{(4)}=0$   
 $Performingwish^{(4)}=0$   
 $sinceburn^{(4)}=1$   
 $sincewish^{(4)}=2$   
 $now^{(4)}=6$

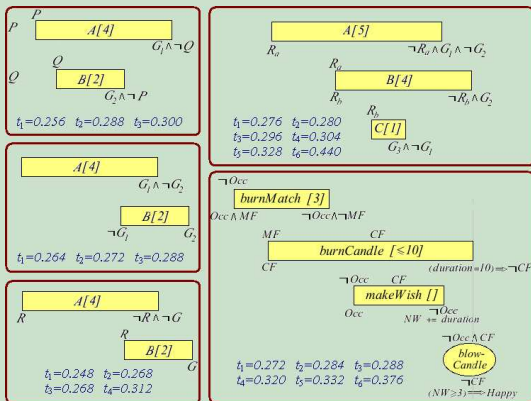
PDDL plan:  $\{(1 : burnCandle[5]), (2 : makeWish[3]), (6 : blowCandle)\}$



# Other Experiments



# Other Experiments



- No comparison with *state-of-the-art* planners available yet.
- Presumably slower on existing (temporally-simple) benchmarks, since we have only focused on generality, not yet efficiency.

# Conclusion and Future Work

We have

- 1 extended a declarative semantics of PDDL with true concurrency
- 2 proposed a general solution to PDDL planning problems based on a CSP encoding of the declarative semantics
  - Handles arbitrary PDDL temporal annotations
  - Determines happening times by satisfying constraints
  - Solves **temporally-expressive** PDDL problems (possibly with **duration-related constraints** and **effects**) in a **unified search**.

# Conclusion and Future Work

We have

- 1 extended a declarative semantics of PDDL with true concurrency
- 2 proposed a general solution to PDDL planning problems based on a CSP encoding of the declarative semantics
  - Handles arbitrary PDDL temporal annotations
  - Determines happening times by satisfying constraints
  - Solves **temporally-expressive** PDDL problems (possibly with **duration-related constraints** and **effects**) in a **unified search**.

Future work:

- 1 Implementation and optimization of automatic translation
- 2 Constraints and preferences in PDDL 3.0