

A Logic for Decidable Reasoning about Services

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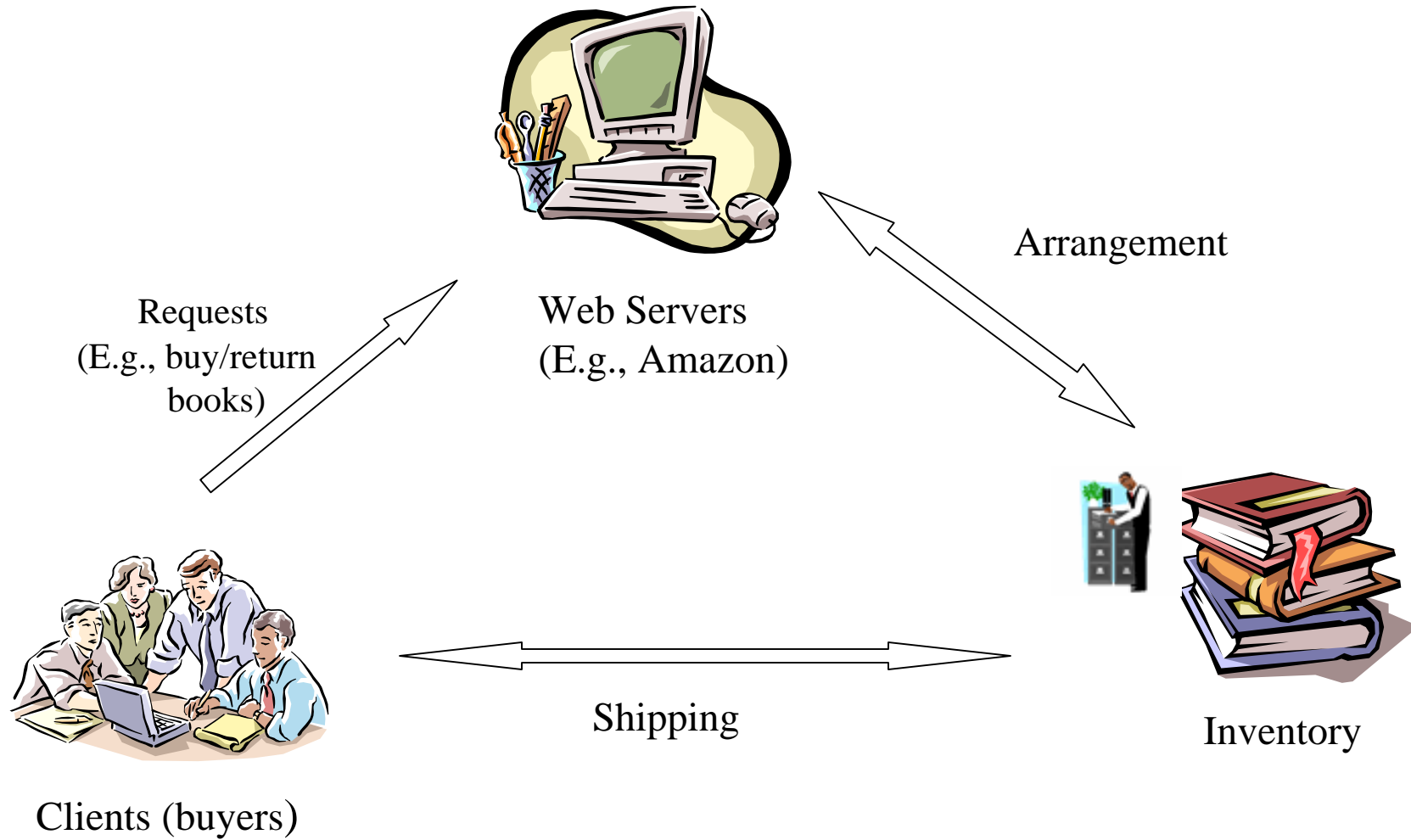
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Outline

- Motivations
- Preliminaries
- Specification of the modified situation calculus for services
- Decidable reasoning about actions in this logic
- Discussions and future work

Shopping Online



Motivations

- Usually suppliers (Web servers) could not get complete information (OWA)
- Need composition of atomic services to achieve the clients' requests
- Integrating Semantic webs with Web services
- Representing the dynamics
 - What needs to be represented?
atomic services (i.e., actions), dynamic environment (such as what books are available currently), the effect of service action
 - Expectations:
 - Represent actions for large/infinite domains (such as people, weight, time)
 - Be able to represent knowledge such as “*there exist some ...*”
- For composite services and the environment,
 - What do we care about ? (reasoning)
 - Whether the composite services are possible to be executed successfully?
 - Whether certain properties/goals can be satisfied after the execution?
 - Expectations: efficient reasoning (here, decidability)

The Situation Calculus

- A first-order logic language
 - Represent actions and effects in a natural way
 - Very compact
- Three sorts:
 - Actions: $buyBook(x,y)$, $returnBook(x,y)$, ...
 - Situations: S_0 , $do(a,s)$, $do([a_1, \dots, a_n], s)$
 - Objects: things other than actions and situations.
E.g., places, names, numbers, etc.
- Fluents: system features whose truth values may vary.
E.g., $instore(x,s)$, $boughtBook(x,y,s)$, $bought(x,y,s)$...

Basic action theory \mathcal{D}

- A set of first-order axioms to model actions and effects

- Precondition axioms for actions \mathcal{D}_{ap} :

$$Poss(buyBook(x,y),s) \equiv client(x) \wedge book(y) \wedge instore(y,s)$$

- Successor state axioms \mathcal{D}_{ss} :

$$bought(x,y,do(a,s)) \equiv a = buybook(x,y) \vee a = buyCD(x,y)$$

$$bought(x,y,s) \wedge \neg (a = returnbook(x,y) \vee a = returnCD(x,y))$$

- Axioms for initial database \mathcal{D}_{S_0} :

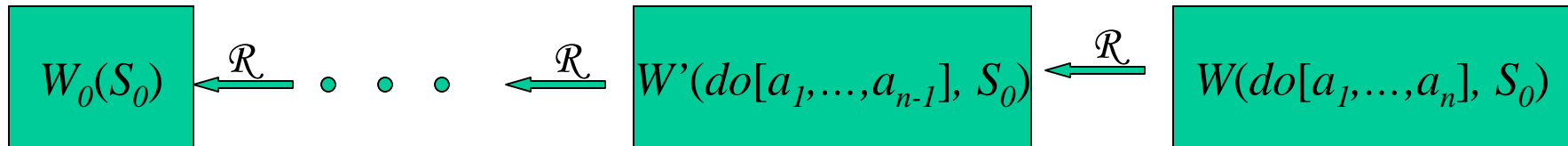
- Knowledge known to be true in the situation S_0
- Non-changeable facts
- Open world assumption

Reasoning about Actions

- E.g., $(\exists x)(\forall y)(\forall y') \text{ boughtBook}(x,y,S) \wedge \text{boughtBook}(x,y',S) \supset y=y'$
- Key reasoning mechanism -- regression operator \mathcal{R}
 - Successor state axioms support regression in a natural way

If $F(x_1, \dots, x_n, do(a, s)) \equiv \Psi_F(x_1, \dots, x_n, a, s)$, then

$$\mathcal{R}[F(t_1, \dots, t_n, do(a', S))] = \mathcal{R}[\Psi_F(t_1, \dots, t_n, a', S)].$$



- Important properties for regression
 - $\mathcal{D} \models W \equiv \mathcal{R}[W]$
 - $\mathcal{D} \models W$ iff $\mathcal{D}_{S_0} \cup \mathcal{D}_{una} \models \mathcal{R}[W]$

Disadvantages of the Situation Calculus

Advantage: representing actions and effects very compactly.

Disadvantage: reasoning for actions in general is undecidable under the open world assumption (OWA).

1. Can we get rid of the disadvantage?
2. Can we specify the Semantic Web features in a natural way?

Solution: Consider a fragment of first-order logic C^2 .

Description Logics v.s. C^2

- Description logics
 - Base of OWL
 - Different varieties
 - $\mathcal{ALCQIO}(\sqcap, \sqcup, \neg, |, id)$
- C^2 – a fragment of FOL
 - At most two variables x, y
 - No function symbols
 - Add counting quantifiers $\exists^{\geq n}, \exists^{\leq n}$
- $\mathcal{ALCQIO}(\sqcap, \sqcup, \neg, |, id)$ v.s. C^2
 - › Concept names \Leftrightarrow unary predicates *instore* ----- *instore(x)*
 - › Role names \Leftrightarrow binary predicates *boughtBook* ----- *boughtBook(x,y)*
 - › E.g., $\exists^{\geq n}R.C \Leftrightarrow \exists^{\geq n} y.R(x,y) \wedge C(y)$, $\forall R.C \Leftrightarrow \forall y.R(x,y) \supset C(y)$
 $\neg C \Leftrightarrow \neg C(x)$, $C1 \sqcap C2 \Leftrightarrow C1(x) \wedge C2(x)$

Decidability of DLs and C^2

- [Borgida1996] $\mathcal{ALCQIO}(\sqcap, \sqcup, \neg, |, id)$ plus cross-product $\Leftrightarrow C^2$.
- We showed that: $C1 \times C2 = (R \sqcup \neg R)|_{C2} \sqcap ((R \sqcup \neg R)|_{C1})^-$.
- [Grädel et al., Pacholski et al. 1997] C^2 is decidable even under OWA.

$\mathcal{ALCQIO}(\sqcap, \sqcup, \neg, |, id) \Leftrightarrow C^2$, the translation algorithm is linear to the size of the given formula.

$\mathcal{ALCQIO}(\sqcap, \sqcup, \neg, |, id)$ is decidable even under OWA.

- Other advantages
 - The features in Semantic Webs can be easily represented in C^2 .
 - The reasoning in C^2 can also be easily translated into DLs.
 - May use current existing efficient DL reasoners for C^2 formulas.

The Decidable Situation Calculus \mathcal{L}_{SC}^{DL}

- Sorts:
 - Terms of *objects* are either variable x , variable y , or constants
 - Action functions have at most two arguments
 - Variable symbol a of sort *action* and symbol s of sort *situation* are the only additional variable symbols
- Fluents with either two or three arguments:
 - (Dynamic) concepts $instore(x,s), \dots$
 - (Dynamic) roles $boughtBook(x,y,s), boughtCD(x,y,s), bought(x,y,s), \dots$
- Facts with either one or two arguments:
 - (Static) concepts $person(x), client(x), book(x), cd(x), \dots$
 - (Static) roles $hasCreditcard(x,y), \dots$
- Logic: add counting quantifiers $\exists^{\geq n}, \exists^{\leq n}$

The Basic Action Theory of \mathcal{L}_{SC}^{DL}

- Precondition axioms:
 - The RHS is C^2 if the situation argument s is suppressed
- Success state axioms:
 - Allow counting quantifiers
 - Variables a and s are free in the RHS of the axioms
 - Moreover, x, y, a and s are the only variables (both free and quantified)
- Axioms for initial databases: (with OWA)
 - Each axiom is C^2 if S_0 is suppressed

Purpose: to ensure the regression result is C^2 regardless S_0 .

Extensions of the Basic Action Theory

- Allowing specify certain features similar to DLs
- Acyclic TBox axioms:
 - Dynamic ones: $C(x,s) \equiv \Phi_C(x,s)$ (C – defined dynamic concept)
 - Static ones: $C(x) \equiv \Phi_C(x)$ (provided in the \mathcal{D}_{S0})
 - The RHS is C^2 when the situation argument s is suppressedE.g., $valCust(x,s) \equiv person(x) \wedge (\exists^{\geq 3} y) bought(x,y,s)$
 $client(x) \equiv person(x) \wedge (\exists y) hasCreditcard(x,y)$
 - Reasoning: use lazy unfolding for Dynamic ones
- RBox axioms:
 - For taxonomic reasoning purpose
 - $R1 \supset R2$ for role $R1, R2$E.g., $boughtBook(x,y,s) \supset bought(x,y,s)$, $boughtCD(x,y,s) \supset bought(x,y,s)$
 - Correctly compiled in \mathcal{D}_{SS} . I.e., $\mathcal{D} \models (\forall x,y,s).R1(x,y)[s] \supset R2(x,y)[s]$

Reasoning: Regression + Lazy Unfolding

- Expectations
 - Resulting formula should be C^2 if S_0 is suppressed
 - Be able to handle dynamic TBox axioms
- Reiter's regression operator is not suitable: introduce new variables
- Formula W that is regressable in \mathcal{L}_{SC}^{DC}
 - The situation term in W are ground
 - Variables in W can only include x, y
- Modified regression operator \mathcal{R}
 - When W is not atomic, the operator is still defined recursively
 - E.g., $\mathcal{R}[W1 \wedge W2] = \mathcal{R}[W1] \wedge \mathcal{R}[W2], \dots$
 - Add $\mathcal{R}[\exists^{\geq n} v.W] = \exists^{\geq n} v.\mathcal{R}[W]$
 - Reuse variables x and y when W is atomic ([examples on the next slide](#))
 - When W is a defined dynamic concept, use TBox axioms (lazy unfolding)

A Regression Example in \mathcal{L}_{SC}^{DL}

$A1 = buyCD(\text{Tom}, \text{BackStreetBoys}),$

$A2 = buyBook(\text{Tom}, \text{HarryPotter}),$

$A3 = buyBook(\text{Tom}, \text{TheFirm})$

$\mathcal{R}[(\exists x).valCust(x, do([A1,A2,A3],S_0))]$

$= \mathcal{R}[(\exists x). person(x) \wedge (\exists^{\geq 3} y) bought(x, y, do([A1,A2,A3], S_0))] \text{ (lazy unfolding)}$

$= (\exists x). person(x) \wedge (\exists^{\geq 3} y) \mathcal{R}[bought(x, y, do([A1,A2,A3], S_0))]$

$= \dots \text{ (recursively do regression using the successor state axioms)}$

$= (\exists x). person(x) \wedge (\exists^{\geq 3} y) [(x=\text{Tom} \wedge y = \text{TheFirm}) \vee$

$(x=\text{Tom} \wedge y = \text{HarryPotter}) \vee$

$(x=\text{Tom} \wedge y = \text{HarryPotter}) \vee$

$bought(x,y,S_0)]$

Important Properties

Suppose W is a regressable formula of \mathcal{L}_{SC}^{DL} with the basic action theory \mathcal{D}

- The regression $\mathcal{R}[W]$ terminates in a finite number of steps.
- $\mathcal{R}[W]$ is a C^2 formula if S_0 is suppressed
- $\mathcal{D} \models W \equiv \mathcal{R}[W]$
- $\mathcal{D} \models W$ iff $\mathcal{D}_{S_0} \cup \mathcal{D}_{una} \models \mathcal{R}[W]$
- The problem whether is $\mathcal{D} \models W$ decidable
 - $\mathcal{D}_{S_0} \cup \mathcal{D}_{una} \models \mathcal{R}[W]$ is a decidable reasoning in C^2 when S_0 is suppressed everywhere
- The executability problems and projection problems are decidable in \mathcal{L}_{SC}^{DL}

Discussions and Future Work

- Conclusions
 - Formalize a decidable language suitable for Web services
 - Have compact powerful expression power
- Other related researches
 - [McIlraith and Son 2002] assumes that sufficient information is available
 - [Berardi et al. 2003] uses propositional dynamic logic to model services
e-services \rightarrow constants, fluents $\rightarrow F(s)$ (propositional fragment of the situation calculus)
 - [Artale & Franconi 2001, Baader et al. 2003] extend DLs with temporal logics to capture the change of the world over time instead of caused by actions
 - [Baader et al. 2005] defines a service using a triple of sets of DL formulas
- Possible future work
 - Implementations
 - Consider the knowledge base progression/update problem in \mathcal{L}_{SC}
 - Etc.

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