

# Tutorial 2

Week of January 22, 2007

# 1 Scheme on CDF

Invoking: `scheme`

Exiting: `(exit)` or `Ctrl-D`

Loading `filename.scm`: `(load ‘‘filename’’)`  
or  
`(load ‘‘filename.scm’’)`

Tracing: `(trace proc_name)`

Transcript:

```
(transcript-on <my_transcript>)  
(transcript-off)
```

will save a transcript of a session to `<my_transcript>`.

Debugger:

```
-start: (debug)  
-help: ?  
-go back (read-eval-print level): (restart 1) or Ctrl-C Ctrl-C  
-quit: q
```

## 2 Read-Eval-Print loop

1. Read input from user
2. Evaluate input
3. Print return value

HOW DOES EVALUATION GO? Let's see...

```
1 ]=> 1  
;Value: 1
```

this means: '1' evaluates to 1

```
1 ]=> +  
;Value 1: #[arity-dispatched-procedure 1]
```

This means '+' evaluates to some addition procedure. STRESS THIS.

1 ]=> (+ 1 2)

---

| Attention!!! I see an opening bracket!!!  
|  
| Can I find a matching closing bracket? Yes. OK. Moving along.  
| (If I don't find a matching closing bracket, I'll just complain  
| and do nothing)  
|  
| Here's what I do next:  
|  
| - I look at the first thing after an opening bracket and I  
| evaluate it. I NEED it to evaluate to a procedure. If it does  
| not, I complain and do nothing.  
| Here I am looking at the string '+'. I can evaluate it to a  
| procedure which does the addition (as above). OK. Moving along.  
|  
| - Now I look at the rest of the things before the closing bracket  
| and evaluate them. If I cannot evaluate something along the way,  
| I complain and do nothing. Otherwise, move along.  
| Well, in our case '1' evaluates to 1 and '2' evaluates to 2.  
| So far, so good.  
|  
| - Now I attempt to apply the addition procedure (which I got from  
| evaluating '+') to 1 and 2, which I got from evaluating '1'  
| and '2'. It happens so that the addition procedure is defined  
| in such a way that it successfully does the job of adding 1 and  
| 2 and it says that 3 is the answer. I am done. I have evaluated  
| the expression '(+ 1 2)' to 3. So I say so.

---

;Value: 3

### 3 Defining Procedures

Not all procedures are pre-defined (like addition in the previous example). Here's how we define procedures. In what follows, we show how to define a procedure, *without* associating a name with it. There are many interesting uses for such "unnamed procedures" which we will see later in the course. Later in this tutorial we will see how to associate a name with such a procedure definition.

```
(lambda (var1 var2 ... varN) exp1 exp2 ... expM)
```

'lambda' is a {\em syntactic form}, a special keyword that tells the interpreter that what follows is a procedure definition.

var1, var2, ..., varN are the arguments to the procedure.

The rest is the body of the procedure.

All of exp1, exp2, ..., expM are evaluated; value of expM only is returned.

```
1 ]=> (lambda (x) x)
;Value 1: #[compound-procedure 1]
```

'(lambda (x) x)' got evaluated to a procedure that takes something as an argument and returns it.

```
1 ]=> ( lambda (x) (+ x x) )
;Value 1: #[compound-procedure 2]
```

'( lambda (x) (+ x x) )' got evaluated to a procedure that takes something as an argument and returns the result of evaluating '(+ something something)'.  
'(+ something something)'

```
1 ]=>( (lambda (x) (+ x x)) 5 )
;Value: 10
```

“( (lambda (x) (+ x x)) 5 )” is evaluated as follows:

- see an opening bracket
- evaluate the first thing after it:
  - “(lambda (x) (+ x x))” evaluates to a procedure as above
- “5” evaluates to 5
- the procedure is applied to 5:
  - evaluate “(+ 5 5)” :
    - evaluate “+” to an addition procedure
    - evaluate 5 to 5
    - get 10
  - return the result of evaluating “(+ 5 5)”, i.e. return 10
- print the result: “;Value: 10”

BUT NOTE:

```
1 ]=> ( (lambda (x y) (* x y) (+ x y)) 3 5)
; Value: 8
```

## 4 Giving a name to a procedure

```
(define proc-name (lambda (var1 var2 ... varN) exp1 exp2 ... expM))
```

“define”, like “lambda” is a *syntactic form*.

When I see ‘‘define’’, I know I will have 2 things following it, before the closing bracket.

Then the following happens:

- the second thing is evaluated
- the first thing gets the value obtained from evaluation

```
1 ]=> (define my_proc (lambda (x) (+ x x)))  
;Value: my_proc
```

This means:

- ‘‘(lambda (x) (+ x x))’’ got evaluated to a procedure as before
- my\_proc now has a value: it is that procedure

```
1 ]=> (my_proc 5)  
;Value: 10
```

this means: - ‘‘my\_proc’’ got evaluated to a procedure above  
- apply that procedure to 5  
- get 10 and print it

IMPORTANT NOTE:

This is no different from:

```
1 ]=> (define x 10)
;Value: x
```

```
1 ]=> x
;Value: 10
```

We evaluate “10” to 10 and assign it to x. In case of procedure definition the second thing just happens to evaluate to a procedure.

Shortcut:

```
(define (proc-name var1 ... varN) exp1 ... expM)
```

It is THE SAME AS

```
(define proc-name (lambda (var1 var2 ... varN) exp1 exp2 ... expM))
```

```
1 ]=> (define (my_proc x) (+ x x))
;Value: my_proc
```

```
1 ]=> (my_proc 5)
;Value: 10
```



## 5 Example

We saw this example in lecture. In the file called “myfile.scm” you have the following:

```
-----  
; my comments are here  
(define (increment n)  
  (+ n 1))  
  
; some other comments are here  
(define foobar 21)  
-----
```

```
werewolf:~\% scheme  
Scheme Microcode Version 14.9  
MIT Scheme running under GNU/Linux  
Type '^C' (control-C) followed by 'H' to obtain information about interrupts.  
Scheme saved on Monday June 17, 2002 at 10:03:44 PM  
  Release 7.7.1  
  Microcode 14.9  
  Runtime 15.1
```

```
1 ]=> (load "myfile")  
  
;Loading "myfile.scm" -- done  
;Value: foobar  
  
1 ]=> (increment foobar)  
  
;Value: 22
```

OK. What happened there?

- 1) Loading
- 2) Evaluation

- 1) Loading:
  - same as typing in the stuff in the interpreter:
  - we have defined the procedure ‘‘increment’’ and
  - we have defined foobar to be 21
  
- 2) Evaluating ‘‘(increment foobar)’’:
  - lookup procedure value corresponding to ‘‘increment’’:  
(lambda (n) (+ n 1))
  
  - lookup ‘‘foobar’’: 21
  
  - evaluate ‘‘( (lambda (n) (+ n 1)) 21 )’’:
    - evaluate (+ 21 1)
      - evaluate ‘‘+’’ to an addition procedure
      - evaluate 21 to 21 and 1 to 1
      - apply addition procedure to 21 and 1, get 22
  - print 22

And Scheme interpreter sort of tells you this:

```
1 ]=> (trace increment)
;Unspecified return value <---- just means ‘‘nothing is returned’’
                                it is not an error, we can procede
```

```
1 ]=> (increment foobar)
```

```
[Entering #[compound-procedure 1 increment]
  Args: 21]
[22
  <== #[compound-procedure 1 increment]
  Args: 21]
;Value: 22
```

Well, you have to learn its language...

## 6 Quoting

`quote` is another *syntactic form*. It inhibits the normal evaluation of the expression that follows it.

`(quote <expression>)` is same as `'<expression>`

`'<expression>` evaluates to `<expression>`

Some examples:

```
1 ]=> (define x 1)
;Value: x
```

```
1 ]=> 'x
;Value: x
```

```
1 ]=> x
;Value: 1
```

```
1 ]=> '1
;Value: 1
```

```
1 ]=> 1
;Value: 1
```

More interesting:

```
(define x 'y)
(define y 5)
(define z y)
```

```
(+ y 1)    evals to 6
(+ z 1)    evals to 6
(+ x 1)    evals to Error
```

EXPLAIN WHAT HAPPENS HERE IN DETAIL