

## ML Lectures (cont.)

Winter 2007

### Acknowledgments:

1. Standard ML of New Jersey website: [www.smlnj.org](http://www.smlnj.org)
2. Programming in Standard ML. by Robert Harper.
3. Concept in Programming by John C. Mitchell

### Type Synonym

We can give existing types new names.

Syntax: `type tycon = ty`

*tycon* becomes an alias (synonym) for the existing type *ty*.

```
1 -type float = real;
    type float = real

2 -type count = int and average=real;
    type count = int

3 type average = real
    ??

4 -val f: float = 2.3;
    val f=2.3: float

5 -val i:count = 3;
    val i = 3: count
```

## Type Synonym (continue...)

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But notice `float`, `real`, and `average` are all of the same base type, i.e. `real`!

```
6 -val a:average = f;  
    val a = 2.3: average
```

```
7 -val res = a+f;  
    val res = 4.6: average
```

Type synonyms make program more readable.

```
8 -type car= {make:string, built:int};  
    ??
```

```
9 -val c1: car = {make="Toyota", built=2001};  
    ??
```

```
10 -fun nextModel ({make=n,built=y}:car)= y+1;  
    val newxtModel = fn : car -> int
```

```
11 - nextModel c1;  
    ??
```

## User defined datatypes

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General Syntax:

```
datatype tycon = cons1 of ty1  
                | cons2 of ty2  
                ...  
                | consn of tyn
```

- Defines a **new** type called `tycon`.
- `tyi`'s are previously defined types..
- `consi`'s are *constructors*. They are used to create a value of `tycon` type

**Note:** "of `tyi`" is omitted if a constructor does not need any argument (such constructors are called *constants*).

## Enumerated Types

When all constructors are constants (no argument).

Example:

```
1 -datatype color = Red|Blue| Green;  
    datatype color = Blue | Green | Red  
  
2 -val c=Red; (*calling constructor Red*)  
    val c=Red: color;  
  
3 -fun colorStr(Red)= "Red"  
    | colorStr(Blue)= "Blue"  
    | colorStr(Green)= "Green";  
    val colorStr = fn : ??  
  
6 -colorStr(c);  
    val it= ??
```

## Variant Types

Can create union of different types:

```
1 -datatype number = r of real  
    | i of int;  
    datatype number= i of int | r of real  
  
3 -val n1 = i 2;  
    val n1 = i 2 : number  
  
4 -val n2 = r 3.0;  
    ??  
  
5 -val lst=[r 2.2, i 3, i 4, r 0.1];  
    val lst = [r 2.2, i 3, i 4,r 0.1]: ??  
  
6 -fun sumInts ([])=0  
    | sumInts (i x::rest)= x+ sumInts rest  
    | sumInts (r x::rest)= sumInts rest;  
    val sumInts = fn : ??  
  
9 -sumInts lst;  
    ??
```

## Recursive Types

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A datatype can be recursive: e.g. **linked list**.

```
1 -datatype llist= Nil | Node of int*llist;
  datatype llist = Nil | Node of int*llist

2 -val x = Nil;
  val x=Nil: ??

3 -val y = Node (5, Nil);
  ??
4 -val z = Node(3, Node(2,Node(1,Nil)));
  ???
(*computing the length of a linked list*)
5 -fun len Nil =0
6   |len(Node(_,rest))= 1 + len rest;
  val len = fn : ??

7 -len z;
  ??
```

## Recursive Types (continue...)

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Example: a *polymorphic* linked list

```
1 -datatype 'a llist= Nil|Node of 'a*('a llist);

2 -val x = Nil;
  val x=Nil: ??

3 -val y = Node (5, Nil);
  val y = Node (5,Nil) : ??

4 -val z = Node("Test", Node("B",Nil));
  ???
```

A binary tree where only leaves have data:

```
6 -datatype 'a tree= L of 'a
           | N of ('a tree)*('a tree);
7 -val mytree= N(L(1),N(L(2),L(3)));

8 -fun max (x,y)= if x>y then x else y;
9 -fun depth(L _)=0
10  |depth(N(ltree,rtree))=
    1+max (depth ltree, depth rtree);
```

## Mutual Recursive Types

Want to represent a tree with arbitrary #of branches.

See the diagram first ...

Defining mutually recursive datatypes (using **and**).

```
1 -datatype tree = Empty | Node of int*forest
2     and forest= Nil | Cons of tree*forest
datatype tree = Empty | Node of int * forest
datatype forest = Cons of tree * forest | Nil

3 -val t1=Node(2,Nil);
 ???
4 -val t2=Node(3,Nil);
 ???
5 -val t3=Node(7,Cons(t1,Cons(t2,Nil)));
 ???
6 -val t4=Node(5,Nil);
 ???
7 -val t5=Node(1,Nil);
 ???
8 -val t6=Node(2,Cons(t5,Cons(t4,Cons(t3,Nil))));
```

## Mutual Recursive Types: function example...

We want to count how many nodes are in a tree.

solution:  $1 + \# \text{of nodes in its subtrees (i.e. forest)}$

```
1 -fun numnodeT (Empty)=0
2     | numnodeT (Node(data,f))= 1+ numnodeF(f)
3     and
4         numnodeF(Nil) = 0
5         |numnodeF(Cons(t,f))= ???
```

```
val numnodeT = fn : tree -> int
val numnodeF = fn : forest -> int
```

(\* Note that numnodeT and numnodeF are  
mutually recursive.\*)

```
6 -numnodeT(t6)
 ???
```