Lists Revisited

Recall the Cons Cell Representation:

The *pair* or *cons cell* is the most fundamental of Scheme's structured object types.

A **list** is a sequence of **pairs**; each pair's cdr is the next pair in the sequence.

The cdr of the last pair in a **proper list** is the empty list. Otherwise the sequence of pairs forms an **improper list**. I.e., an empty list is a proper list, and and any pair whose cdr is a proper list is a proper list.

An improper list is printed in **dotted-pair notation** with a period (dot) preceding the final element of the list. A pair whose cdr is not a list is often called a **dotted pair**

Creating lists

cons vs. list: The procedure cons actually builds pairs, and there is no reason that the cdr of a pair must be a list.

The procedure list is similar to cons, except that it takes an arbitrary number of arguments and always builds a proper list.

E.g., (list 'a 'b 'c)
$$\rightarrow$$
 (a b c)

Testing for Equality

- (eq? a b): Returns #t iff a and b are the same Scheme object. (Don't use eq? with numbers!)
- (= a b): Returns #t iff a and b are numerically equal. Pre: a and b must evaluate to numbers.
- (eqv? a b): Similar to eq?, but works for numbers and characters. More expensive than eq?, however.
- (equal? a b): Returns #t iff a and b have the same structure and contents. Thus, equal? recursively tests for equality. The most expensive equality predicate.

Recommended Reading:

Dybvig $\S 6.1$, 2nd ed. (available online), or Dybvig $\S 6.2$, 3rd ed.

Testing for Equality (cont.)

The eq? predicate doesn't work for lists.

Why not?

- 1. (cons 'a '()) makes a new list
- 2. (cons 'a '()) makes a(nother) new list
- 3. eq? checks if its two args are the same
- 4. (eq? (cons 'a '()) (cons 'a '())) evaluates to () (ie, #f)

Lists are stored as pointers to the first element (car) and the rest of the list (cdr).

Symbols are stored uniquely, so eq? works on them.

Equality Checking for Lists

For lists, need a comparison procedure to check for the same **structure** in two lists. How might you write such a procedure?

- (equal? 'a 'a) evaluates to #t
- (equal? 'a 'b) evaluates to ()
- (equal? '(a) '(a)) evaluates to #t
- (equal? '((a)) '(a)) evaluates to ()

Does this really work? Hint: atoms are numbers, does this work for numbers? Play around with it and with the built-in predicate procedure equal?.

Other Useful Predicates

- (null? a): Returns #t iff a is the empty list (or #f, depending on the implementation).
- (pair? a): Returns #t iff a is a pair, *i.e.*, a cons cell.
- (number? a): Returns #t iff a is a number.
- (min list): Returns the minimum of a list of numbers.
- (max list): Returns the maximum of a list of numbers.
- (even? a): Returns #t iff a is even.

Lots more in Dybvig §6.

Recursive Procedures: Counting

```
(define (atomcount x)
  (cond ((null? x) 0)
         ((atom? x) 1)
         (else (+ (atomcount (car x))
                   (atomcount (cdr x)))))
 • (atomcount '(1 2)) \Rightarrow 2
 • (atomcount '(1 (2 (3)) (5))) \Rightarrow 4:
(at '(1 (2 (3)) (5)))
(+ (at 1) (at ((2 (3)) (5))))
(+ 1 (+ (at (2 (3))) (at ((5)))))
(+ 1 (+ (at 2) (at ((3)))) (+ (at (5)) (at ()))))
(+ 1 (+ (+ 1 (+ (at (3)) (at ()))) (+ (+ (at 5) (at ())) 0)))
(+ 1 (+ (+ 1 (+ (at 3) (at ())) 0)) (+ (+ 1 0) 0)))
(+ 1 (+ (+ 1 (+ (+ 1 0) 0)) (+ 1 0)))
(+1 (+ (+1 (+10)) 1))
(+ 1 (+ (+ 1 1) 1))
(+1(+21))
(+13)
```

This is called "car-cdr-recursion."

Efficiency Issues

Problem: Evaluating the same expression twice.

Example:

What can you do if there is no assignment statement?

Efficiency Issues

Solution 1: Bind values to parameters in a helper procedure.

Note: There is a built-in max function.

Note 2: Helper procedures are an important and useful tool!

Efficiency Issues

Solution 2: Use a let or let* construct, that binds variables to expression results.

Polymorphic and Monomorphic Functions

- Polymorphic functions can be applied to arguments of many forms
- The function length is polymorphic: it works on lists of numbers, lists of symbols, lists of lists, lists of anything
- The function square is monomorphic: it only works on numbers

Higher-Order Procedures

Procedures as input values:

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Higher-Order Procedures

Procedures as returned values:

Built-In Higher-Order Procedures:

There is a built-in procedure map. Let's define our own restricted version first....

- mymap takes two arguments: a function and a list
- mymap builds a new list whose elements are the result of applying the function to each element of the (old) list

Higher-order Procedures: map

• Example:

```
(mymap abs '(-1 2 -3 4)) \Rightarrow (1 2 3 4) (mymap (lambda (x) (+ 1 x)) '(-1 2 -3)) \Rightarrow (0 3 -2)
```

• The built-in map will produce the same results, but note that the built-in map can take more than two arguments:

```
(map cons '(a b c) '((1) (2) (3))) \Rightarrow ((a 1) (b 2) (c 3))
```

What's Wrong Here??

Why doesn't this work?

Using eval to Correct the Problem

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Limitations of Using eval

BUT: eval only works in the current definition of atomcount because numbers evaluate to themselves.

```
1 ]=> (+ 1 2 3)
; Value: 6

1 ]=> (cons '+ '(1 2 3))
; Value 12: (+ 1 2 3)

1 ]=> (eval (cons '+ '(1 2 3)) '())
; Value: 6
```

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Using eval to Evaluate Expressions

Too complicated!!

Applying Procedures with apply