CSC2542
State-Space Planning

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Acknowledgements

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Motivation

- Nearly all planning procedures are search procedures
- Different planning procedures have different search spaces
  - Two examples:
- *State-space planning*
  - Each node represents a state of the world
    - A plan is a path through the space
- *Plan-space planning*
  - Each node is a set of partially-instantiated operators, plus some constraints
    - Impose more and more constraints, until we get a plan

Outline

- State-space planning
  - Forward search*
  - Backward search*
  - Lifting
  - STRIPS
  - Block-stacking

* You'll sometimes see these referred to as "progression" and "regression" in some of the automated planning literature, but these are alternative uses (some would argue misuses) relative to existing formal definitions within the RAC literature

Forward-search(O, s_0, g)

$s \leftarrow s_0$

$\pi \leftarrow$ the empty plan

loop

if s satisfies g then return π

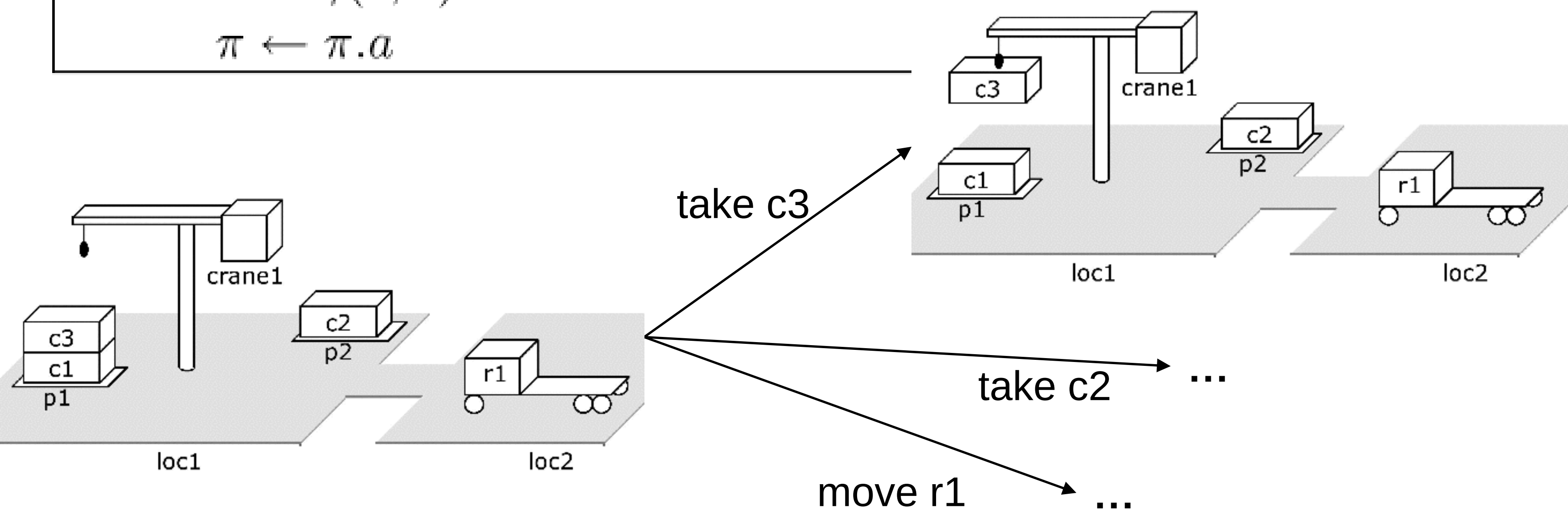
$E \leftarrow \{a | a$ is a ground instance an operator in O, and $\mathrm{precond}(a)$ is true in $s\}$

if $E = \emptyset$ then return failure

nondeterministically choose an action $a \in E$

$s \leftarrow \gamma(s, a)$

$\pi \leftarrow \pi.a$

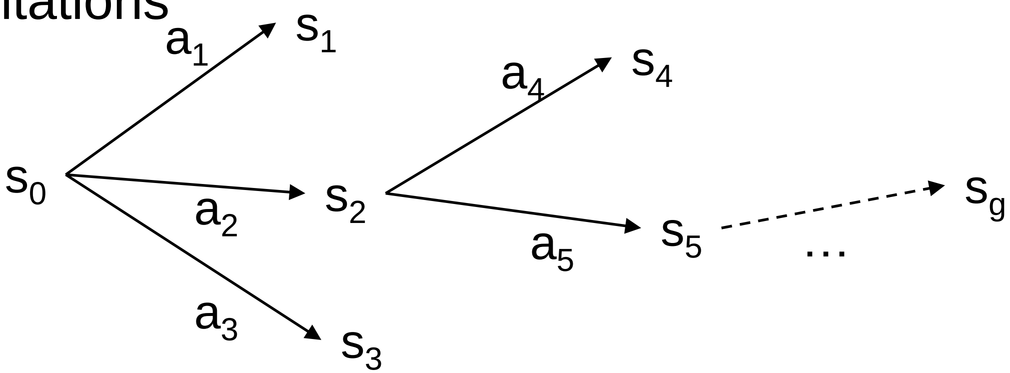


Properties

- Forward-search is *sound*
  - for any plan returned by any of its nondeterministic traces, this plan is guaranteed to be a solution
- Forward-search also is *complete*
  - if a solution exists then at least one of Forward-search's nondeterministic traces will return a solution.

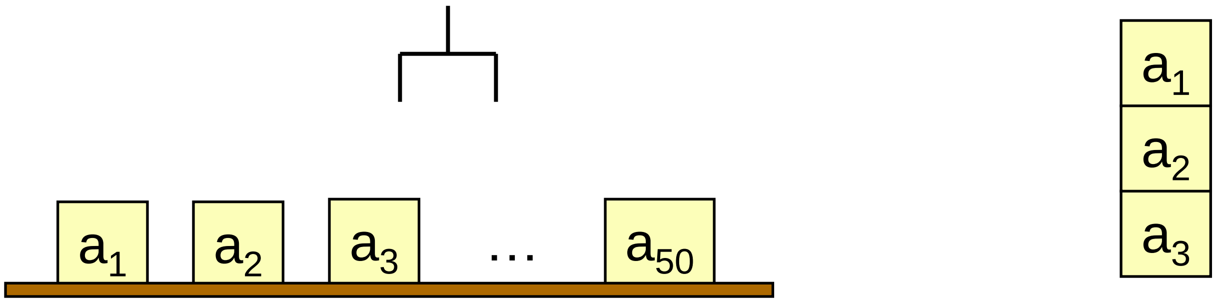
Deterministic Implementations

- Some deterministic implementations of forward search:
  - breadth-first search
  - depth-first search
  - best-first search (e.g., A*)
  - greedy search



- Breadth-first and best-first search are sound and complete
  - But they usually aren't practical, requiring too much memory
  - Memory requirement is exponential in the length of the solution
- In practice, more likely to use depth-first search or greedy search
  - Worst-case memory requirement is linear in the length of the solution
  - In general, sound but not complete
    - But classical planning has only finitely many states
    - Thus, can make depth-first search complete by doing cycle-checking

Branching Factor of Forward Search



initial state

goal

- Forward search can have a very large branching factor
  - Can have many applicable actions that don't progress toward goal
- Why this is bad:
  - Deterministic implementations can waste time trying lots of irrelevant actions
- Need a good heuristic function and/or pruning procedure

(This will be a focus of later discussion)

Backward Search

- For forward search, we started at the initial state and computed state transitions
  - new state = $\gamma(s,a)$
- For backward search, we start at the goal and compute inverse state transitions
  - new set of subgoals = $\gamma^{-1}(g,a)$
- To define $\gamma^{-1}(g,a)$, must first define *relevance*:
  - An action *a* is relevant for a goal *g* if
    - *a* makes at least one of *g*'s literals true
      - $g \cap \text{effects}(a) \neq \emptyset$
    - *a* does not make any of *g*'s literals false
      - $g^{+} \cap \text{effects}^{-}(a) = \emptyset$ and $g^{-} \cap \text{effects}^{+}(a) = \emptyset$

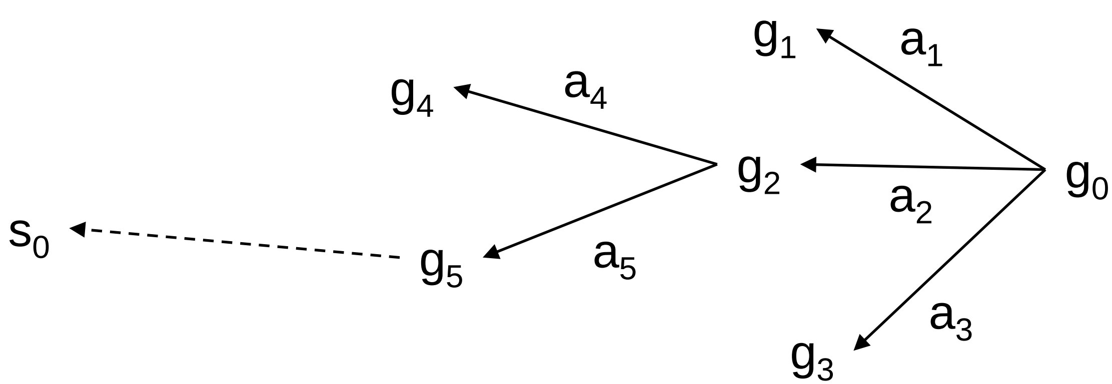
Inverse State Transitions

- If a is relevant for g, then
  - $\gamma^{-1}(g,a) = (g - \text{effects}(a)) \cup \text{precond}(a)$
- Otherwise $\gamma^{-1}(g,a)$ is undefined

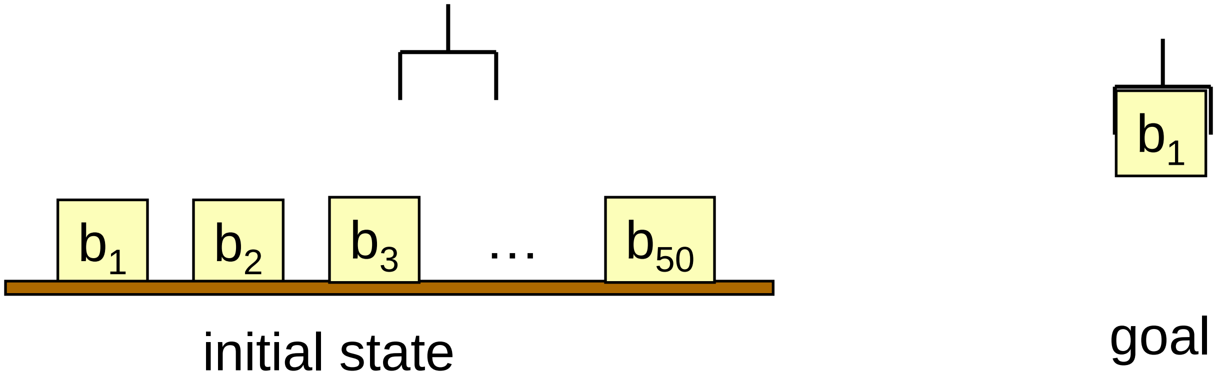
- Example: suppose that
  - g = {on(b1,b2), on(b2,b3)}
  - a = stack(b1,b2)

- What is $\gamma^{-1}(g,a)$?

```
Backward-search(O, s_0, g)
    π ← the empty plan
    loop
        if s_0 satisfies g then return π
        A ← {a | a is a ground instance of an operator in O
                  and γ^{-1}(g, a) is defined}
        if A = ∅ then return failure
        nondeterministically choose an action a ∈ A
        π ← a.π
        g ← γ^{-1}(g, a)
```



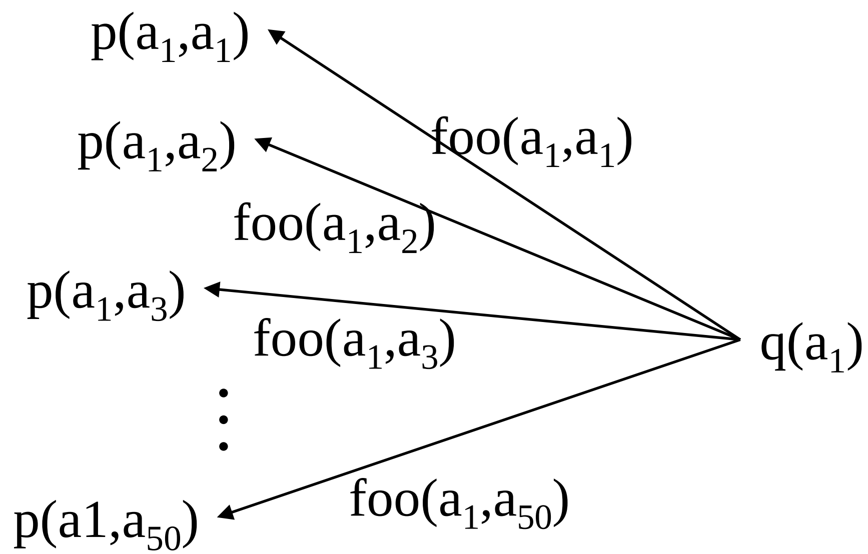
Efficiency of Backward Search



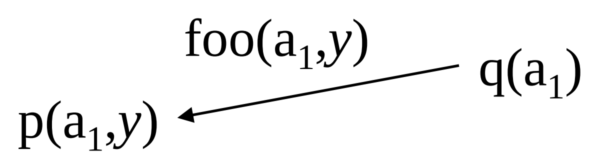
- Backward search can *also* have a very large branching factor
  - E.g., an operator o that is relevant for g may have many ground instances a_1, a_2, …, a_n such that each a_i's input state might be unreachable from the initial state
- As before, deterministic implementations can waste lots of time trying all of them

Lifting

foo(x,y)
  precond: p(x,y)
  effects: q(x)



- Can reduce the branching factor of backward search if we *partially* instantiate the operators
  - this is called *lifting*



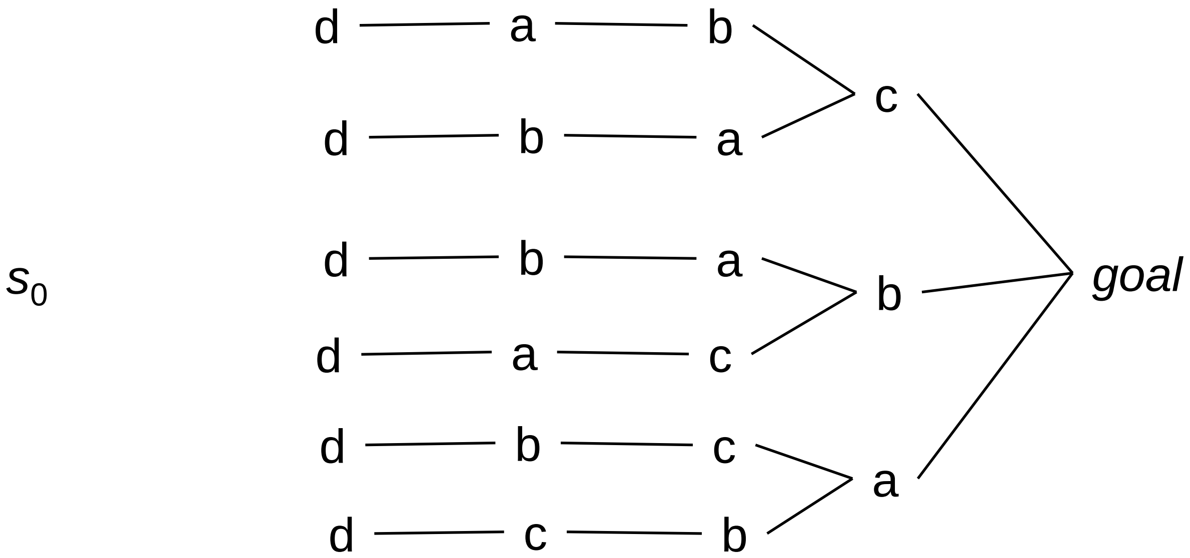
Lifted Backward Search

- More complicated than Backward-search
  - Have to keep track of what substitutions were performed
- But it has a much smaller branching factor

```
Lifted-backward-search(O, s_0, g)
    π ← the empty plan
    loop
        if s_0 satisfies g then return π
        A ← {(o, θ) | o is a standardization of an operator in O,
                  θ is an mgu for an atom of g and an atom of effects⁺(o),
                  and γ⁻¹(θ(g), θ(o)) is defined}
        if A = ∅ then return failure
        nondeterministically choose a pair (o, θ) ∈ A
        π ← the concatenation of θ(o) and θ(π)
        g ← γ⁻¹(θ(g), θ(o))
```

The Search Space is Still Too Large

- Lifted-backward-search generates a smaller search space than Backward-search, but it still can be quite large
  - Suppose actions *a*, *b*, and *c* are independent, action *d* must precede all of them, and there's no path from s_0 to *d*'s input state
  - We'll try all possible orderings of *a*, *b*, and *c* before realizing there is no solution
  - Plan-space planning or the use of "landmarks" can help with this problem



Pruning the Search Space

Pruning the search space can really help.

Two techniques we will discuss:

- Sound pruning using branch-and-bound heuristic search
- Domain customization that prunes actions and states

For now, just two examples:

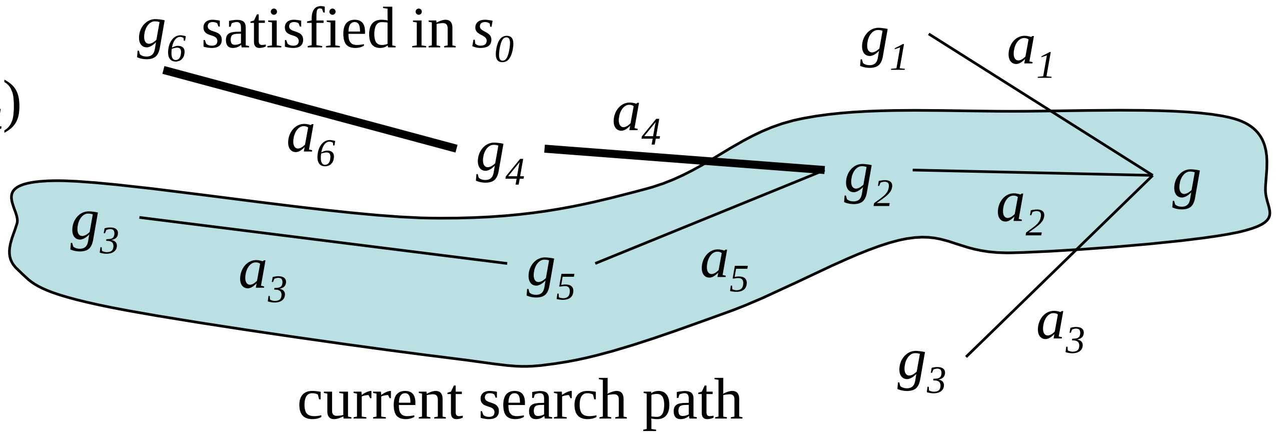
- STRIPS
- Block stacking

STRIPS

- $\pi \leftarrow$ the empty plan
- do a modified backward search from *g*
  - instead of $\gamma^{-1}(s,a)$, each new set of subgoals is just precond(*a*)
  - when you find an action that's executable in the current state, then go forward on the current search path as far as possible, executing actions and appending them to π
  - repeat until all goals are satisfied

$\pi = \langle a_6, a_4 \rangle$

$s = \gamma(\gamma(s_0,a_6),a_4)$



```
Ground-STRIPS(O, s, g)
    π ← the empty plan
    loop
        if s satisfies g then return π
        A ← {a | a is a ground instance of an operator in O,
                  and a is relevant for g}
        if A = ∅ then return failure
        nondeterministically choose any action a ∈ A
        π' ← Ground-STRIPS(O, s, precond(a))
        if π' = failure then return failure
        ;; if we get here, then π' achieves precond(a) from s
        s ← γ(s, π')
        ;; s now satisfies precond(a)
        s ← γ(s, a)
        π ← π.π'.a
```

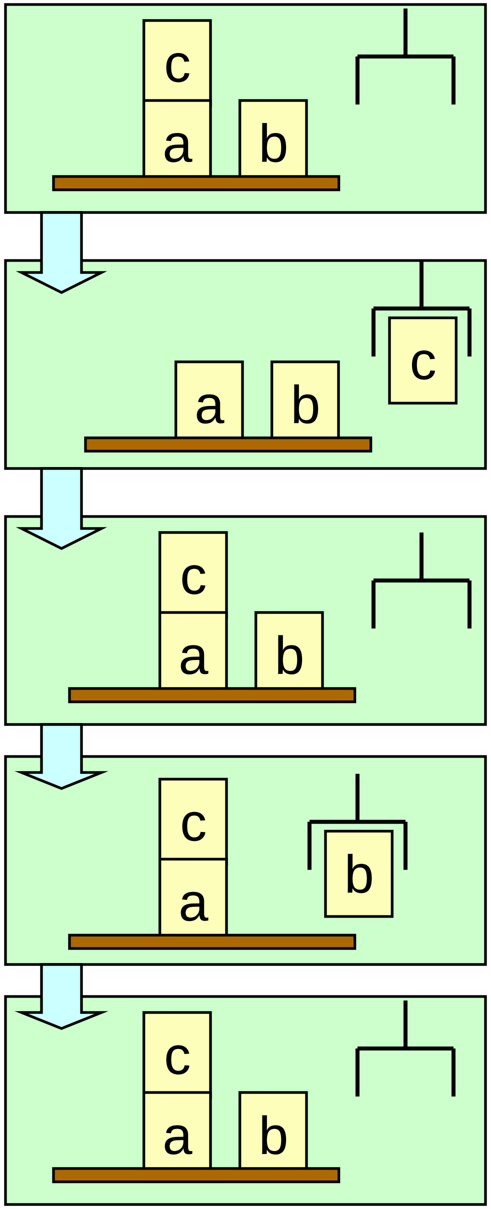
Quick Review of Blocks World

unstack(x,y)
Pre: on(x,y), clear(x), handempty
Eff: ~on(x,y), ~clear(x), ~handempty,
holding(x), clear(y)

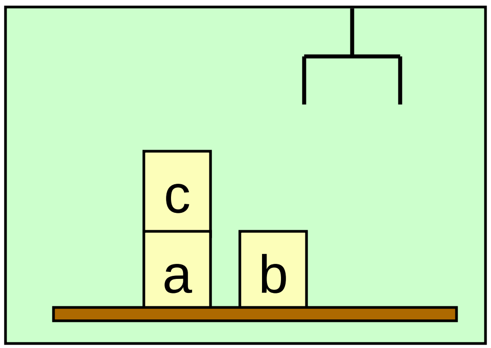
stack(x,y)
Pre: holding(x), clear(y)
Eff: ~holding(x), ~clear(y),
on(x,y), clear(x), handempty

pickup(x)
Pre: ontable(x), clear(x), handempty
Eff: ~ontable(x), ~clear(x), ~handempty, holding(x)

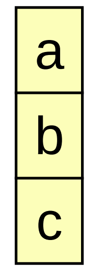
putdown(x)
Pre: holding(x)
Eff: ~holding(x), ontable(x), clear(x), handempty



The Sussman Anomaly



Initial state



goal

- On this problem, STRIPS can’t produce an irredundant solution
  - Try it and see

The Register Assignment Problem

- State-variable formulation:

  Initial state: {value(r1)=3, value(r2)=5, value(r3)=0}

  Goal: {value(r1)=5, value(r2)=3}

  Operator: assign(*r*,*v*,*r*′,*v*′)

  precond: value(*r*)=*v*, value(*r*′)=*v*′

  effects: value(*r*)=*v*′

- STRIPS cannot solve this problem at all

How to Handle Problems like These?

Several ways:

- Do something other than state-space search
  - e.g., Chapters 5–8

- Use forward or backward state-space search, with *domain-specific* knowledge to prune the search space
  - Can solve both problems quite easily this way
  - Example: block stacking using forward search

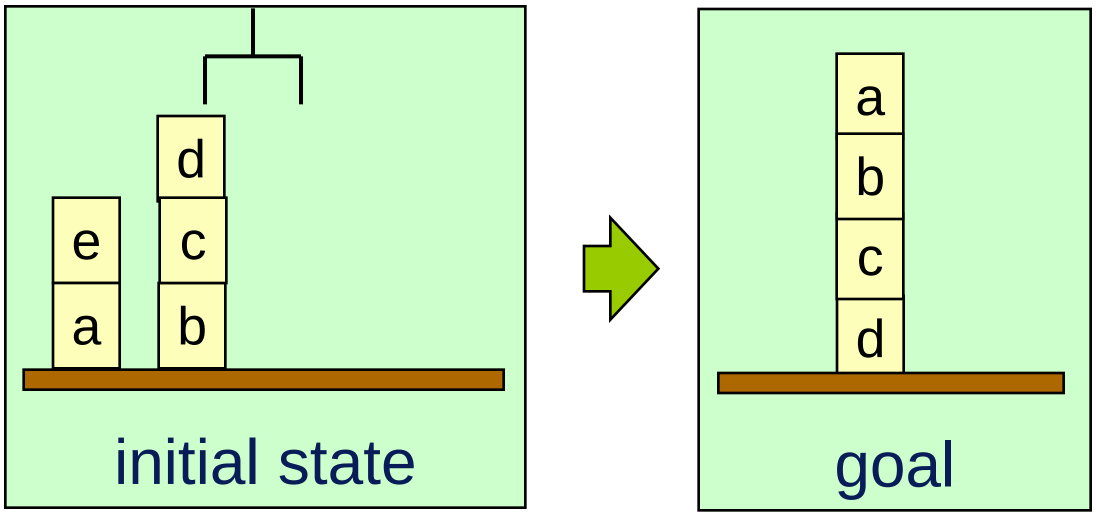
Domain-Specific Knowledge

- A blocks-world planning problem $P = (O,s_0,g)$ is solvable if s_0 and g satisfy some simple consistency conditions
    - g should not mention any blocks not mentioned in s_0
    - a block cannot be on two other blocks at once
    - etc.
        - Can check these in time O(n log n)
- If P is solvable, can easily construct a solution of length O(2m), where m is the number of blocks
    - Move all blocks to the table, then build up stacks from the bottom
        - Can do this in time O(n)
- With additional domain-specific knowledge can do even better …

Additional Domain-Specific Knowledge

A block *x* needs to be moved if any of the following is true:

- *s* contains ontable(*x*) and *g* contains on(*x,y*) - see a below
- *s* contains on(*x,y*) and *g* contains ontable(*x*) - see d below
- *s* contains on(*x,y*) and *g* contains on(*x,z*) for some *y*≠*z* - see c below
- *s* contains on(*x,y*) and *y* needs to be moved - see e below



Domain-Specific Algorithm

loop

  if there is a clear block *x* such that

    x needs to be moved **and**

    x can be moved to a place where it won't need to be moved

      then move *x* to that place

  else if there is a clear block *x* such that *x* needs to be moved

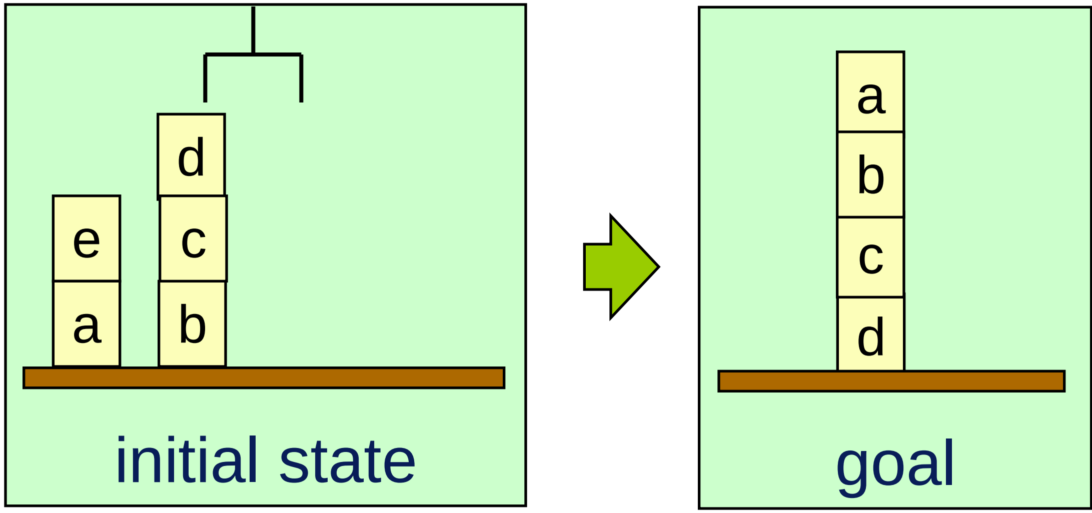
      then move *x* to the table

  else if the goal is satisfied

      then return the plan

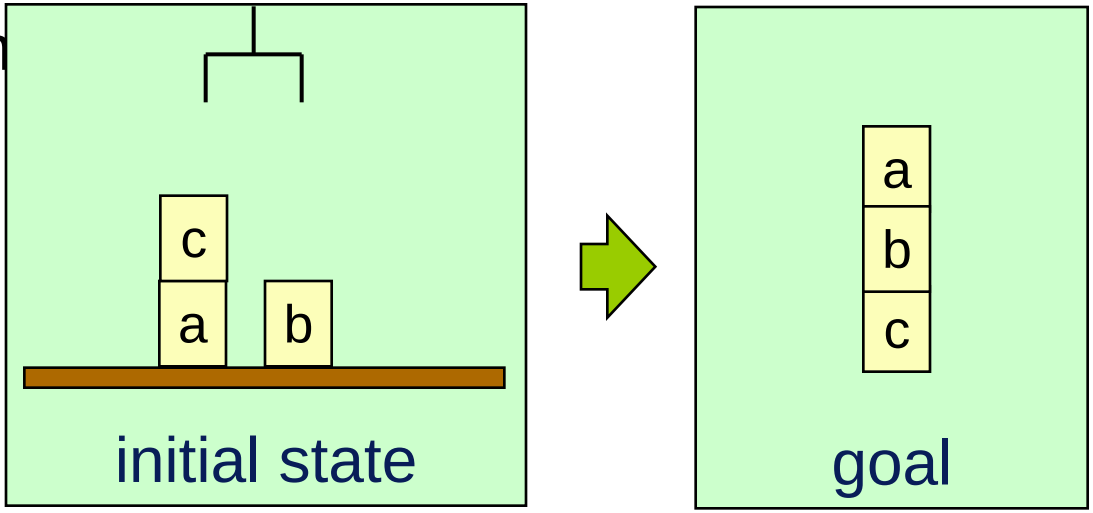
  else return failure

repeat



Easily Solves the Sussman Anomaly

loop
    if there is a clear block *x* such that
        x needs to be moved **and**
        x can be moved to a place where it won’t need to be moved
            then move *x* to that place
    else if there is a clear block *x* such that
                x needs to be moved
            then move *x* to the table
    else if the goal is satisfied
            then return the plan
    else return failure
repeat



Properties

The block-stacking algorithm:

- Sound, complete, guaranteed to terminate
- Runs in time $O(n^3)$
  - Can be modified to run in time $O(n)$
- Often finds optimal (shortest) solutions
- But sometimes only near-optimal (Exercise 4.22 in the book)

  (Note: PLAN LENGTH for the blocks world is NP-complete)