## CSC2542

Representations

## for (Classical) Planning

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## Recall:

Planning Problem
$\boldsymbol{P}=\left(\Sigma, s_{0}, G\right)$
$\Sigma$ : System Description
$s_{0}$ : Initial state(s)
E.g., Initial state $=s_{0}$

G: Objective
Goal state,
Set of goal states,
Set of tasks, "trajectory" of states,
Objective function, ...
E.g., Goal state $=s_{5}$


The Dock Worker Robots (DWR) domain

## Further Recall:

## System Description (as a state transition system)

$\Sigma=(S, A, E, \gamma)$

- $S=\{$ states $\}$
- $A=\{$ actions $\}$
- $E=\{$ exogenous events $\}$
- State-transition function $\gamma: S \times(A \cup E) \rightarrow 2^{S}$

Example: Dock Workers Robots from previous slide

- $S=\left\{\mathrm{s}_{0}, \ldots, \mathrm{~s}_{5}\right\}$
- $A=\{$ move1, move2, put, take, load, unload $\}$
- $E=\{ \}$
- $\gamma$ : as captured by the arrows mapping states and actions to successor states


## Representational Challenge

- How do we represent our planning problem is a way that supports exploration of the principles and practice of automated planning?


## Approach:

- There isn't one answer.
- The [GNT04] proposes representations that are suitable for generating classical plans.
[GNT04] = Ghallab, Nau, Traverso, Automated Planning: Theory and Practice, 2004


## Broad Perspective on Plan Representation

The right representation for the right objective.
Distinguish representation schemes for:

1. studying the principles of planning and related tasks.
2. specifying planning domains
3. direct use within (classical) planners

## Summary: Broad Perspective

1. Studying the formal principles of planning and other related task

- (First-order) logical languages
(e.g., situation calculus, $A$ languages, event calculus, fluent calculus, PDL) Properties:
- well-defined semantics, representational issues must be addressed in the language (not in the algorithm that interprets and manipulates them)
- excellent for study and proving properties. Not ideal for 3 below.

2. Specifying planning domains

- PDDL-n (PDDL2.1, PDDL2.2, PDDL3, ....)

Properties:

- (reasonably) well-defined semantics
- designed for input to planners - translate to an internal representation for specific planners. Translators exist for most state-of-the-art planners

3. Direct use within (classical) planners

- Classical representation (e.g., STRIPS)
- Set-theoretic representation (basis for rep'ns used w/ SAT solvers)
- State-variable representation (aka "Finite Domain Repn' (FDR)*") (e.g., SAS, SAS+)
Variants of these exist for particular planners (e.g., SAT solvers, model checkers, etc.)
* [Helmert, AlJ 2009]


## This Lecture:

1. Studying the formal principles of planning and other related task

- (First-order) logical languages


Properties:

- (reasonably) well-defined semantics
- designed for input to planners - translate to an internal representation for specific planners. Translators exist for most state-of-the-art planners

3. Direct use within (classical) planners (what's in the text)

- Classical representation (e.g., STRIPS)
- Set-theoretic representation (basis for rep'ns used w/ SAT solvers)
- State-variable representation (aka "FDR") (e.g., SAS, SAS+)

Variants of these exist for particular planners (e.g., SAT solvers, model checkers, etc.)

## Outline

- Representation schemes for classical planning

1. Classical representation
2. Set-theoretic representation
3. State-variable representation

- Examples: DWR and the Blocks World
- Comparisons


## Quick Review of Classical Planning

8 restrictive assumptions req'd: A0: Finite
A1: Fully observable
A2: Deterministic
A3: Static
A4: Attainment goals
A5: Sequential plans A6: Implicit time
A7: Offline planning


## Representation: Motivation for Approach

## Default view:

- represent state explicitly
- represent actions as a transition system (e.g., as an incidence matrix)


## Problem:

- explicit graph corresponding to transition system is huge
- direct manipulation of transition system is cumbersome Solution:

Provide compact representation of transition system \& induced graph

1. Explicate the structure of the "states"

- e.g., states specified in terms of state variables

2. Represent actions not as transition system/incidence matrices but as functions (e.g., operators) specified in terms of the state variables

- An action is applicable to a state when some state variables have certain values. When applicable, it will change the values of certain (other) state variables

3. To plan,

- Just give the initial state
- Use the operators to generate the other states as needed


## Why is this more compact?

## Why is this more compact than an explicit transition system?

- In an explicit transition system, actions are represented as state-tostate transitions. Each action will be represented by an incidence matrix of size $|\mathrm{S}| \mathrm{x}|\mathrm{S}|$
- In the proposed model, actions are represented only in terms of state variables whose values they care about, and whose value they affect. (It exploits the structure of the problem!)
- Consider a state space of 1024 states. It can be represented by $\log _{2} 1024=10$ state variables. If an action needs variable v 1 to be true and makes v7 to be false, it can be represented by just 2 bits (instead of a 1024x1024 matrix)
- Of course, if the action has a complicated mapping from states to states, in the worst case the action rep will be just as large
- The assumption being made here is that the actions will have effects on a small number of state variables.


## 1. Classical Representation

- Start with a function-free first-order language
- Finitely many predicate symbols and constant symbols, but no function symbols
- Example: the DWR domain
- Locations: I1, I2, ...
- Containers: c1, c2, ...
- Piles: p1, p2, ...
- Robot carts: r1, r2, ...
- Cranes: k1, k2, ...



## Quick review of terminology

- Atom: predicate symbol and args
- Use these to represent both fixed and dynamic ("fluent") relations

| $\operatorname{adjacent}\left(l, l^{\prime}\right)$ | $\operatorname{attached}(p, l)$ |
| :--- | :--- |
| $\operatorname{occupied}(\Lambda)$ | at $(r, l)$ |
| belong $(k, l)$ |  |
| loaded $(r, c)$ | unloaded $(r)$ |
| holding $(k, c)$ | empty $(k)$ |
| in $(c, p)$ | on $\left(c, c^{\prime}\right)$ |
| top $(c, p)$ | top(pallet, $p)$ |

- Ground expression: contains no variable symbols - e.g., in(c1,p3)
- Unground expression: at least one variable symbol - e.g., in(c1,x)
- Substitution: $\theta=\left\{x_{1} \leftarrow t_{1}, x_{2} \leftarrow t_{2}, \ldots, x_{n} \leftarrow t_{n}\right\}$
- Each $x_{i}$ is a variable symbol; each $t_{i}$ is a term
- Instance of $e$ : result of applying a substitution $\theta$ to $e$
- Replace variables of e simultaneously, not sequentially


## States

- State: a set $s$ of ground atoms
- The atoms represent the things that are true in one of $\Sigma$ 's states
- Only finitely many ground atoms, so only finitely many possible states

\{attached(p1,loc1), in(c1,p1), in(c3,p1), top ( $\mathrm{c} 3, \mathrm{p} 1$ ), on( $\mathrm{c} 3, \mathrm{c} 1$ ), on( c 1, pallet), attached( $\mathrm{p} 2, \mathrm{loc} 1$ ), in( $\mathrm{c} 2, \mathrm{p} 2)$, top( $\mathrm{c} 2, \mathrm{p} 2$ ), on(c2,pallet), belong(crane1,loc1), empty(crane1), adjacent(loc1,loc2), adjacent(loc2,loc1), at(r1,loc2), occupied(loc2), unloaded(r1)\}.


## Operators

- Operator: a triple $0=($ name $(o)$, precond(o), effects(o))
- name(o) is a syntactic expression of the form $n\left(x_{1}, \ldots, x_{k}\right)$
- $n$ : operator symbol - must be unique for each operator
- $x_{1}, \ldots, x_{k}$ : variable symbols (parameters)
- must include every variable symbol in o
- precond(o): preconditions
- literals that must be true in order to use the operator
- effects(o): effects
- literals the operator will make true
$\operatorname{take}(k, l, c, d, p)$
;; crane $k$ at location $l$ takes $c$ off of $d$ in pile $p$
precond: belong $(k, l)$, attached $(p, l)$, empty $(k), \operatorname{top}(c, p)$, on $(c, d)$
effects: $\quad$ holding $(k, c), \neg \operatorname{empty}(k), \neg \operatorname{in}(c, p), \neg \operatorname{top}(c, p), \neg \operatorname{on}(c, d), \operatorname{top}(d, p)$


## Actions

- Action: ground instance (via substitution) of an operator

take $(k, l, c, d, p)$
;; crane $k$ at location $l$ takes $c$ off of $d$ in pile $p$
precond: belong $(k, l)$, attached $(p, l)$, empty $(k)$, top $(c, p)$, on $(c, d)$
effects: $\quad \operatorname{holding}(k, c), \neg \operatorname{empty}(k), \neg \operatorname{in}(c, p), \neg \operatorname{top}(c, p), \neg \operatorname{on}(c, d), \operatorname{top}(d, p)$
take(crane1,loc1,c3,c1,p1)
;; crane crane1 at location loc1 takes c3 off c1 in pile p1 precond: belong(crane1,loc1), attached(p1,loc1),
empty(crane1), top(c3,p1), on(c3,c1)
effects: holding(crane1,c3), ᄀempty(crane1), ᄀin(c3,p1),
ᄀtop(c3,p1), ᄀon(c3,c1), top(c1,p1)


## Notation

- Let a be an operator or action. Then
- precond ${ }^{+}(a)=$ \{atoms that appear positively in a's preconditions $\}$
- precond-(a) = \{atoms that appear negatively in a's preconditions\}
- effects ${ }^{+}(a)=\{$ atoms that appear positively in a's effects $\}$
- effects $^{-}(a)=$ \{atoms that appear negatively in a's effects\}


## E.g.,

take $(k, l, c, d, p)$
;; crane $k$ at location $l$ takes $c$ off of $d$ in pile $p$
precond: belong $(k, l)$, attached $(p, l)$, empty $(k), \operatorname{top}(c, p)$, on $(c, d)$
effects: $\quad \operatorname{holding}(k, c), \neg \operatorname{empty}(k), \neg \operatorname{in}(c, p), \neg \operatorname{top}(c, p), \neg \operatorname{on}(c, d), \operatorname{top}(d, p)$

- effects ${ }^{+}$(take $(k, l, c, d, p)=\{$ holding $(k, c)$, top $(d, p)\}$
- effects- ${ }^{-(\operatorname{take}(k, l, c, d, p)}=\{\operatorname{empty}(k)$, in $(c, p), \operatorname{top}(c, p)$, on $(c, d)\}$


## Aside: Some things to note

- The state only explicitly represents what is true. The semantics of this representation is that any fluent not included in the state is false - just like a database. (Recall that one of the assumptions of classical planning is complete initial (and subsequent) state. The problem would be a lot harder w/o this assumption!!)
- Terminology: an action is a ground operator. In the Knowledge Representation (KR) literature the concept of an "operator" is not used. Actions may be ground or unground.
- Classical planners generally operate over ground actions.


## Applicability

- An action $a$ is applicable to a state $s$ if $s$ satisfies precond(a),
- i.e., if precond ${ }^{+}(a) \subseteq s$ and precond-(a) $\cap s=\varnothing$
- Here are an action and a state that it's applicable to:

take(crane1,loc1,c3,c1,p1)
;; crane crane1 at location loc1 takes c3 off c1 in pile p1 precond: belong(crane1,loc1), attached(p1,loc1),
empty(crane1), top(c3,p1), on(c3,c1)
effects: holding(crane1,c3), ᄀempty(crane1), $\neg \mathrm{in}(\mathrm{c} 3, \mathrm{p} 1)$,
$\neg$ top (c3,p1), ᄀon(c3, c1), top(c1,p1)


## Result of Performing an Action

- If $a$ is applicable to $s$, the result of performing it is

$$
\gamma(s, a)=\left(s-\operatorname{effects}^{-}(\mathrm{a})\right) \cup \text { effects }^{+}(\mathrm{a})
$$

Set of things that are true.

- Delete negative effects, and add positive ones (if not in set then false)
take(crane1,loc1,c3,c1,p1)
;; crane crane1 at location loc1 takes c3 off c1 in pile p1 precond: belong(crane1,loc1), attached(p1,loc1),

$$
\text { empty(crane1), top(c3,p1), on }(c 3, c 1)
$$

effects: holding(crane1,c3), ᄀempty(crane1), $\neg \mathrm{in}(\mathrm{c} 3, \mathrm{p} 1)$,

$$
\neg \operatorname{top}(\mathrm{c} 3, \mathrm{p} 1), \neg \mathrm{on}(\mathrm{c} 3, \mathrm{c} 1), \text { top }(\mathrm{c} 1, \mathrm{p} 1)
$$



## $\operatorname{move}(r, l, m)$ <br> Operators for the DWR Domain

;; robot $r$ moves from location $l$ to location $m$
precond: $\operatorname{adjacent}(l, m)$, at $(r, l), \neg \operatorname{occupied}(m)$ effects: $\quad$ at $(r, m)$, occupied $(m), \neg \operatorname{occupied}(l), \neg$ at $(r, l)$

- Planning domain:
language \& operators
- Operators corresponds to a set of state-transition systems
;; crane $k$ at location $l$ loads container $c$ onto robot $r$ precond: belong $(k, l)$, holding $(k, c)$, at $(r, l)$, unloaded $(r)$ effects: $\quad$ empty $(k), \neg$ holding $(k, c)$, loaded $(r, c), \neg$ unloaded $(r)$
unload $(k, l, c, r)$
;; crane $k$ at location $l$ takes container $c$ from robot $r$ precond: belong $(k, l)$, at $(r, l)$, loaded $(r, c)$, empty $(k)$ effects: $\quad \neg$ empty $(k)$, holding $(k, c)$, unloaded $(r), \neg$ loaded

```
put(k,l,c,d,p)
```

;; crane $k$ at location $l$ puts $c$ onto $d$ in pile $p$
 precond: belong $(k, l)$, attached $(p, l)$, holding $(k, c), \operatorname{top}(d, p)$
effects: $\quad \neg \operatorname{holding}(k, c)$, empty $(k), \operatorname{in}(c, p), \operatorname{top}(c, p)$, on $(c, d), \neg \operatorname{top}(d, p)$
$\operatorname{take}(k, l, c, d, p)$
;; crane $k$ at location $l$ takes $c$ off of $d$ in pile $p$
precond: belong $(k, l)$, attached $(p, l)$, empty $(k), \operatorname{top}(c, p)$, on $(c, d)$
effects: $\quad$ holding $(k, c), \neg \operatorname{empty}(k), \neg \operatorname{in}(c, p), \neg \operatorname{top}(c, p), \neg \operatorname{on}(c, d), \operatorname{top}(d, p)$

## Planning Problems

Given a planning domain (language $L$, operators $O$ )

- Encoding of a planning problem: a triple $P=\left(O, s_{0}, g\right)$
- $O$ is the collection of operators
- $s_{0}$ is a state (the initial state)
- $g$ is a set of literals (the goal formula)
- The actual planning problem: $\mathcal{P}=\left(\Sigma, s_{0}, g\right)$
- $s_{0}$ and $g$ are as above
- $\Sigma=(S, A, \gamma)$ is a state-transition system
- $S=\{$ all sets of ground atoms in $L\}$
- $A=\{$ all ground instances of operators in $O\}$
- $\gamma=$ state-transition function determined by the operators


## Plans and Solutions

- Plan*: any sequence of actions $\sigma=\left\langle a_{1}, a_{2}, \ldots, a_{n}\right\rangle$ such that each $a_{i}$ is a ground instance of an operator in $O$
- The plan is a solution for $P=\left(O, s_{0}, g\right)$ if it is executable and achieves $g$
- i.e., if there are states $s_{0}, s_{1}, \ldots, s_{n}$ such that
- $\gamma\left(s_{0}, a_{1}\right)=s_{1}$
- $\gamma\left(s_{1}, a_{2}\right)=s_{2}$
- $\gamma\left(s_{n-1}, a_{n}\right)=s_{n}$
- $s_{n}$ satisfies $g$
*Recall that we are restricting our attention to "Classical Planning"


## Example

- Let $P_{1}=\left(O, s_{1}, g_{1}\right)$, where
$g_{1}=\{\operatorname{loaded}(\mathrm{r} 1, \mathrm{c} 3)$,
- $O$ is the set of operators given earlier

loc1

loc2

$$
s_{1}=\{\operatorname{attached}(\mathrm{p} 1, \operatorname{loc} 1), \quad \text { in }(\mathrm{c} 1, \mathrm{p} 1), \quad \text { in }(\mathrm{c} 3, \mathrm{p} 1),
$$ top(c3,p1), on(c3,c1), on(c1,pallet), attached(p2,loc1), in(c2,p2), top(c2,p2), on(c2,pallet), belong(crane1,loc1), empty(crane1), adjacent(loc1,loc2), adjacent(loc2,loc1), at(r1,loc2), occupied(loc2), unloaded(r1)\}.

## Example GOAL STATE:

$g_{1}=\{\operatorname{loaded}(\mathrm{r} 1, \mathrm{c} 3), \mathrm{at}(\mathrm{r} 1, \mathrm{loc} 2)\}$

## INITIAL STATE:



The DWR state $s_{1}=\{\operatorname{attached}(\mathrm{p} 1, \mathrm{loc} 1)$, in( $\mathrm{c} 1, \mathrm{p} 1$ ), in( $\mathrm{c} 3, \mathrm{p} 1$ ), top(c3,p1), on(c3,c1), on(c1,pallet), attached(p2,loc1), in(c2,p2), top(c2,p2), on(c2,pallet), belong(crane1,loc1), empty(crane1), adjacent(loc1,loc2), adjacent(loc2,loc1), at(r1,loc2), occupied(loc2), unloaded(r1)\}.

## Example（cont．）


－Here are three solutions for $P_{1}$ ：
－〈take（crane1，loc1，c3，c1，p1），move（r1，loc2，loc1），move（r1，loc1，loc2）， move（r1，loc2，loc1），load（crane1，loc1，c3，r1），move（r1，loc1，loc2）$)$
－〈take（crane1，loc1，c3，c1，p1），move（r1，loc2，loc1）， load（crane1，loc1，c3，r1），move（r1，loc1，loc2）〉
－〈move（r1，loc2，loc1），take（crane1，loc1，c3，c1，p1）， load（crane1，loc1，c3，r1），move（r1，loc1，loc2）＞
－Each produces：


## Example（cont．）

－First is redundant：can remove actions and still have a solution
1．〈take（crane1，loc1，c3，c1，p1），move（r1，loc2，loc1），move（r1，loc1，loc2）， move（r1，loc2，loc1），load（crane1，loc1，c3，r1），move（r1，loc1，loc2）$)$

2．〈take（crane1，loc1，c3，c1，p1），move（r1，loc2，loc1）， load（crane1，loc1，c3，r1），move（r1，loc1，loc2）〉

3．〈move（r1，loc2，loc1），take（crane1，loc1，c3，c1，p1）， load（crane1，loc1，c3，r1），move（r1，loc1，loc2）〉
$\bullet 2^{\text {nd }}$ and $3^{\text {rd }}$ are irredundant and shortest


## 2. Set-Theoretic Representation

Like classical rep'n, but restricted to propositional logic.


- States:
- Instead of a collection of ground atoms ...
\{on(c1,pallet), on(c1,r1), on(c1,c2), ..., at(r1,|1), at(r1,|2), ...\}
... use a collection of propositions (boolean variables): \{on-c1-pallet, on-c1-r1, on-c1-c2, ..., at-r1-I1, at-r1-I2, ...\}

Instead of operators like this one,
take $(k, l, c, d, p)$
;; crane $k$ at location $l$ takes $c$ off of $d$ in pile $p$
precond: belong $(k, l)$, $\operatorname{attached}(p, l)$, empty $(k), \operatorname{top}(c, p)$, on $(c, d)$
effects: $\quad$ holding $(k, c), \neg \operatorname{empty}(k), \neg \operatorname{in}(c, p), \neg \operatorname{top}(c, p), \neg \operatorname{on}(c, d), \operatorname{top}(d, p)$
Take all of the operator instances, E.g.:
take(crane1,loc1,c3,c1,p1)
;; crane crane1 at location loc1 takes c3 off c1 in pile p1
precond: belong(crane1,loc1), attached(p1,loc1),
empty(crane1), top(c3,p1), on(c3,c1)
effects: holding(crane1,c3), ᄀempty(crane1), $\neg \mathrm{in}(\mathrm{c} 3, \mathrm{p} 1)$, $\neg$ top (c3,p1), ᄀon(c3,c1), top(c1,p1)

## And rewrite ground atoms as propositions, E.g.:

## take-crane1-loc1-c3-c1-p1

precond: belong-crane1-loc1, attached-p1-loc1,empty-crane1, top-c3-p1, on-c3-c1 delete: empty-crane1, in-c3-p1, top-c3-p1, on-c3-p1 add: holding-crane1-c3, top-c1-p1

## Comparison

A set-theoretic representation is equivalent to a classical representation in which all of the atoms are ground

## Problem: Exponential blowup

- If a classical operator contains $n$ atoms and each atom has arity $k$, then it corresponds to $c^{n k}$ actions where $c=|\{c o n s t a n t ~ s y m b o l s\}|$


## 3. State-Variable Representation (aka FDR)

- Non-fluents (properties that don't change) are ground relations:
e.g., adjacent(loc1,loc2)
- Fluents are functions:
i.e., for properties that can change, assign values to state variables
- Classical and state-variable rep'ns take similar amounts of space each can be translated into the other in low-order polynomial time

```
move(r,l,m)
    ; robot rat location l moves to an adjacent location m
    precond: rloc}(r)=l,\operatorname{adjacent}(l,m
    effects: rloc(r)\leftarrowm
```

```
{top(p1)=c3,
    cpos(c3)=c1,
    cpos(c1)=pallet,
    holding(crane1)=nil,
    rloc(r1)=loc2,
    loaded(r1)=nil, ...}
```



## State-Variable Representation (cont.)

- Captures further information about the state. E.g., that state variables can only take on one of the values in the domain. This helps reduce the search space.
- Basis for the SAS and SAS+ formalisms (used most recently in the FastDownward Planner (FD) and its descendents (e.g., LAMA, etc)
- Basis for encodings further plan properties such as domain transition graphs (DTGs) and causal graphs (CG)


## Example: The Blocks World (Review on your own)

## Example: The Blocks World

- Infinitely wide table, finite number of children's blocks
- Ignore where a block is located on the table
- A block can sit on the table or on another block
- Want to move blocks from one configuration to another
- e.g.,

- Classical, set-theoretic, and state-variable formulations for the case of FIVE BLOCKS follow.


## 1. Example Classical Representation

- Constant symbols:
- The blocks: a, b, c, d, e
- Predicates:

- ontable $(x)$ - block $x$ is on the table
- on $(x, y) \quad$ - block $x$ is on block $y$
- clear $(x)$ - block $x$ has nothing on it
- holding $(x)$ - the robot hand is holding block $x$
- handempty - the robot hand isn't holding anything


## Classical Operators

## unstack $(x, y)$



Effects: $\sim \operatorname{on}(x, y), \sim \operatorname{clear}(x), \sim$ handempty, holding $(x)$, clear $(y)$

## $\operatorname{stack}(x, y)$

Precond: holding $(x)$, clear $(y)$
Effects: ~holding $(x), \sim$ clear $(y)$, on $(x, y)$, clear $(x)$, handempty

## $\operatorname{pickup}(x)$

Precond: ontable $(x)$, clear $(x)$, handempty
Effects: ~ontable $(x), \sim \operatorname{clear}(x)$, $\sim$ handempty, holding $(x)$

## putdown $(x)$

Precond: holding $(x)$
Effects: ~holding $(x)$, ontable $(x)$, clear $(x)$, handempty

## 2. Example Set-Theoretic Representation

For five blocks, 36 propositions, 50 actions

## E.g.,

 ontable-a - block a is on the table on-c-a $\quad$ - block c is on block a clear-c - block $c$ has nothing on it holding-d - the robot hand is holding block d handempty - the robot hand isn't holding anything
... (31 more)

## Set-Theoretic Actions

## E.g.,



## 3. Example State-Variable Representation

- Constant symbols:
a, b, c, d, e of type block
0,1 , table, nil of type other
- State variables:
$\operatorname{pos}(x)=y \quad$ if block $x$ is on block $y$
 $\operatorname{pos}(x)=$ table if block $x$ is on the table $\operatorname{pos}(x)=$ nil if block $x$ is being held clear $(x)=1 \quad$ if block $x$ has nothing on it $\operatorname{clear}(x)=0 \quad$ if block $x$ is being held or has a block on it holding $=x \quad$ if the robot hand is holding block $x$ holding $=$ nil $\quad$ if the robot hand is holding nothing


## State-Variable Operators

```
unstack(x : block, y : block)
Precond: \(\operatorname{pos}(x)=y, \operatorname{clear}(y)=0, \operatorname{clear}(x)=1\), holding=nil Effects: \(\operatorname{pos}(x)=\operatorname{nil}\), clear \((x)=0\), holding \(=x\), clear \((y)=1\)
```


## $\operatorname{stack}(x:$ block, $y$ : block)

Precond: holding $=x, \operatorname{clear}(x)=0, \operatorname{clear}(y)=1$
Effects: holding=nil, clear $(y)=0, \operatorname{pos}(x)=y, \operatorname{clear}(x)=1$


## Representational Equivalence

- Any problem that can be represented in one representation can also be represented in the other two
- Can convert in linear time and space, except when converting to set-theoretic (where we get an exponential blowup)

${ }^{* * *}$ ) trivially, or there can be a more parsimonious problem-specific encoding that ignores irrelevant variables


## Comparison

- Classical representation
- Most popular for classical planning, basis of PDDL
- Set-theoretic representation
- Can take much more space than classical representation
- Useful in algorithms that manipulate ground atoms directly
- e.g., planning graphs, SAT
- Useful for certain kinds of theoretical studies
- State-variable representation (e.g., SAS, SAS+,"FDR")
- Equivalent to classical representation in expressive power
- Arguably less natural to conceive
- Clever problem-specific encodings can be much more compact and embed critical info (e.g., one-of constraints)
- Leveraged in many of the state-of-the-art heuristic search classical planners (e.g., FD, LAMA, etc)
- Useful in non-classical planning problems as a way to handle numbers, functions, time


## Extending Expressivity: ADL*

- Previous representations were so-called "STRIPS" rep'ns. These have useful properties for automatically generating classical plans, but are not always sufficient to express the behaviour of more complex domains.
- ADL is a richer, and thus more compact, representation language that allows for
- Disjunction and Quantification in preconditions and goals
- Effects that are Quantified, and/or Conditional (effect is conditioned on state)
- PDDL supports STRIPS and ADL, but not all planners support ADL, and not all planners even support a so-called Classical Representation
- In the KR community ADL or greater is common.
* ADL = "Action Description Language", [Pednault, KR89]


## Pros/Cons: Compiling to Canonical Action Rep'n

Possible to compile down ADL actions into STRIPS actions

- Quantification -> conjunctions/disjunctions over finite universes
- Actions with conditional effects -> multiple (exponentially more) actions without conditional effects
- Actions with disjunctive effects -> multiple actions, each of which take one of the disjuncts as their preconditions (called "determinization")
- Domain axioms (ramifications) -> the individual effects of the actions; so all actions satisfy STRIPS assumption
Compilation is not always a win-win.
- By compiling down to canonical form, we can concentrate on highly efficient planning for canonical actions
- However, often compilation leads to an exponential blowup and makes it harder to exploit the structure of the domain
- By leaving actions in non-canonical form, we can often do more compact encoding of the domains as well as more efficient search
- However, we will have to continually extend planning algorithms to handle these representations

