CSC2542 Representations for (Classical) Planning

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Acknowledgements

Some the slides used in this course are modifications of Dana Nau's lecture slides for the textbook Automated Planning, licensed under the Creative Commons Attribution-NonCommercial-ShareAlike License: http://creativecommons.org/licenses/by-nc-sa/2.0/

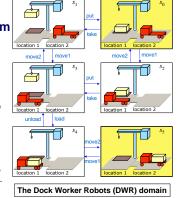
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Recall: Planning Problem

- $\boldsymbol{P} = (\Sigma, \boldsymbol{S}_0, \boldsymbol{G})$
- Σ: System Description
- $s_{o:}$ Initial state(s) E.g., Initial state = s_0
- G: Objective Goal state, Set of goal states, Set of tasks, "trajectory" of states, Objective function, ... E.g., Goal state = s₅



Further Recall: System Description (as a state transition system)

 $\boldsymbol{\Sigma} = (\boldsymbol{S}, \boldsymbol{A}, \boldsymbol{E}, \boldsymbol{\gamma})$

- S = {states}
- A = {actions}
- E = {exogenous events}
- State-transition function γ : S x (A \cup E) \rightarrow 2^S

Example: Dock Workers Robots from previous slide

• $S = \{s_0, ..., s_5\}$

- A = {move1, move2, put, take, load, unload}
- *E* = {}
- γ: as captured by the arrows mapping states and actions to successor states

Representational Challenge

 How do we represent our planning problem is a way that supports exploration of the principles and practice of automated planning?

Approach:

- There isn't one answer.
- The [GNT04] proposes representations that are suitable for generating classical plans.

[GNT04] = Ghallab, Nau, Traverso, Automated Planning: Theory and Practice, 2004

Broad Perspective on Plan Representation

The right representation for the right objective. Distinguish representation schemes for:

- 1. studying the principles of planning and related tasks.
- 2. specifying planning domains
- 3. direct use within (classical) planners

Summary: Broad Perspective

- 1. Studying the formal principles of planning and other related task
 - (First-order) logical languages (e.g., situation calculus, A languages, event calculus, fluent calculus, PDL) Properties:
 - well-defined semantics, representational issues must be addressed in the language (not in the algorithm that interprets and manipulates them)
 - excellent for study and proving properties. Not ideal for 3 below.
- 2. Specifying planning domains PDDL-n (PDDL2.1, PDDL2.2, PDDL3,)
 - Properties:
 - (reasonably) well-defined semantics
 - designed for input to planners translate to an internal representation for specific planners. Translators exist for most state-of-the-art planners

3. Direct use within (classical) planners • Classical representation (e.g., STRIPS)

- Set-theoretic representation (basis for rep'ns used w/ SAT solvers)
- State-variable representation (aka "Finite Domain Repn' (FDR)*")(e.g., SAS, SAS+) Variants of these exist for particular planners (e.g., SAT solvers, model

checkers, etc. * [Helmert, AIJ 2009]

This Lecture:

- 1. Studying the formal principles of planning and other related task ٠ (First-order) logical languages
- calculus PDL) Pr dressed in the es them) WILL COVER LATER elow 2. Spec PDDL-n (PDDL2.1, PDDL2.2, PDDL3, Properties:
 - · (reasonably) well-defined semantics
 - designed for input to planners translate to an internal representation for specific planners. Translators exist for most state-of-the-art planners
- 3. Direct use within (classical) planners (what's in the text) Classical representation (e.g., STRIPS)

 - State-variable representation (basis for rep'ns used w/ SAT solvers)
 State-variable representation (aka "FDR") (e.g., SAS, SAS+)
 Variants of these exist for particular planners (e.g., SAT solvers, model checkers. etc.)

Outline

- Representation schemes for classical planning •
 - 1. Classical representation
 - 2. Set-theoretic representation
 - State-variable representation
- Examples: DWR and the Blocks World
- Comparisons

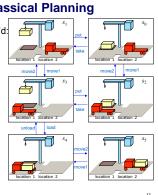
Quick Review of Classical Planning



- A0: Finite
- A1: Fully observable
- A2: Deterministic
- A3: Static

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- A4: Attainment goals
- A5: Sequential plans
- A6: Implicit time
- A7: Offline planning



Representation: Motivation for Approach Default view:

- · represent state explicitly
- represent actions as a transition system (e.g., as an incidence matrix) Problem:
 - · explicit graph corresponding to transition system is huge
 - · direct manipulation of transition system is cumbersome

Solution:

- Provide compact representation of transition system & induced graph 1. Explicate the structure of the "states"
 - · e.g., states specified in terms of state variables
- 2. Represent actions not as transition system/incidence matrices but as functions (e.g., operators) specified in terms of the state variables
 - An action is applicable to a state when some state variables have certain values. When applicable, it will change the values of certain (other) state variables
- 3. To plan,
 - Just give the initial state
 - · Use the operators to generate the other states as needed

Why is this more compact?

Why is this more compact than an explicit transition system?

- · In an explicit transition system, actions are represented as state-tostate transitions. Each action will be represented by an incidence matrix of size |S|x|S|
- In the proposed model, actions are represented only in terms of state variables whose values they care about, and whose value they affect. (It exploits the structure of the problem!)
- Consider a state space of 1024 states. It can be represented by log₂1024=10 state variables. If an action needs variable v1 to be true and makes v7 to be false, it can be represented by just 2 bits (instead of a 1024x1024 matrix)
 - · Of course, if the action has a complicated mapping from states to states, in the worst case the action rep will be just as large
 - The assumption being made here is that the actions will have effects on a small number of state variables.

1. Classical Representation

- Start with a function-free first-order language
 - Finitely many predicate symbols and constant symbols, but *no* function symbols
- Example: the DWR domain
 - Locations: I1, I2, ...
 - Containers: c1, c2, ..
 - Piles: p1, p2, ...
 - Robot carts: r1, r2, ...
 - Cranes: k1, k2, ...

Quick review of terminology

- Atom: predicate symbol and args
 - Use these to represent both fixed and dynamic ("fluent") relations adjacent(*l*,*l*) attached(*p*,*l*) belong(*k*,*l*)
 - occupied(*I*) loaded(*r*,*c*)
 - $\begin{array}{ll} \text{holding}(k,c) & \text{empty}(k) \\ \text{in}(c,p) & \text{on}(c,c') \end{array}$
 - top(c,p) top(pallet,p)
- Ground expression: contains no variable symbols e.g., in(c1,p3)
- Unground expression: at least one variable symbol e.g., in(c1,x)
- Substitution: θ = {x₁ ← t₁, x₂ ← t₂, ..., x_n ← t_n}
 Each x_i is a variable symbol; each t_i is a term
 - **Instance** of *e*: result of applying a substitution θ to *e*

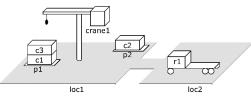
at(*r,l*)

unloaded(r)

• Replace variables of e simultaneously, not sequentially

States

- State: a set s of ground atoms
 - The atoms represent the things that are **true** in one of Σ 's states
 - Only finitely many ground atoms, so only finitely many possible states



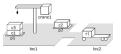
 $\label{eq:constraint} \begin{array}{l} \{\texttt{attached}(p1,loc1), \ \texttt{in}(c1,p1), \ \texttt{in}(c3,p1), \\ \texttt{top}(c3,p1), \ \texttt{on}(c3,c1), \ \texttt{on}(c1,pallet), \ \texttt{attached}(p2,loc1), \ \texttt{in}(c2,p2), \ \texttt{top}(c2,p2), \\ \texttt{on}(c2,pallet), \ \texttt{belong}(\texttt{crane1},loc1), \ \texttt{empty}(\texttt{crane1}), \ \texttt{adjacent}(loc1,loc2), \ \texttt{adjacent}(loc1,loc2), \ \texttt{adjacent}(loc2,loc1), \ \texttt{attached}(p1,loc1), \ \texttt{adjacent}(loc1,loc2), \ \texttt{adjacent}($

Operators

- **Operator:** a triple *o*=(name(*o*), precond(*o*), effects(*o*))
 - name(o) is a syntactic expression of the form n(x₁,...,x_k)
 n: operator symbol must be unique for each operator
 - x₁,...,x_k: variable symbols (parameters)
 must include every variable symbol in o
 - precond(o): preconditions
 - literals that must be true in order to use the operator
 - effects(o): effects
 - literals the operator will make true
- $\mathsf{take}(k,l,c,d,p)$

;; crane k at location l takes c off of d in pile p precond: belong(k, l), attached(p, l), empty(k), top(c, p), on(c, d) effects: holding(k, c), \neg empty(k), \neg in(c, p), \neg top(c, p), \neg on(c, d), top(d, p)

Actions



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 Action: ground instance (via substitution) of an operator

 $\mathsf{take}(k,l,c,d,p)$

;; crane k at location l takes c off of d in pile p precond: belong(k,l), attached(p,l), empty(k), top(c, p), on(c, d)

```
effects: holding(k, c), \neg empty(k), \neg in(c, p), \neg top(c, p), \neg on(c, d), top(d, p)
```

take(crane1,loc1,c3,c1,p1)

;; crane cranel at location loc1 takes c3 off c1 in pile p1 precond: belong(crane1,loc1), attached(p1,loc1), empty(crane1), top(c3,p1), on(c3,c1) effects: holding(crane1,c3), -empty(crane1), -in(c3,p1), -top(c3,p1), -on(c3,c1), top(c1,p1)

Notation

- Let a be an operator or action. Then
 - precond*(a) = {atoms that appear positively in a's preconditions}
 - precond⁻(a) = {atoms that appear negatively in a's preconditions}
 - effects+(*a*) = {atoms that appear positively in *a*'s effects}
 - effects⁻(a) = {atoms that appear negatively in a's effects}

E.g.,

 $\mathsf{take}(k, l, c, d, p)$

;; crane k at location l takes c off of d in pile p precond: belong(k, l), attached(p, l), empty(k), top(c, p), on(c, d)effects: holding(k, c), $\neg empty(k), \neg in(c, p), \neg top(c, p), \neg on(c, d), top(d, p)$

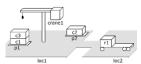
- effects+(take(k,l,c,d,p) = {holding(k,c), top(d,p)}
- effects⁻(take(k,l,c,d,p) = {empty(k), in(c,p), top(c,p), on(c,d)}

Aside: Some things to note

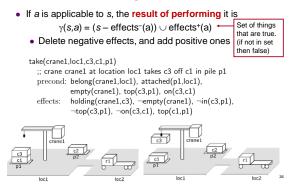
- The state only explicitly represents what is **true**. The semantics of this representation is that any fluent not included in the state is **false** just like a database. (Recall that one of the assumptions of classical planning is complete initial (and subsequent) state. The problem would be a lot harder w/o this assumption!!)
- **Terminology:** an action is a ground operator. In the Knowledge Representation (KR) literature the concept of an "operator" is not used. Actions may be ground or unground.
- Classical planners generally operate over ground actions.

Applicability

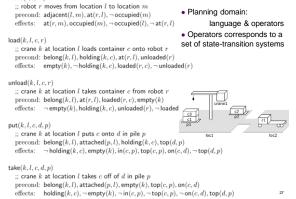
- An action a is applicable to a state s if s satisfies precond(a),
 - i.e., if precond⁺(a) \subseteq s and precond⁻(a) \cap s = \emptyset
- Here are an action and a state that it's applicable to:



Result of Performing an Action



move(r, l, m) Operators for the DWR Domain



Planning Problems

Given a planning domain (language *L*, operators *O*)

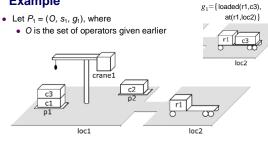
- *Encoding* of a planning problem: a triple *P*=(*O*,*s*₀,*g*)
 - O is the collection of operators
 - s₀ is a state (the initial state)
 - g is a set of literals (the goal formula)
- The *actual planning problem*: $\mathcal{P} = (\Sigma, s_0, g)$
 - s_0 and g are as above
 - $\Sigma = (S, A, \gamma)$ is a state-transition system
 - S = {all sets of ground atoms in L}
 - A = {all ground instances of operators in O}
 - γ = state-transition function determined by the operators

Plans and Solutions

- Plan*: any sequence of actions σ = (a₁, a₂, ..., a_n) such that each a_i is a ground instance of an operator in O
- The plan is a *solution* for *P*=(*O*,*s*₀,*g*) if it is executable and achieves *g*
 - i.e., if there are states s_0, s_1, \ldots, s_n such that
 - $\gamma(s_0, a_1) = s_1$
 - $\gamma(s_1, a_2) = s_2$
 - •...
 - ...,
 - $\gamma(s_{n-1},a_n) = s_n$
 - s_n satisfies g

* Recall that we are restricting our attention to "Classical Planning"

Example

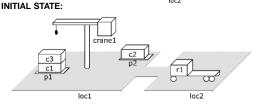


 $s_1 = \{\texttt{attached(p1,loc1)}, \texttt{ in(c1,p1)}, \texttt{ in(c3,p1)}, \texttt{ top(c3,p1)}, \texttt{ on(c3,c1)}, \texttt{ on(c1,pallet)}, \texttt{ attached(p2,loc1)}, \texttt{ in(c2,p2)}, \texttt{ top(c2,p2)}, \texttt{ on(c2,pallet)}, \texttt{ belong(crane1,loc1)}, \texttt{ empty(crane1)}, \texttt{ adjacent(loc1,loc2)}, \texttt{ adjacent(loc1,l$ cent(loc2,loc1), at(r1,loc2), occupied(loc2), unloaded(r1)}.

Example

GOAL STATE: $g_1 = \{ loaded(r1,c3), at(r1, loc2) \}$



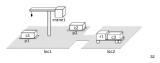


The DWR state $s_1 = \{ \mathsf{attached}(\mathsf{p1}, \mathsf{loc1}), \mathsf{in}(\mathsf{c1}, \mathsf{p1}), \mathsf{in}(\mathsf{c3}, \mathsf{p1}),$ top(c3,p1), on(c3,c1), on(c1,pallet), attached(p2,loc1), in(c2,p2), top(c2,p2), on(c2,pallet), belong(crane1,loc1), empty(crane1), adjacent(loc1,loc2), adjacent(loc2,loc1), at(r1,loc2), occupied(loc2), unloaded(r1)}.

Example (cont.)

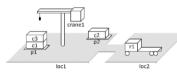


- Here are three solutions for P1:
 - $\bullet \quad \ \ \langle take(crane1,loc1,c3,c1,p1), \ move(r1,loc2,loc1), \ move(r1,loc1,loc2), \\$ $move(r1,loc2,loc1), \ load(crane1,loc1,c3,r1), \ move(r1,loc1,loc2)\rangle$
 - (take(crane1,loc1,c3,c1,p1), move(r1,loc2,loc1), ٠ load(crane1,loc1,c3,r1), move(r1,loc1,loc2)>
 - (move(r1,loc2,loc1), take(crane1,loc1,c3,c1,p1), load(crane1,loc1,c3,r1), move(r1,loc1,loc2))
- · Each produces:



2. Set-Theoretic Representation

Like classical rep'n, but restricted to propositional logic.



States:

- Instead of a collection of ground atoms ... {on(c1,pallet), on(c1,r1), on(c1,c2), ..., at(r1,l1), at(r1,l2), ...}
- ... use a collection of propositions (boolean variables): {on-c1-pallet, on-c1-r1, on-c1-c2, ..., at-r1-l1, at-r1-l2, ...}

move(r1,loc2,loc1), load(crane1,loc1,c3,r1), move(r1,loc1,loc2))

1.

Example (cont.)

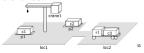
 $\label{eq:crane1,loc1,c3,c1,p1} $$ (take(crane1,loc1,c3,c1,p1), move(r1,loc2,loc1), $$$ 2. load(crane1,loc1,c3,r1), move(r1,loc1,loc2))

• First is redundant: can remove actions and still have a solution

(take(crane1,loc1,c3,c1,p1), move(r1,loc2,loc1), move(r1,loc1,loc2),

load(crane1,loc1,c3,r1), move(r1,loc1,loc2))

2nd and 3rd are irredundant and shortest



Instead of operators like this one,

	a, p) : at location l takes c off of d in pile p belong (k, l) , attached (p, l) , empty (k) , top (c, p) , on((c, d)
	$holding(k,c), \neg empty(k), \neg in(c,p), \neg top(c,p), \neg cop(c,p), \neg cop(c,$	
ake all of t	he operator instances, E.g.:	
take(crane)	.,loc1,c3,c1,p1)	
;; crane	crane1 at location loc1 takes c3 off c1 in pile p1	
precond	<pre>belong(crane1,loc1), attached(p1,loc1), empty(crane1), top(c3,p1), on(c3,c1)</pre>	
effects:	$\label{eq:cranel} \begin{array}{l} \mbox{holding(cranel,c3), } \neg\mbox{empty(cranel), } \neg\mbox{in(c3,p1), } \\ \neg\mbox{top(c3,p1), } \neg\mbox{on(c3,c1), top(c1,p1)} \end{array}$	
nd rewrite	ground atoms as propositions, E.g.:	

precond: belong-crane1-loc1, attached-p1-loc1, empty-crane1, top-c3-p1, on-c3-c1 empty-crane1, in-c3-p1, top-c3-p1, on-c3-p1 delete: add: holding-crane1-c3, top-c1-p1

Comparison

A set-theoretic representation is equivalent to a classical representation in which all of the atoms are ground

Problem: Exponential blowup

 If a classical operator contains n atoms and each atom has arity k, then it corresponds to c^{nk} actions where $c = |\{\text{constant symbols}\}|$

3. State-Variable Representation (aka FDR)

- Non-fluents (properties that don't change) are ground relations: e.g., adjacent(loc1,loc2)
- · Fluents are functions:

{top

- i.e., for properties that can change, assign values to state variables
- Classical and state-variable rep'ns take similar amounts of space
- each can be translated into the other in low-order polynomial time

```
move(r, l, m)
  ;; robot r at location l moves to an adjacent location m
   precond: rloc(r) = l, adjacent(l, m)
   effects: rloc(r) \leftarrow m
```

top(p1)=c3, cpos(c3)=c1, cpos(c1)=pallet, holding(crane1)=nil, rloc(r1)=loc2, loaded(r1)=nil,}	crane1 ci pl	n m
	loc1	loc2

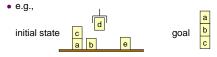
State-Variable Representation (cont.)

- Captures further information about the state. E.g., that state variables can only take on one of the values in the domain. This helps reduce the search space.
- · Basis for the SAS and SAS+ formalisms (used most recently in the FastDownward Planner (FD) and its descendents (e.g., LAMA, etc)
- Basis for encodings further plan properties such as domain transition graphs (DTGs) and causal graphs (CG)

Example: The Blocks World (Review on your own)

Example: The Blocks World

- Infinitely wide table, finite number of children's blocks
- Ignore where a block is located on the table
- A block can sit on the table or on another block
- · Want to move blocks from one configuration to another



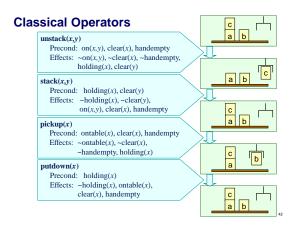
 Classical, set-theoretic, and state-variable formulations for the case of FIVE BLOCKS follow.

1. Example Classical Representation

- Constant symbols:
- The blocks: a, b, c, d, e

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- Predicates:
 - ontable(x) block x is on the table
 - on(*x,y*) - block x is on block y
 - clear(x) - block x has nothing on it
 - holding(x) the robot hand is holding block x
 - handempty- the robot hand isn't holding anything



2. Example Set-Theoretic Representation

For five blocks, 36 propositions, 50 actions

E.

	a b e
.g.,	
ontable-a	 block a is on the table
on-c-a	- block c is on block a
clear-c	 block c has nothing on it
holding-d	- the robot hand is holding block d

holding-d - the robot hand is holding block d
 handempty - the robot hand isn't holding anything
 ... (31 more)

Set-Theoretic Actions С а b unstack-c-a E.g., Pre: on-c,a, clear-c, handempty Del: on-c,a, clear-c, handempty holding-c, clear-a Add: С a b stack-c-a Pre: holding-c, clear-a holding-c, clear-a Del: Add: on-c-a, clear-c, handempty c a b pickup-c Pre: ontable-c, clear-c, handempty Del: ontable-c, clear-c, handempty Add: holding-c С Ь а putdown-c Pre: holding-c holding-c Del: Add: ontable-c, clear-c, handempty c a b ... (46 more)

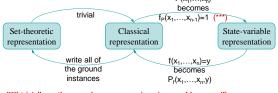
3. Example State-Variable Representation

•	Constant symbols:		
	a, b, c, d, e	of type block	
	0, 1, table, nil		
٠	State variables:		
	pos(x) = y	if block x is on block y	
	pos(x) = table	if block x is on the table	
	pos(x) = nil	if block x is being held	
	clear(x) = 1	if block x has nothing on it	
	clear(x) = 0	if block x is being held or has a block on it	
	holding = x	if the robot hand is holding block x	
	holding = nil	if the robot hand is holding nothing	

State-Variable Operators С Г а b unstack(x : block, y : block) Precond: pos(x)=y, clear(y)=0, clear(x)=1, holding=nil Effects: pos(x)=nil, clear(x)=0, holding=x, clear(y)=1 с a b stack(x : block, y : block) Precond: holding=x, clear(x)=0, clear(y)=1 Effects: holding=nil, clear(y)=0, pos(x)=y, clear(x)=1 С а b pickup(x : block) Precond: pos(x)=table, clear(x)=1, holding=nil Effects: pos(x)=nil, clear(x)=0, holding=xС b а putdown(x : block)Precond: holding=x Effects: holding=nil, pos(x)=table, clear(x)=1 с а b

Representational Equivalence

- Any problem that can be represented in one representation can also be represented in the other two
- Can convert in linear time and space, except when converting to set-theoretic (where we get an exponential blowup)
 P(x₁,...,x_n)



(***) trivially, or there can be a more parsimonious problem-specific encoding that ignores irrelevant variables

Comparison

- Classical representation
 - Most popular for classical planning, basis of PDDL
- Set-theoretic representation
 - Can take much more space than classical representation
 - Useful in algorithms that manipulate ground atoms directly
 e.g., planning graphs, SAT
 - Useful for certain kinds of theoretical studies
- State-variable representation (e.g., SAS, SAS+, "FDR")
 - · Equivalent to classical representation in expressive power
 - · Arguably less natural to conceive
 - Clever problem-specific encodings can be much more compact and embed critical info (e.g., one-of constraints)
 - Leveraged in many of the state-of-the-art heuristic search classical planners (e.g., FD, LAMA, etc)
 - Useful in non-classical planning problems as a way to handle numbers, functions, time

Extending Expressivity: ADL*

- Previous representations were so-called "STRIPS" rep'ns. These have useful properties for automatically generating classical plans, but are not always sufficient to express the behaviour of more complex domains.
- ADL is a richer, and thus more compact, representation language that allows for
 - Disjunction and Quantification in preconditions and goals
 - Effects that are Quantified, and/or Conditional (effect is conditioned on state)
- PDDL supports STRIPS and ADL, but not all planners support ADL, and not all planners even support a so-called Classical Representation
- In the KR community ADL or greater is common.

* ADL = "Action Description Language", [Pednault, KR89]

Pros/Cons: Compiling to Canonical Action Rep'n

Possible to compile down ADL actions into STRIPS actions

- Quantification -> conjunctions/disjunctions over finite universes
- Actions with conditional effects -> multiple (exponentially more) actions without conditional effects
- Actions with disjunctive effects -> multiple actions, each of which take one of the disjuncts as their preconditions (called "determinization")
- Domain axioms (ramifications) -> the individual effects of the actions; so all actions satisfy STRIPS assumption
- Compilation is not always a win-win.
 - By compiling down to canonical form, we can concentrate on highly efficient planning for canonical actions
 - However, often compilation leads to an exponential blowup and makes it harder to exploit the structure of the domain
 - By leaving actions in non-canonical form, we can often do more compact encoding of the domains as well as more efficient search
 - However, we will have to continually extend planning algorithms to handle these representations