CSC2542 Representations for (Classical) Planning

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Acknowledgements

Some the slides used in this course are modifications of Dana Nau's lecture slides for the textbook *Automated Planning*, licensed under the Creative Commons Attribution-NonCommercial-ShareAlike License:

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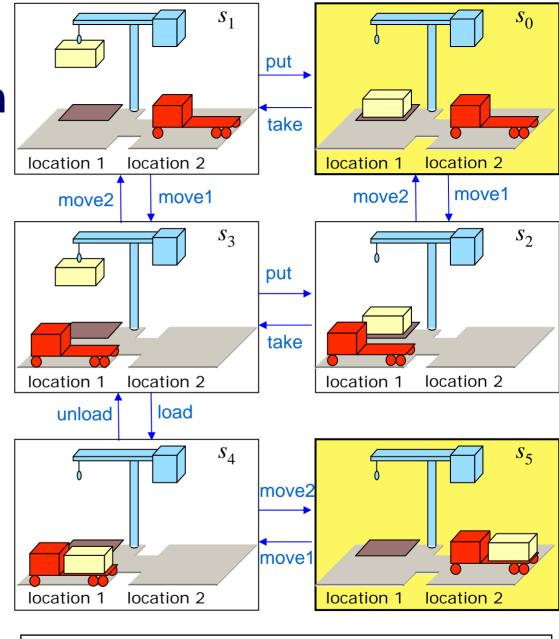
Recall: Planning Problem

$$P = (\Sigma, s_0, G)$$

Σ: System Description

 s_{o} : Initial state(s) E.g., Initial state = s_0

G: Objective
Goal state,
Set of goal states,
Set of tasks,
"trajectory" of states,
Objective function, ...
E.g., Goal state = s₅



The Dock Worker Robots (DWR) domain

Further Recall: System Description (as a state transition system)

$$\Sigma = (S, A, E, \gamma)$$

- *S* = {states}
- *A* = {actions}
- *E* = {exogenous events}
- State-transition function $\gamma: S \times (A \cup E) \rightarrow 2^S$

Example: Dock Workers Robots from previous slide

- $S = \{s_0, ..., s_5\}$
- A = {move1, move2, put, take, load, unload}
- *E* = {}
- γ: as captured by the arrows mapping states and actions to successor states

Representational Challenge

 How do we represent our planning problem is a way that supports exploration of the principles and practice of automated planning?

Approach:

- There isn't one answer.
- The textbook proposes representations that are suitable for generating classical plans.

Broad Perspective on Plan Representation

The right representation for the right objective.

Distinguish representation schemes for:

- studying the principles of planning and related tasks.
- 2. specifying planning domains
- 3. direct use within (classical) planners

Summary: Broad Perspective

1. Studying the formal principles of planning and other related task

- (First-order) logical languages
 (e.g., situation calculus, A languages, event calculus, fluent calculus, PDL)
 Properties:
- well-defined semantics, representational issues must be addressed in the language (not in the algorithm that interprets and manipulates them)
- excellent for study and proving properties. Not ideal for 3 below.

2. Specifying planning domains

PDDL-n (PDDL2.1, PDDL2.2, PDDL3,)

Properties:

- (reasonably) well-defined semantics
- designed for input to planners translate to an internal representation for specific planners. Translators exist for most state-of-the-art planners

3. Direct use within (classical) planners

- Classical representation (e.g., STRIPS)
- Set-theoretic representation (basis for rep'ns used w/ SAT solvers)
- State-variable representation (e.g., SAS, SAS+)

Variants of these exist for particular planners (e.g., SAT solvers, model checkers, etc.)

This Lecture:

- 1. Studying the formal principles of planning and other related task
 - (First-order) logical languages

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calculus, PDL)

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2. Spec

Pro

PDDL-n (PDDL2.1, PDDL2.2, PDDL3,)

Properties:

- (reasonably) well-defined semantics
- designed for input to planners translate to an internal representation for specific planners. Translators exist for most state-of-the-art planners
- 3. Direct use within (classical) planners (what's in the text)
 - Classical representation (e.g., STRIPS)
 - Set-theoretic representation (basis for rep'ns used w/ SAT solvers)
 - State-variable representation (e.g., SAS, SAS+)

Variants of these exist for particular planners (e.g., SAT solvers, model checkers, etc.)

Outline

- Representation schemes for classical planning
 - Classical representation
 - 2. Set-theoretic representation
 - 3. State-variable representation
- Examples: DWR and the Blocks World
- Comparisons

Quick Review of Classical Planning

8 restrictive assumptions req'd:

A0: Finite

A1: Fully observable

A2: Deterministic

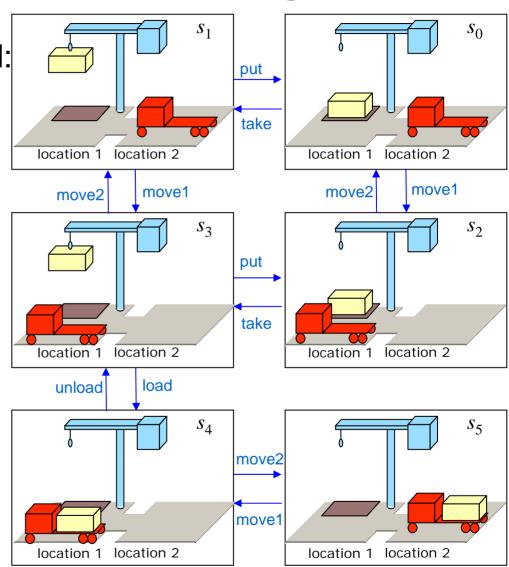
A3: Static

A4: Attainment goals

A5: Sequential plans

A6: Implicit time

A7: Offline planning



Representation: Motivation for Approach

Default view:

- represent state explicitly
- represent actions as a transition system (e.g., as an incidence matrix)

Problem:

- explicit graph corresponding to transition system is huge
- direct manipulation of transition system is cumbersome

Solution:

Provide compact representation of transition system & induced graph

- Explicate the structure of the "states"
 - e.g., states specified in terms of state variables
- 2. Represent actions not as transition system/incidence matrices but as functions (e.g., operators) specified in terms of the state variables
 - An action is applicable to a state when some state variables have certain values. When applicable, it will change the values of certain (other) state variables
- 3. To plan,
 - Just give the initial state
 - Use the operators to generate the other states as needed

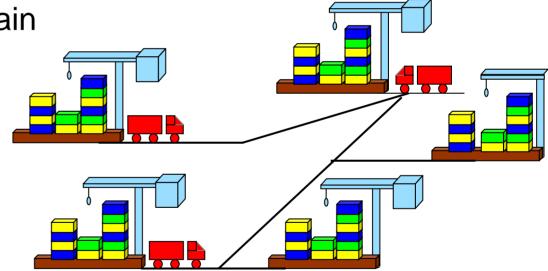
Why is this more compact?

Why is this more compact than an explicit transition system?

- In an explicit transition system, actions are represented as state-tostate transitions. Each action will be represented by an incidence matrix of size |S|x|S|
- In the proposed model, actions are represented only in terms of state variables whose values they care about, and whose value they affect. (It exploits the structure of the problem!)
- Consider a state space of 1024 states. It can be represented by log₂1024=10 state variables. If an action needs variable v1 to be true and makes v7 to be false, it can be represented by just 2 bits (instead of a 1024x1024 matrix)
 - Of course, if the action has a complicated mapping from states to states, in the worst case the action rep will be just as large
 - The assumption being made here is that the actions will have effects on a small number of state variables.

1. Classical Representation

- Start with a function-free first-order language
 - Finitely many predicate symbols and constant symbols, but no function symbols
- Example: the DWR domain
 - Locations: I1, I2, ...
 - Containers: c1, c2, ...
 - Piles: p1, p2, ...
 - Robot carts: r1, r2, ...
 - Cranes: k1, k2, ...



Quick review of terminology

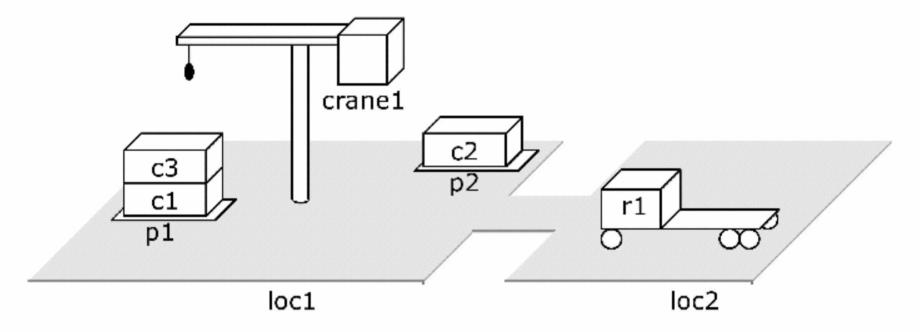
- Atom: predicate symbol and args
 - Use these to represent both fixed and dynamic ("fluent") relations

```
\begin{array}{lll} \operatorname{adjacent}(\textit{I},\textit{I}') & \operatorname{attached}(\textit{p},\textit{I}) & \operatorname{belong}(\textit{k},\textit{I}) \\ \operatorname{occupied}(\textit{I}) & \operatorname{at}(\textit{r},\textit{I}) \\ \operatorname{loaded}(\textit{r},\textit{c}) & \operatorname{unloaded}(\textit{r}) \\ \operatorname{holding}(\textit{k},\textit{c}) & \operatorname{empty}(\textit{k}) \\ \operatorname{in}(\textit{c},\textit{p}) & \operatorname{on}(\textit{c},\textit{c}') \\ \operatorname{top}(\textit{c},\textit{p}) & \operatorname{top}(\text{pallet},\textit{p}) \end{array}
```

- Ground expression: contains no variable symbols e.g., in(c1,p3)
- Unground expression: at least one variable symbol e.g., in(c1,x)
- Substitution: $\theta = \{x_1 \leftarrow v_1, x_2 \leftarrow v_2, ..., x_n \leftarrow v_n\}$
 - Each x_i is a variable symbol; each v_i is a term
- Instance of e: result of applying a substitution θ to e
 - Replace variables of e simultaneously, not sequentially

States

- State: a set s of ground atoms
 - The atoms represent the things that are **true** in one of Σ 's states
 - Only finitely many ground atoms, so only finitely many possible states



 $\{ attached(p1,loc1), \ in(c1,p1), \ in(c3,p1), \\ top(c3,p1), \ on(c3,c1), \ on(c1,pallet), \ attached(p2,loc1), \ in(c2,p2), \ top(c2,p2), \\ on(c2,pallet), \ belong(crane1,loc1), \ empty(crane1), \ adjacent(loc1,loc2), \ adjacent(loc2,loc1), \ at(r1,loc2), \ occupied(loc2), \ unloaded(r1) \}.$

Operators

- Operator: a triple o=(name(o), precond(o), effects(o))
 - name(o) is a syntactic expression of the form $n(x_1,...,x_k)$
 - n: operator symbol must be unique for each operator
 - $x_1, ..., x_k$: variable symbols (parameters)
 - must include every variable symbol in o
 - precond(o): preconditions
 - literals that must be true in order to use the operator
 - effects(o): effects
 - literals the operator will make true

```
take(k,l,c,d,p)

;; crane k at location l takes c off of d in pile p

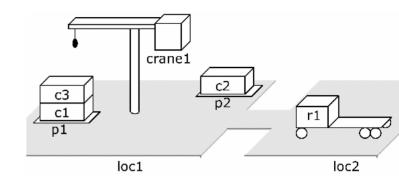
precond: belong(k,l), attached(p,l), empty(k), top(c,p), on(c,d)

effects: holding(k,c), \neg empty(k), \neg in(c,p), \neg top(c,p), \neg on(c,d), top(d,p)
```

Actions

take(k, l, c, d, p)

Action: ground instance (via substitution) of an operator



```
;; crane k at location l takes c off of d in pile p
 precond: belong(k, l), attached(p, l), empty(k), top(c, p), on(c, d)
             \mathsf{holding}(k,c), \neg \mathsf{empty}(k), \neg \mathsf{in}(c,p), \neg \mathsf{top}(c,p), \neg \mathsf{on}(c,d), \mathsf{top}(d,p)
 effects:
take(crane1,loc1,c3,c1,p1)
   ;; crane crane1 at location loc1 takes c3 off c1 in pile p1
   precond: belong(crane1,loc1), attached(p1,loc1),
              empty(crane1), top(c3,p1), on(c3,c1)
              holding(crane1,c3), \negempty(crane1), \negin(c3,p1),
   effects:
               \neg top(c3,p1), \neg on(c3,c1), top(c1,p1)
```

Notation

- Let a be an operator or action. Then
 - precond*(a) = {atoms that appear positively in a's preconditions}
 - precond⁻(a) = {atoms that appear negatively in a's preconditions}
 - effects+(a) = {atoms that appear positively in a's effects}
 - effects⁻(a) = {atoms that appear negatively in a's effects}

E.g.,

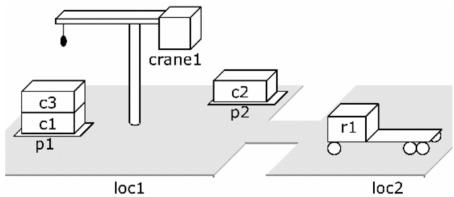
- effects⁺(take(k,l,c,d,p) = {holding(k,c), top(d,p)}
- effects⁻(take(k,l,c,d,p) = {empty(k), in(c,p), top(c,p), on(c,d)}

Aside: Some things to note

- The state only explicitly represents what is true. The semantics of this representation is that any fluent not included in the state is false just like a database.
 (Recall that one of the assumptions of classical planning is complete initial (and subsequent) state.
 The problem would be a lot harder w/o this assumption!!)
- **Terminology:** an action is a ground operator. In the Knowledge Representation (KR) literature the concept of an "operator" is not used. Actions may be ground or unground.
- Classical planners generally operate over ground actions.

Applicability

- An action a is applicable to a state s if s satisfies precond(a),
 - i.e., if precond⁺(a) \subseteq s and precond⁻(a) \cap s = \emptyset
- Here are an action and a state that it's applicable to:



```
take(crane1,loc1,c3,c1,p1)

;; crane crane1 at location loc1 takes c3 off c1 in pile p1

precond: belong(crane1,loc1), attached(p1,loc1),

empty(crane1), top(c3,p1), on(c3,c1)

effects: holding(crane1,c3), ¬empty(crane1), ¬in(c3,p1),

¬top(c3,p1), ¬on(c3,c1), top(c1,p1)
```

Result of Performing an Action

- If a is applicable to s, the result of performing it is
 γ(s,a) = (s − effects⁻(a)) ∪ effects⁺(a)
 - Delete negative effects, and add positive ones

Set of things that are true. (if not in set then false)

```
take(crane1,loc1,c3,c1,p1) 

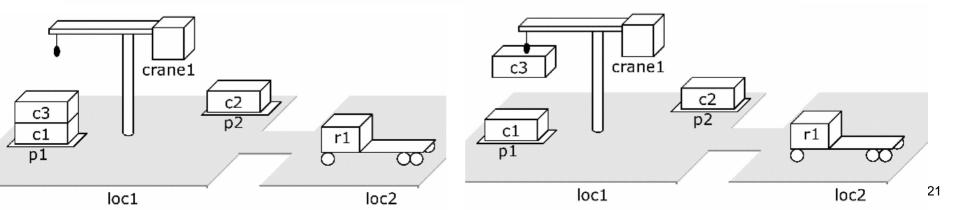
;; crane crane1 at location loc1 takes c3 off c1 in pile p1 

precond: belong(crane1,loc1), attached(p1,loc1), 

empty(crane1), top(c3,p1), on(c3,c1) 

effects: holding(crane1,c3), \negempty(crane1), \negin(c3,p1), 

\negtop(c3,p1), \negon(c3,c1), top(c1,p1)
```



Operators for the DWR Domain

;; robot r moves from location l to location m precond: $\operatorname{adjacent}(l,m),\operatorname{at}(r,l),\neg\operatorname{occupied}(m)$ effects: $\operatorname{at}(r,m),\operatorname{occupied}(m),\neg\operatorname{occupied}(l),\neg\operatorname{at}(r,l)$

move(r, l, m)

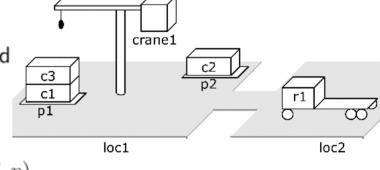
effects:

- effects: $\mathsf{at}(r,m), \mathsf{occupied}(m), \neg \, \mathsf{occupied}(l), \neg \, \mathsf{at}(r, m)$
- $\mathsf{load}(k,l,c,r)$;; crane k at location l loads container c onto robot r
 - precond: belong(k, l), holding(k, c), at(r, l), unloaded(r)
 - effects: empty(k), \neg holding(k, c), loaded(r, c), \neg unloaded(r)
- unload(k, l, c, r);; crane k at location l takes container c from robot r
 - precond: $\mathsf{belong}(k, l), \mathsf{at}(r, l), \mathsf{loaded}(r, c), \mathsf{empty}(k)$ effects: $\neg \, \mathsf{empty}(k), \mathsf{holding}(k, c), \mathsf{unloaded}(r), \neg \, \mathsf{loaded}$
- $\mathsf{put}(k,l,c,d,p)$
 - ;; crane k at location l puts c onto d in pile p
 - precond: belong(k, l), attached(p, l), holding(k, c), top(d, p) effects: \neg holding(k, c), empty(k), in(c, p), top(c, p), on(c, d), \neg top(d, p)

 $\mathsf{holding}(k,c), \neg \mathsf{empty}(k), \neg \mathsf{in}(c,p), \neg \mathsf{top}(c,p), \neg \mathsf{on}(c,d), \mathsf{top}(d,p)$

- take(k, l, c, d, p)
 - ;; crane k at location l takes c off of d in pile p precond: belong (k, l), attached (p, l), empty (k), top (c, p), on (c, d)

- Planning domain:
- Ianguage & operators
 Operators corresponds to a set of state-transition systems



Planning Problems

Given a planning domain (language L, operators O)

- Encoding of a planning problem: a triple $P=(O,s_0,g)$
 - O is the collection of operators
 - s_0 is a state (the initial state)
 - g is a set of literals (the goal formula)
- The actual planning problem: $\mathcal{P} = (\Sigma, s_0, g)$
 - s_0 and g are as above
 - $\Sigma = (S, A, \gamma)$ is a state-transition system
 - S = {all sets of ground atoms in L}
 - A = {all ground instances of operators in O}
 - γ = state-transition function determined by the operators

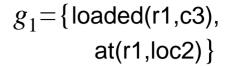
Plans and Solutions

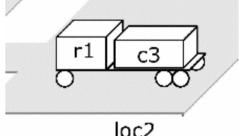
- *Plan*: any sequence of actions $\sigma = \langle a_1, a_2, ..., a_n \rangle$ such that each a_i is a ground instance of an operator in O
- The plan is a solution for P=(O,s₀,g) if it is executable and achieves g
 - i.e., if there are states $s_0, s_1, ..., s_n$ such that
 - $\gamma(s_0, a_1) = s_1$
 - $\gamma(s_1, a_2) = s_2$
 - . . .

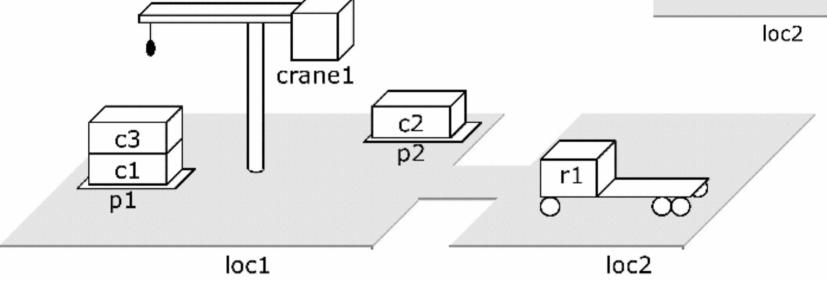
 - s_n satisfies g

Example

- Let $P_1 = (O, s_1, g_1)$, where
 - O is the set of operators given earlier





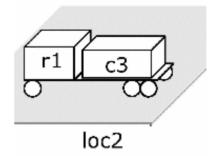


 $s_1=\{attached(p1,loc1), in(c1,p1), in(c3,p1), top(c3,p1), on(c3,c1), on(c1,pallet), attached(p2,loc1), in(c2,p2), top(c2,p2), on(c2,pallet), belong(crane1,loc1), empty(crane1), adjacent(loc1,loc2), adjacent(loc2,loc1), at(r1,loc2), occupied(loc2), unloaded(r1)\}.$

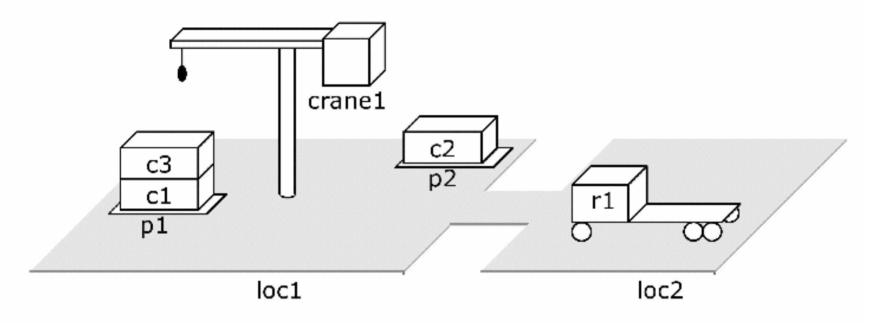
Example

GOAL STATE:

 $g_1 = \{ loaded(r1,c3), at(r1,loc2) \}$

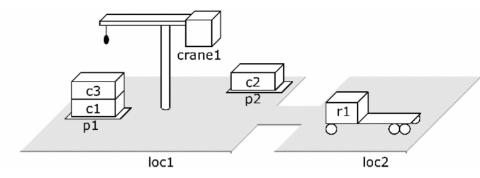


INITIAL STATE:

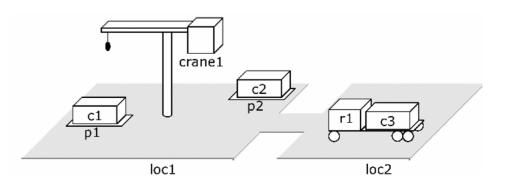


The DWR state s_1 ={attached(p1,loc1), in(c1,p1), in(c3,p1), top(c3,p1), on(c3,c1), on(c1,pallet), attached(p2,loc1), in(c2,p2), top(c2,p2), on(c2,pallet), belong(crane1,loc1), empty(crane1), adjacent(loc1,loc2), adjacent(loc2,loc1), at(r1,loc2), occupied(loc2), unloaded(r1)}.

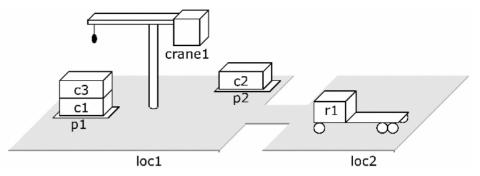
Example (cont.)



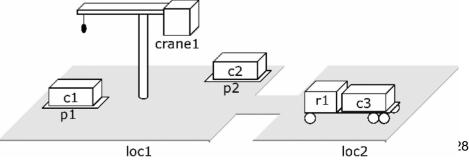
- Here are three solutions for P_1 :
 - (take(crane1,loc1,c3,c1,p1), move(r1,loc2,loc1), move(r1,loc1,loc2), move(r1,loc2,loc1), load(crane1,loc1,c3,r1), move(r1,loc1,loc2))
 - (take(crane1,loc1,c3,c1,p1), move(r1,loc2,loc1), load(crane1,loc1,c3,r1), move(r1,loc1,loc2))
 - (move(r1,loc2,loc1), take(crane1,loc1,c3,c1,p1), load(crane1,loc1,c3,r1), move(r1,loc1,loc2))
- Each produces:



Example (cont.)

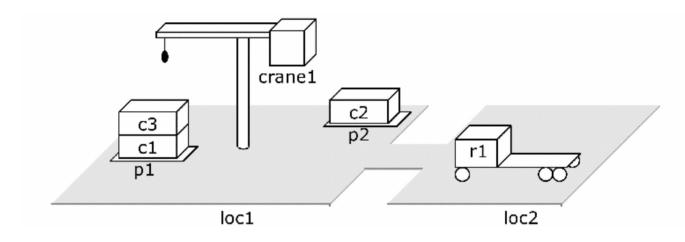


- First is redundant: can remove actions and still have a solution
 - 1. \(\take(\text{crane1,loc1,c3,c1,p1})\), \(\text{move}(\text{r1,loc2,loc1})\), \(\text{move}(\text{r1,loc1,loc2})\), move(r1,loc2,loc1), load(crane1,loc1,c3,r1), move(r1,loc1,loc2)
 - 2. \(\rangle \take(\text{crane1,loc1,c3,c1,p1})\), \(\text{move}(\text{r1,loc2,loc1})\), load(crane1,loc1,c3,r1), move(r1,loc1,loc2)>
 - 3. $\langle move(r1,loc2,loc1), take(crane1,loc1,c3,c1,p1), \rangle$ load(crane1,loc1,c3,r1), move(r1,loc1,loc2)
- •2nd and 3rd are *irredundant* and shortest



2. Set-Theoretic Representation

Like classical rep'n, but restricted to propositional logic.



- States:
 - Instead of a collection of ground atoms ...
 {on(c1,pallet), on(c1,r1), on(c1,c2), ..., at(r1,l1), at(r1,l2), ...}
 - ... use a collection of propositions (boolean variables): {on-c1-pallet, on-c1-r1, on-c1-c2, ..., at-r1-l1, at-r1-l2, ...}

Instead of operators like this one,

```
 \begin{array}{l} \mathsf{take}(k,l,c,d,p) \\ \mathsf{;;} \ \mathsf{crane} \ k \ \mathsf{at} \ \mathsf{location} \ l \ \mathsf{takes} \ c \ \mathsf{off} \ \mathsf{of} \ d \ \mathsf{in} \ \mathsf{pile} \ p \\ \mathsf{precond:} \ \mathsf{belong}(k,l), \mathsf{attached}(p,l), \mathsf{empty}(k), \mathsf{top}(c,p), \mathsf{on}(c,d) \\ \mathsf{effects:} \ \ \mathsf{holding}(k,c), \neg \, \mathsf{empty}(k), \neg \, \mathsf{in}(c,p), \neg \, \mathsf{top}(c,p), \neg \, \mathsf{on}(c,d), \mathsf{top}(d,p) \\ \end{array}
```

Take all of the operator instances, E.g.:

```
take(crane1,loc1,c3,c1,p1) 

;; crane crane1 at location loc1 takes c3 off c1 in pile p1 

precond: belong(crane1,loc1), attached(p1,loc1), 

empty(crane1), top(c3,p1), on(c3,c1) 

effects: holding(crane1,c3), \negempty(crane1), \negin(c3,p1), 

\negtop(c3,p1), \negon(c3,c1), top(c1,p1)
```

And rewrite ground atoms as **propositions**, E.g.:

```
take-crane1-loc1-c3-c1-p1
precond: belong-crane1-loc1, attached-p1-loc1, empty-crane1, top-c3-p1, on-c3-c1
delete: empty-crane1, in-c3-p1, top-c3-p1, on-c3-p1
add: holding-crane1-c3, top-c1-p1
```

Comparison

A set-theoretic representation is equivalent to a classical representation in which all of the atoms are ground

Problem: Exponential blowup

• If a classical operator contains n atoms and each atom has arity k, then it corresponds to c^{nk} actions where $c = |\{\text{constant symbols}\}|$

3. State-Variable Representation

Non-fluents (properties that don't change) are ground relations:
 e.g., adjacent(loc1,loc2)

- Fluents are functions:
 - i.e., for properties that can change, assign values to state variables
- Classical and state-variable rep'ns take similar amounts of space each can be translated into the other in low-order polynomial time

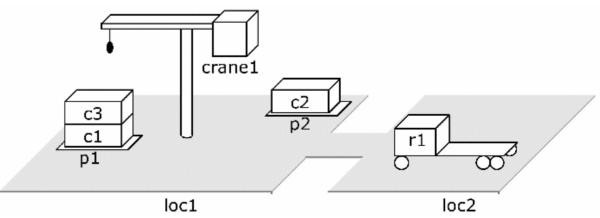
```
move(r, l, m)

;; robot r at location l moves to an adjacent location m

precond: rloc(r) = l, adjacent(l, m)

effects: rloc(r) \leftarrow m
```

```
{top(p1)=c3,
  cpos(c3)=c1,
  cpos(c1)=pallet,
  holding(crane1)=nil,
  rloc(r1)=loc2,
  loaded(r1)=nil, ...}
```



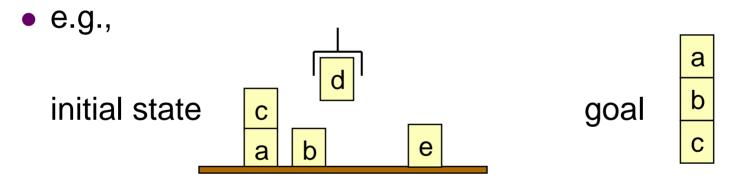
State-Variable Representation (cont.)

- Captures further information about the state. E.g., that state variables can only take on one of the values in the domain. This helps reduce the search space.
- Basis for the SAS and SAS+ formalisms (used most recently in the FastDownward Planner (FD)
- Basis for encodings further plan properties such as domain transition graphs (DTGs) and causal graphs (CG)

Example: The Blocks World (Review on your own)

Example: The Blocks World

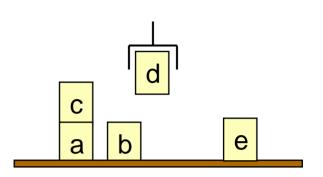
- Infinitely wide table, finite number of children's blocks
- Ignore where a block is located on the table
- A block can sit on the table or on another block
- Want to move blocks from one configuration to another



 Classical, set-theoretic, and state-variable formulations for the case of FIVE BLOCKS follow.

1. Example Classical Representation

- Constant symbols:
 - The blocks: a, b, c, d, e
- Predicates:
 - ontable(x) block x is on the table
 - on(x,y) block x is on block y
 - clear(x) block x has nothing on it
 - holding(x) the robot hand is holding block x
 - handempty the robot hand isn't holding anything



Classical Operators

unstack(x,y)

Precond: on(x,y), clear(x), handempty

Effects: \sim on(x,y), \sim clear(x), \sim handempty,

holding(x), clear(y)

stack(x,y)

Precond: holding(x), clear(y)

Effects: \sim holding(x), \sim clear(y),

on(x,y), clear(x), handempty

pickup(x)

Precond: ontable(x), clear(x), handempty

Effects: \sim ontable(x), \sim clear(x),

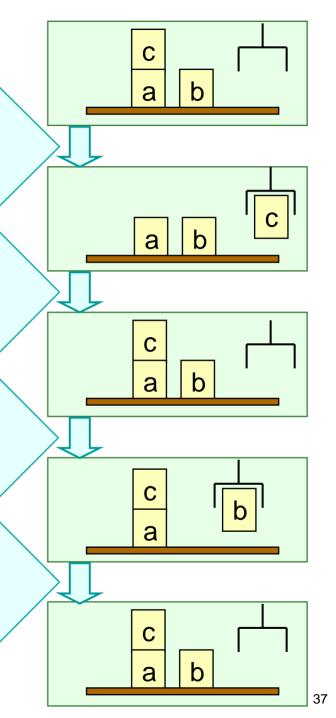
~handempty, holding(*x*)

putdown(x)

Precond: holding(x)

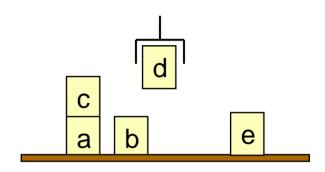
Effects: \sim holding(x), ontable(x),

clear(x), handempty



2. Example Set-Theoretic Representation

For five blocks, 36 propositions, 50 actions



E.g.,

ontable-a - block a is on the table

on-c-a - block c is on block a

clear-c - block c has nothing on it

holding-d - the robot hand is holding block d

handempty - the robot hand isn't holding anything

... (31 more)

Set-Theoretic Actions

E.g.,

unstack-c-a

Pre: on-c,a, clear-c, handempty

Del: on-c,a, clear-c, handempty

Add: holding-c, clear-a

stack-c-a

Pre: holding-c, clear-a

Del: holding-c, clear-a

Add: on-c-a, clear-c, handempty

pickup-c

Pre: ontable-c, clear-c, handempty

Del: ontable-c, clear-c, handempty

Add: holding-c

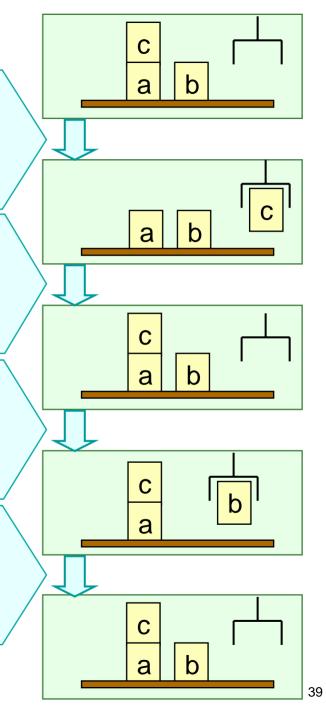
putdown-c

Pre: holding-c

Del: holding-c

Add: ontable-c, clear-c, handempty

... (46 more)



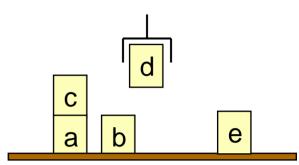
3. Example State-Variable Representation

Constant symbols:

- a, b, c, d, e of type block
- 0, 1, table, nil of type other

State variables:

- pos(x) = y if block x is on block y
- pos(x) = table if block x is on the table
- pos(x) = nil if block x is being held
- clear(x) = 1 if block x has nothing on it
- clear(x) = 0 if block x is being held or has a block on it
- holding = x if the robot hand is holding block x
- holding = nil if the robot hand is holding nothing



State-Variable Operators

unstack(x : block, y : block)

Precond: pos(x)=y, clear(y)=0, clear(x)=1, holding=nil

Effects: pos(x)=nil, clear(x)=0, holding=x, clear(y)=1

stack(x : block, y : block)

Precond: holding=x, clear(x)=0, clear(y)=1

holding=nil, clear(y)=0, pos(x)=y, clear(x)=1

pickup(x : block)

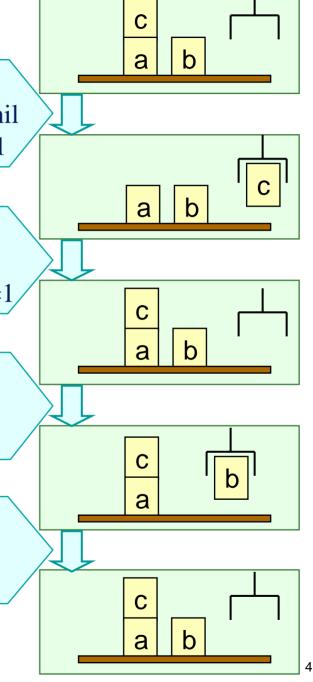
Precond: pos(x)=table, clear(x)=1, holding=nil

Effects: pos(x)=nil, clear(x)=0, holding=x

putdown(x : block)

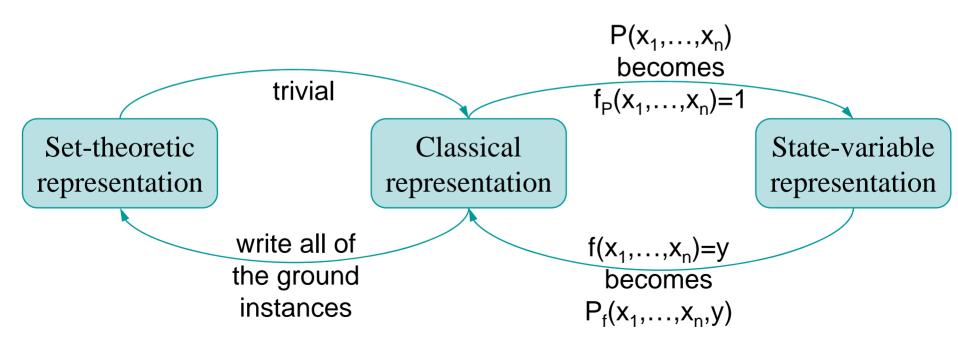
Precond: holding=x

Effects: holding=nil, pos(x)=table, clear(x)=1



Representational Equivalence

- Any problem that can be represented in one representation can also be represented in the other two
- Can convert in linear time and space, except when converting to set-theoretic (where we get an exponential blowup)



Comparison

- Classical representation
 - Most popular for classical planning, basis of PDDL
- Set-theoretic representation
 - Can take much more space than classical representation
 - Useful in algorithms that manipulate ground atoms directly
 - e.g., planning graphs, SAT
 - Useful for certain kinds of theoretical studies
- State-variable representation (e.g., SAS, SAS+)
 - Equivalent to classical representation in expressive power
 - Arguably less natural to conceive
 - Useful in non-classical planning problems as a way to handle numbers, functions, time

Extending Expressivity: ADL

- Previous representations were so-called "STRIPS" rep'ns.
 These have useful properties for automatically generating
 classical plans, but are not always sufficient to express the
 behaviour of more complex domains.
- ADL is a richer, and thus more compact, representation language that allows for
 - Disjunction and Quantification in preconditions and goals
 - Effects that are Quantified, and/or Conditional (effect is conditioned on state)
- PDDL supports STRIPS and ADL, but not all planners support ADL, and not all planners even support a so-called Classical Representation
- In the KR community ADL or greater is common.

Pros/Cons: Compiling to Canonical Action Rep'n

Possible to compile down ADL actions into STRIPS actions

- Quantification -> conjunctions/disjunctions over finite universes
- Actions with conditional effects -> multiple (exponentially more) actions without conditional effects
- Actions with disjunctive effects -> multiple actions, each of which take one of the disjuncts as their preconditions
- Domain axioms (ramifications) -> the individual effects of the actions; so all actions satisfy STRIPS assumption

Compilation is not always a win-win.

- By compiling down to canonical form, we can concentrate on highly efficient planning for canonical actions
 - However, often compilation leads to an exponential blowup and makes it harder to exploit the structure of the domain
- By leaving actions in non-canonical form, we can often do more compact encoding of the domains as well as more efficient search
 - However, we will have to continually extend planning algorithms to handle these representations