

Question 1. [10 MARKS]**Part (a)** [2 MARKS]

`list[s...m]` and `list[m+1...e]` are sorted in non-descending order and `list != null` and list contents are mutually Comparable.

Part (b) [5 MARKS]

```
private static void merge(Comparable[] list, int s, int m, int e) {  
  
    _F_ // index of current candidate in list[s...m]  
    _G_ // index of current candidate in list[m+1...e]  
  
    _K_ // Temporary storage to accumulate the merged result  
    int p=0; // index to put next element into merged  
  
    // merged[0...p] contains list[s...p1-1] merged with list[m+1...p2-1]  
    _J_  
    _E_  
    _H_  
    _B_  
} else {  
    _I_  
    _C_  
    _M_  
    _A_  
}  
_D_  
System.arraycopy(list, p1, list, s+p, m+1-p1);  
_N_  
_L_
```

Part (c) [3 MARKS]

Answer:

```
+---+---+---+---+---+---+---+---+  
| 0 | 1 | 4 | 5 | 8 | 2 | 3 | 6 | 7 | 9 |  
+---+---+---+---+---+---+---+---+
```

Question 2. [10 MARKS]

```
public class ArrayBST {  
  
    /** REPRESENTATION INVARIANT:  
  
        The left and right children of keys[i] are in keys[2i+1] and keys[2i+2], and for any node  
        the children in its left subtree are less than that node and the children in the right subtree are  
        greater.  
  
     */  
    private Comparable[] keys;  
  
    /** An empty ArrayBST that will always represent a tree of height <= maxHeight.  
     * Requires: maxHeight >= 0. */  
    public ArrayBST(int maxHeight) {  
  
        keys= new Comparable[((int) Math.pow(2, maxHeight)) - 1];  
  
        /** Return the root of a BST made of BSTNodes that represents the  
         * same tree as this tree does. */  
        public BSTNode linkedRepresentation() {  
  
            return linkedHelper(0);  
        }  
  
        private BSTNode linkedHelper(int i) {  
            if (i >= keys.length) {  
                return null;  
            }  
  
            if (keys[i] == null) {  
                return null;  
            }  
  
            BSTNode temp= new BSTNode();  
            temp.left= linkedHelper(i * 2 + 1);  
            temp.right= linkedHelper(i * 2 + 2);  
            temp.key= keys[i];  
  
            return temp;  
        }  
    }  
}
```

Question 3. [10 MARKS]**Rewritten comment:**

```
/**  
 * Return the first position of name in this list if name appears  
 * within the range beginning one position after start and ending at  
 * end (inclusive); return -1 otherwise.  
 * Requires: -1 <= start && end < size() && name != null  
 * @param name the name to search for  
 * @param start the position in OrderedList after which search starts  
 * @param end last position in OrderedList at which search stops  
 * @return position of name in OrderedList.  
 */
```

Question 4. [10 MARKS]**Part (a)** [3 MARKS]

```
private static void checkDigits(String s) throws WrongDigitException {
    for (int i= 0; i != s.length(); i++) {
        if (s.charAt(i) > '7') {
            throw new WrongDigitException("'" + s.charAt(i));
        }
    }
}
```

Part (b) [3 MARKS]

```
public class WrongDigitException extends Exception {
    public WrongDigitException(String m) {
        super(m);
    }
}
```

Part (c) [4 MARKS]

```
public static String getInput(BufferedReader br) throws IOException {
    String result= br.readLine();
    try {
        checkDigits(result);
    } catch (WrongDigitException e) {
        result= e.getMessage() + " is not a proper digit in base 8";
    }

    return result;
}
```

Question 5. [10 MARKS]**Part (a)** [4 MARKS]

Solution: The second, third and fourth statements should be circled.

Part (b) [6 MARKS]

Consider the following theorem:

2^n is always less than $n!$ for integer values of n greater than or equal to 4.

1. Fill in the blank with the smallest correct value.
2. Use induction to prove this theorem.

Let $S(n) = 2^n < n!$.

Prove: $S(n) \forall n \geq 4$.

Base case: Prove $S(4)$

When $n = 4$, $2^n = 16 < 24 = n!$. Thus, $S(n)$ holds for the base case.

Let $k \geq 4$ be an arbitrary integer.

Induction Hypothesis: Assume $S(k - 1)$: $2^{k-1} < (k - 1)!$.

Induction Step: Prove $S(k)$: $2^k < k!$.

$2^k = 2 * 2^{k-1}$. From the induction hypothesis, $2^{k-1} < (k - 1)!$. So $2^k < 2 * (k - 1)!$.

Because $k \geq 4$, $k \geq 2$. So $2^k < 2 * (k - 1)! < k * (k - 1)! = k!$.

Conclusion: $S(n)$ for all $n \geq 4$.

Question 6. [10 MARKS]**Part (a)** [6 MARKS]

What is the running time for the following operations, in $O(\dots)$ notation? Give as small and simple a bound as possible.

1. Finding out whether a linked list with n nodes contains any duplicate elements: $O(n^2)$
2. Printing the n th element of a sorted array: $O(1)$
3. Calculating the height of a binary tree with n nodes: $O(n)$
4. Printing the contents of a binary tree with n nodes in ascending order, without using any kind of list or another tree: $O(n^2)$
5. Printing the contents of a binary search tree with n nodes in ascending order: $O(n)$
6. Finding the smallest element in a full binary search tree with n nodes: $O(\log n)$

Part (b) [4 MARKS]

A method used to find all prime numbers that are less than n runs as follows:

1. Make a table of integers from 2 to n .
2. For every value of i from 2 to \sqrt{n} , cross out the entries $2 * i, 3 * i, 4 * i$ up to n .
3. Once $i > \sqrt{n}$, print out all entries that have not been crossed out.

Circle the tightest bound for the questions below.

A) What is the running time for Step 1?

$O(n)$

B) What is the running time for Step 2?

$O(n * \sqrt{n})$

C) What is the running time for Step 3?

$O(n)$

D) What is the total running time for this method? $O(n * \sqrt{n})$

Question 7. [10 MARKS]

Consider this code.

```

public class R {
    private R next;

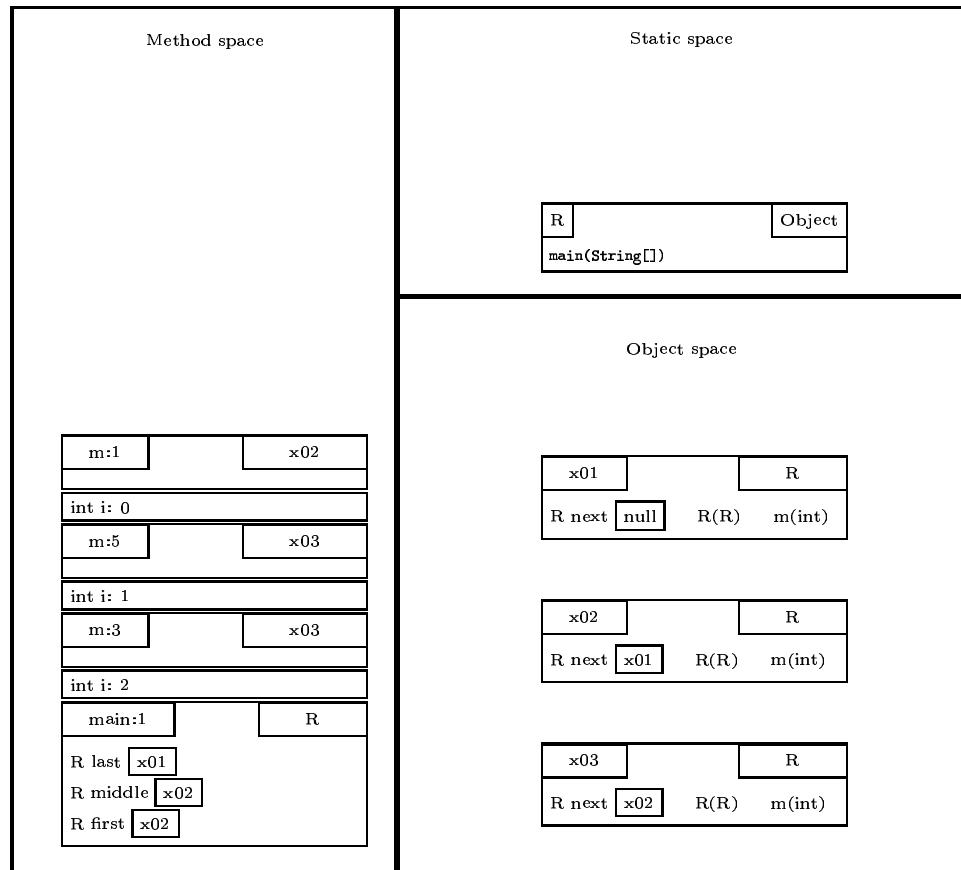
    public R(R n) {
        next= n;           // 1
    }

    public void m(int i) {
        if (i != 0) {      // 1
            if (i % 2 == 0) { // 2
                m(i - 1);   // 3
            } else {          // 4
                next.m(i - 1); // 5
            }
        }
    }

    public static void main(String[] args) {
        R last= new R(null); // 1
        R middle= new R(last); // 2
        R first= new R(middle); // 3
        first.m(2);           // 4
    }
}

```

Below is a memory model trace paused just after `main` has been called. Show the state of memory when line 1 of method `m` has been reached and parameter `i` is 0.



Question 8. [10 MARKS]**Part (a)** [5 MARKS]

```
/** Reverse the order of the nodes in the linked list <code>list</code> and return the list.
 * @param list first node of a linked list, null if the list is empty.
 * @return the first node of the modified list. */
public static Node reverse(Node list) {

    Node previous= null;

    while (list != null) {
        Node temp= list.next;
        list.next= previous;
        previous= list;
        list= temp;
    }

    return previous;
}
```

Part (b) [5 MARKS]

```
/** Reverse the order of the nodes in the linked list <code>list</code> and return the list.
 * @param list first node of a linked list, null if the list is empty.
 * @return the first node of the modified list. */
public static DoubleNode reverse(DoubleNode list) {
```

Here are two solutions:

```
DoubleNode previous= null;           DoubleNode t1= list;

while (list != null) {               while (t1 != null){
    list.previous= list.next;       DoubleNode t2= temp.next;
    list.next= previous;          t1.next= temp.previous;
    previous= list;              t1.previous= t2;
    list= list.previous;         list= t1;
}                                     temp= t1.previous;
return previous;                   }

                                         }
```

return list;

Total Marks = 80