



**Part I (10 marks)**

(a) Let  $L$  be a language. Circle one response for each statement below.

If  $L$  is decidable, then  $\bar{L}$  is decidable.

 TRUE

 FALSE

 DON'T KNOW

If  $L$  is recognizable, then  $\bar{L}$  is recognizable.

 TRUE

 FALSE

 DON'T KNOW

If  $L$  is recognizable but not decidable, then  $\bar{L}$  is not recognizable

 TRUE

 FALSE

 DON'T KNOW

If  $L$  is not recognizable, then  $\bar{L}$  is not recognizable.

 TRUE

 FALSE

 DON'T KNOW

If  $L$  is decidable and  $L' \subseteq L$  then  $L'$  is decidable

 TRUE

 FALSE

 DON'T KNOW

(b) For each language below, circle the **strongest** true statement.

$L_1 = \{\langle M, w \rangle \mid M \text{ accepts } w\}$

 DECIDABLE

 RECOGNIZABLE

 UNRECOGNIZABLE

 DON'T KNOW

$L_2 = \{\langle M, w, t \rangle \mid M \text{ accepts } w \text{ within } t \text{ steps}\}$

 DECIDABLE

 RECOGNIZABLE

 UNRECOGNIZABLE

 DON'T KNOW

$L_3 = \{\langle M, w \rangle \mid M \text{ does not accept } w\}$

 DECIDABLE

 RECOGNIZABLE

 UNRECOGNIZABLE

 DON'T KNOW

Student number:

(c) Write down the formal definition of the statement  $A \leq_m B$ .

**Solution:**

There exists a computable function  $f : \Sigma^* \mapsto \Sigma^*$  such that for all  $x \in \Sigma^*$ ,

$$x \in A \iff f(x) \in B$$

**Part II (10 marks)**

Prove that the following language is recognizable.

$$L_4 = \{\langle M_1, M_2 \rangle \mid L(M_1) \cap L(M_2) \neq \emptyset\}$$

Note that your answer will be marked on its structure as well as its content.

**Solution:**

Define  $M =$  “on input  $\langle M_1, M_2 \rangle \dots$

1. for  $t \leftarrow 1, 2, 3, \dots, \infty$  do
2.     for  $i \leftarrow 1, 2, \dots, t$  do
3.         run  $M_1$  on  $s_i$  for  $t$  steps
4.         run  $M_2$  on  $s_i$  for  $t$  steps
5.         if both accepted, then accept

If  $\langle M_1, M_2 \rangle \in L_4$  then there exists  $s_j \in L(M_1) \cap L(M_2)$ , and so there exist  $T_1$  and  $T_2$  such that  $M_1$  accepts  $s_j$  in  $T_1$  steps, and  $M_2$  accepts  $s_j$  in  $T_2$  steps. Let  $T = \max\{j, T_1, T_2\}$ . Since each step of the outer loop is finite, if  $M$  has not already accepted beforehand, then it reaches  $t = T$  and  $i = j$ , and accepts. On the other hand, if  $\langle M_1, M_2 \rangle \notin L_4$  then  $L(M_1) \cap L(M_2) = \emptyset$ , so  $M$  will not find a string that is accepted by both machines and will therefore not accept. Thus  $L(M) = L_4$ .

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**Part III (10 marks)**

Prove that the following language is not recognizable.

$$L_5 = \{\langle M_1, M_2 \rangle \mid L(M_1) \cap L(M_2) = \emptyset\}$$

Note that your answer will be marked on its structure as well as its content.

**Solution:**

Show  $\text{EMPTY} \leq_p L_5$ . Define  $f(\langle M \rangle) = \langle M_1, M_2 \rangle$ , where  $M_1 = M$  and  $M_2$  is a Turing machine that always accepts. Then  $f$  is computable, since it simply appends the description of a particular, fixed Turing machine to its input. Also,

$$\begin{aligned} \langle M \rangle \in \text{EMPTY} &\implies L(M) = \emptyset \\ &\implies L(M) \cap \Sigma^* = \emptyset \\ &\implies L(M_1) \cap L(M_2) = \emptyset \\ &\implies \langle M_1, M_2 \rangle \in L_5 \end{aligned}$$

$$\begin{aligned} \langle M_1, M_2 \rangle \in L_5 &\implies L(M_1) \cap L(M_2) = \emptyset \\ &\implies L(M) \cap \Sigma^* = \emptyset \\ &\implies L(M) = \emptyset \\ &\implies \langle M \rangle \in \text{EMPTY} \end{aligned}$$

So  $\langle M \rangle \in \text{EMPTY} \iff f(\langle M \rangle) \in L_5$ .

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