

## Hint for Question 5

This is a hint to help you come up with a loop invariant. Suppose that  $a$  is a natural number represented in decimal by  $a_{n-1} \cdots a_0$ , and  $b$  is a natural number represented in decimal by  $b_{n-1} \cdots b_0$ . By the definition of “decimal representation”,

$$a = \sum_{i=0}^{n-1} a_i \cdot 10^i$$

$$b = \sum_{i=0}^{n-1} b_i \cdot 10^i$$

Adding the two equations together, we have

$$(a + b) = \sum_{i=0}^{n-1} (a_i + b_i) \cdot 10^i$$

Notice that, if  $0 \leq (a_i + b_i) \leq 9$  for all  $i$ , then the sequence  $(a_{n-1} + b_{n-1}) \cdots (a_0 + b_0)$  is in fact the decimal representation of the sum  $(a + b)$ . Now suppose just that  $0 \leq (a_0 + b_0) \leq 9$ . Then the digit  $z_0$  computed by the algorithm is equal to  $a_0 + b_0$ , and so we have

$$(a + b) = \sum_{i=1}^{n-1} (a_i + b_i) \cdot 10^i + z_0 \cdot 10^0$$

In general, we cannot assume that  $(a_i + b_i)$  is between 0 and 9, which is where the carry  $c$  comes in. In trying to formulate a loop invariant predicate  $L(i)$ , I would suggest trying to come up with an expression which includes the first  $i$  computed digits  $z_{i-1} \cdots z_0$ , as well as the remaining terms  $(a_{n-1} + b_{n-1}) \cdots (a_i + b_i)$ , and the carry bit  $c$ , so that the expression is equal to  $a + b$  (like in the example above).