#### **CONVERGENCE OF ITERATIVE VOTING**

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Color of the new car...

Abel:

Seth:

# Adam: E **Can't we all just get along?** Cain:

# What we know: (Meir et al. – AAAI 2010)

Assuming players play a myopic "best response" – reacting to the current state:

#### **Iterative Plurality converges**

2 cases:

And and omized tie breaking rule: from en entrated between entrated ucestating to which ties are resolved
Deterministic tie breaking rules: from any state (including non-truthful)

# Tie-breaking rules Linear: $\bullet \succ \bullet \succ \bullet \succ \bullet \succ \bullet$ Non-linear:

#### There is no set order between red and orange

**Pastry example:** (thanks to Ilan Nehama)





#### Short aside: What are scoring rules

Scoring rules for *m* candidates define a scoring vector:

$$(lpha_1, lpha_2, lpha_3, \dots, lpha_m)$$

under the condition

 $\alpha_1 \ge \alpha_2 \ge \alpha_3 \ge \ldots \ge \alpha_m = 0$ 

A voter gives  $\alpha_1$  points to his most preferred candidate,  $\alpha_2$  points to his 2<sup>nd</sup> preference, etc.

The winner is the candidate with most points

# Short aside: Examples of scoring rules

**Plurality**:(1,0,...,0,0)

**Veto**: (1,1,...,1,0)

*k*-approval: 
$$(1,1,...,1,0,0,...,0)$$

*k*-veto: 
$$(1,1,...,1,0,0,...,0)$$

#### Theorem I: Tie-breaking rules matter

When using any arbitrary tiebreaking rule (i.e., not necessarily linear ones), every scoring rule & Maximin has tie-breaking rule for which it will not always converge

# Theorem I: Proof sketch (scoring rules)

4 candidates, 2 voters, tie breaking rule makes *c* win if not tied with *b*. *b* wins if not tied with *d*. *d* wins if not tied with *a*.

a > ... > b > c > d c > ... > d > b > a c > ... > d > b > a c > ... > d > b > a d > b > a d > ... > b > c > d b > ... > a > d > c d > ... > a > d > c

#### Theorem II: Borda doesn't work

When using the Borda voting rule, regardless of tie-breaking rules, the iterative process may never converge

#### Theorem II: Proof sketch

4 candidates, 2 voters (tie breaking doesn't matter):

a > b > c > d $\mathbf{b} \geq \mathbf{a} \geq \mathbf{d} \geq \mathbf{c}$ c > d > b > a $\mathbf{c} > \mathbf{d} > \mathbf{b} > \mathbf{a}$ d - 2; a, b - 3; c - 4a - 2; c, d - 3; b - 4a > b > c > d $\mathbf{b} \geq \mathbf{a} \geq \mathbf{d} \geq \mathbf{c}$  $d \geq c \geq a \geq b$ d > c > a > bb-2; c, d-3; a-4c - 2; a, b - 3; d - 4

# Theorem III: Iterative Veto converge*r*

When using linear tie-breaking rules, iterative Veto will always converge – from truthful or nontruthful starting point

#### Theorem III: Proof

"Best response" straight-forwardly defined as vetoing the current (unwanted) winner.

Lemma 1: If there is a cycle, taking a stage in the cycle where there is more than one candidate with the maximal score, suppose winner score is *s*. Then **winning score at any other stage is** *s* **or** s+1. Any stage with s+1 score has only one candidate with that score.

#### Theorem III: Proof Lemma I

The futility of having a single winner – the score can't get higher, and you can't get multiple candidates to share the score:



#### Theorem III: Proof

Lemma 2: If there is a cycle, all stages with more than one candidate with the maximal score have the same number of candidates with maximal score and maximal-1 score, and these are the same candidates in all the cycle.



#### Theorem III: Proof

2 types of player moves:

A candidate with a score of s becomes winner with score of s+1

A candidate with a score of *s*-1 gets point and becomes winner

Previously vetoed candidates become winners (gaining a point), i.e., voters' situation progressively worse. This is a **finite process** 

# Theorem IV: *k*-Approval doesn't work

When using *k*-approval voting rule for *k*≥2, even with linear tiebreaking rule, the iterative process may never converge

#### Theorem IV: Proof sketch

4 candidates, 2 voters, and the tie breaking rule is alphabetical (a > b > c > d)



# Current problems: lazy-best Borda (with Maria Polukarov)

Lazy-best means we put the new winner in 1<sup>st</sup> place, and push everyone else back one spot.

Does this converge with Borda?

Using a simulator, it seems lazy-best Borda converges.

If we don't allow ties, it's easy to prove convergence. **Tie-breaking is key**. Score increase may be high (up to m-1 points), but points are lowered one point at a time – so a cycle has many stages in which maximal score is either static or gets lowered.

# Current problems: Polynomial Veto (with Maria Polukarov)

Plurality converges after a polynomial number of steps.

Does Veto converge in polynomial time?

Many characteristics found in convergence proof apply:

After initial moves, only candidates with top two scores are relevant

4 types of moves:



#### **Future work**

Better understanding of what influences convergence (tie-breaking rules identified, what else?)

What is best-response for complex voting rules?

Moving beyond myopic best-response to more complex and varied responses

Computational complexity issues for best-response in complex voting rules

Weighted games





(guess they decided to compromise on the car colors...)

Thanks for listening!