

## CSC 2410 Assignment #3, Fall 2017

**Due:** Tuesday Dec 5 at 10:00 AM

You may consult the text. You may not consult any other books and materials.

You may consult with each other **only on problem 6**. For the other problems, you can only consult with the instructor and the TA for this course.

Of course, each problem requires a well-written proof. Proofs that are unnecessarily lengthy might not get full marks, even if they are correct. And they might not be read thoroughly by the grader.

1. **(5 pts)** Either find a planar embedding of the Petersen Graph or find a  $K_5$ - or  $K_{3,3}$ -subdivision in the Petersen Graph. (The Petersen Graph is on the cover of the text; see also Def 1.1.36.)
2. **(10 pts)** Prove that a planar embedded graph is bipartite iff every face has even size.
3. **(25 pts)** 6.2.9 from West.
4. **(35 pts)** In this problem, you may not apply the Four Colour Theorem.
  - (a) Prove that if  $G$  is a connected planar graph such that (i)  $G$  has maximum degree at most 5, and (ii)  $G$  has at least one vertex of degree less than 5, then  $\chi(G) \leq 4$ .
  - (b) Prove that if  $G$  is a 3-connected 5-regular planar graph then  $\chi(G) \leq 4$ .

**Hint:** Choose two non-adjacent vertices  $a, b$  and contract them into a single vertex. If you choose  $a, b$  wisely, then the resultant graph will be planar, and you will be able to prove that it cannot contain a 5-critical subgraph. Then argue that this implies  $\chi(G) < 5$ .
  - (c) Prove that if  $G$  is a planar graph with maximum degree at most 5 then  $\chi(G) \leq 4$ .

**Note:** If you do not solve (a) or (b), then you can assume that what they assert is true and still attempt (c).
5. **(15 pts)** You are given a multigraph with a label from  $\{1, \dots, k\}$  on each edge; there are no loops but there may be multiple edges. Your goal is to assign a label from  $\{1, \dots, k\}$  to each vertex so that there is no edge  $uv$  where  $u, v$  and  $uv$  all have the same label. This is a variation of  $k$ -colouring where each edge only forbids one colour from appearing on both endpoints; if  $uv$  is labelled 1 then  $u$  and  $v$  can have the same colour as long as it is not 1.

Use the probabilistic method to prove that if the number of edges is at most  $k^2$  then there is a way to label the vertices as required.
6. **(20 pts)** Suppose that  $G$  has  $n$  vertices and minimum degree  $\delta$ . Suppose further that for every two vertices  $x, y$  in  $G$ ,  $|N(x) \cup N(y)| + \delta \geq n + 10$ . Prove that  $G$  has a Hamilton cycle.

**Remark:** The “+10” term is much higher than necessary.