Basics:

Network Classification Network Architecture Reliable Data Transfer Delay Models Implementation: Protocol Design



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Application Transport Network Data Link Physical

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Layered Architecture



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Functionality

- Reliable Delivery of Frames
- Flow Control
- Error Detection
- Error Correction

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Multiaccess Media

Ethernet

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Multiaccess Media

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Cocktail Party

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- "Raise your hand if you have a question"
- "Give everyone a chance to speak"

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- Channel Partitioning (Cellular Wireless Networks)
- Random Access (Ethernet, WiFi)
- Taking Turns (Token Ring)

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Ethernet

- Hosts broadcast packets
- When a collision occurs, all transmitted packets are lost
- Lost packets have to be retransmitted



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=> Need Multiaccess Protocol

Goal:

- Understand Multiaccess Protocols
- Understand Ethernet and IEEE 802.11 Protocol

Issues:

- How to deal with collisions? (- > Protocol design)
- Maximal traffic load? (-> Protocol performance)

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- Packet arrival rate (over all hosts) of λ packets/time unit
- Collision or Perfect Reception
- Immediate Feedback: 0, 1, e
- (Re-)transmission Probability: q_r
- Infinite number of hosts (i.e. each node has at most one packet to transmit)

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• How to choose q_r ?

- Would $q_r = 1$ work?
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- Average time until (re-)transmission:

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Notation

- λ : aggregated arrival rate
- n: number of backlogged packets
- $G(n) = nq_r$: average number of transmissions per time slot

Want to compute

- *P*_{succ}: probability of successful transmission in a time slot (as a function of *G*(*n*))
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$$P_{succ} = nq_r (1-q_r)^{n-1} = \frac{nq_r}{1-q_r} (1-q_r)^n$$

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For q_r small, we have

$$\left(1-q_r\right)^n \approx e^{-nq_r}$$
 and $\frac{nq_r}{1-q_r} \approx nq_r$,

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 and $\frac{nq_r}{1-q_r} \approx nq_r$,

and we obtain

$$P_{succ} \approx nq_r e^{-nq_r} = G(n)e^{-G(n)}$$

where $G(n) = nq_r$

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Throughput:

 $G(n)e^{-G(n)}$

where $G(n) = nq_r$



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Note: Arrivals according to a Poisson distribution with rate G:

$$p_k = \frac{G^k}{k!} e^{-G}$$

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$$q_r = \frac{1}{n}$$

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- If $G(n)e^{-G(n)} > \lambda$:
- If $G(n)e^{-G(n)} < \lambda$:
- Optimal $G(n) = nq_r = 1$, or

$$q_r = \frac{1}{n}$$

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• If $G(n)e^{-G(n)} < \lambda$:

• Optimal $G(n) = nq_r = 1$, or

$$q_r = \frac{1}{n}$$

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• Optimal
$$G(n) = nq_r = 1$$
, or

$$q_r = \frac{1}{n}$$

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What did we learn?

- $\lambda_{max} = e^{-1} \approx 0.368$
- q_r should dynamically change

Binary Exponential Backoff

• $q_r = 2^{-k}$



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