

## wait and waitpid (11.2)

- Recall from a previous slide: `pid_t wait( int *status )`
- `wait()` can: (a) block; (b) return with status; (c) return with error
- If there is more than one child, `wait()` returns on termination of *any* children
- `waitpid` can be used to wait for a specific child pid
- `waitpid` also has an option to block or not to block

```
pid_t waitpid( pid_t pid, &status, option );
    pid      == -1      waits for any child
    option == NOHANG    non-blocking
    option == 0         blocking
waitpid(-1, &status, 0) equivalent to wait(&status)
```

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## example: wait.c

```
#include <sys/types.h>
#include <sys/wait.h>
void main( void )
{
    int status;
    if( fork() == 0 ) exit( 7 );      /* normal exit */
    wait( &status ); prExit( status );

    if( fork() == 0 ) abort();        /* generates SIGABRT */
    wait( &status ); prExit( status );

    if( fork() == 0 ) status /= 0;    /* generates SIGFPE */
    wait( &status ); prExit( status );
}
```

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## prExit.c

```
#include <sys/types.h>
#include <sys/wait.h>
void prExit( int status )
{
    if( WIFEXITED( status ) )
        printf( "normal termination, exit status = %d\n",
                WEXITSTATUS( status ) );
    else if( WIFSIGNALED( status ) )
        printf( "abnormal termination, signal number = %d\n",
                WTERMSIG( status ) );
    else if( WIFSTOPPED( status ) )
        printf( "child stopped, signal number = %d\n",
                WSTOPSIG( status ) );
}
```

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## exec

- Six versions of exec:

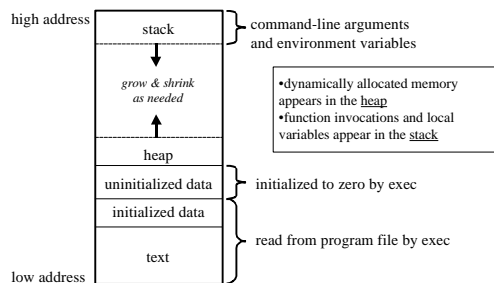
```
execl( char *pathname, char *arg0, ... , (char*) 0 );
execv( char *pathname, char *argv[] );

execlp( char *pathname, char *arg0, ..., (char*) 0,
        char *envp[] );
execvp( char *pathname, char *argv[],
        char *envp[] );

execlp( char *filename, char *arg0, ..., (char*) 0 );
execvp( char *filename, char *argv[] );
```

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## Memory Layout of a C program



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## Miscellaneous: permissions

- Read permissions for a directory and execute permissions for it are not the same:
  - Read**: read directory, obtain a list of filenames
  - Execute**: lets users pass through the directory when it is a component of a pathname being accessed
- Cannot create a new file in a directory unless user has write permissions and execute permission in that directory
- To delete an existing file, the user needs write and execute permissions in the directory containing the file, but does not need read or write permission for file itself (!!!)

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## Miscellaneous: buffering control

```
int setbuffer(FILE *fp, char *buf, int size)
```

- specifies that “**buf**” should be used instead of the default system-allocated buffer, and sets the buffer size to “**size**”
- if “**buf**” is **NULL**, i/o will be unbuffered
- used after stream is opened, but before it is read or written

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```
int setlinebuf( FILE *fp )
```

- used to change **stdout** or **stderr** to line buffered
- can be called anytime

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- A stream can be changed from unbuffered or line buffered to block buffered by using **freopen()**. A stream can be changed from block buffered or line buffered to unbuffered by using **freopen()** followed by **setbuf()** with a buffer argument of **NULL**.

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## Signals

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## Motivation for Signals (11.15)

- When a program forks into 2 or more processes, rarely do they execute independently of each other
- The processes usually require some form of synchronization, and this is typically handled using signals
- Data usually needs to be passed between processes also, and this is typically handled using pipes and sockets, which we'll discuss in detail in a week or two
- Signals are usually generated by
  - machine interrupts
  - the program itself, other programs, or the user (*e.g.* from the keyboard)

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## Introduction

- **<sys/signal.h>** lists the signal types on cdf. Table 11.5 and **signal(5)** give a list of some signal types and their default actions
- When a C program receives a signal, control is immediately passed to a function called a signal handler
- The signal handler function can execute some C statements and exit in three different ways:
  - return control to the place in the program which was executing when the signal occurred
  - return control to some other point in the program
  - terminate the program by calling the **exit** (or **\_exit**) function

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## sigset()

- A default action is provided for each kind of signal, such as terminate, stop, or ignore
- For nearly all signal types, the default action can be changed using the **signal()** function. The exceptions are **SIGKILL** and **SIGSTOP**
- Usage: **signal(int sig, void (\*disp)(int))**
- For each process, UNIX maintains a table of actions that should be performed for each kind of signal. The **signal()** function changes the table entry for the signal named as the first argument to the value provided as the second argument
- The second argument can be **SIG\_IGN** (ignore the signal), **SIG\_DFL** (perform default action), or a pointer to a signal handler function

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## sigset() example

```
#include <stdio.h>
#include <stdlib.h>
#include <sys/signal.h>
int i = 0;
void quit( int code ) {
    fprintf( stderr, "\nInterrupt (code=%d, i=%d)\n", code, i );
    exit( 123 );
}
void main( void ) {
    if ( signal( SIGINT, quit ) == SIG_IGN ) exit( 1 );
    if ( signal( SIGTERM, quit ) == SIG_IGN ) exit( 2 );
    if ( signal( SIGQUIT, quit ) == SIG_IGN ) exit( 3 );
    for(;;)
        if( i++ % 5000000 == 0 ) putc( '.', stderr );
}
```

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## Checking the return value

- The data type that `signal()` returns is:  
*pointer to function with `int` argument returning `void`*
- So, the variable used to hold the result of a call to `signal` should be declared as follows:  

```
void (*signal_result)(int);
```
- It is possible for a child process to accept signals that are being ignored by the parent, which more than likely is undesirable
- Thus, the normal method of installing a new signal handler is:  

```
oldhandler = sigset( SIGHUP, SIG_IGN );  
if( oldhandler != SIG_IGN )  
    sigset( SIGHUP, newhandler );
```

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## Signalling between processes

- One process can send a signal to another process using the misleadingly named function call  

```
kill( int pid, int sig )
```
- This call sends the signal "`sig`" to the process "`pid`"
- Signalling between processes can be used for many purposes:
  - kill errant processes
  - temporarily suspend execution of a process
  - make processes aware of the passage of time
  - synchronize the actions of processes

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## Timer signals

- Three interval timers are maintained for each process:
    - `SIGALRM` (real-time alarm, like a stopwatch)
    - `SIGVTALRM` (virtual-time alarm, measuring CPU time)
    - `SIGPROF` (used for profilers, which we'll cover later)
  - Useful functions to set and get timer info are:
    - `setitimer()`, `getitimer()`
    - `alarm()` (simpler version: only sets `SIGALRM`)
    - `pause()` (suspend until next signal arrives)
    - `sleep()` (caused calling process to suspend)
    - `usleep()` (like `sleep()`, but with finer granularity)
- Note: `sleep()` and `usleep()` are *interruptible* by other signals

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