

# CHAPTER 4

---











## ALGORITHM ANALYSIS AND ASYMPTOTIC NOTATION

Feb 25, 2015

Lisa Yan

# Linear Search

```
def LS(A, x):  
    """ Return index i, x == A[i].  
        Otherwise, return -1 """
```

```
1. i = 0   
2. while i < len(A):     
3.     if A[i] == x:     
4.         return i   
5.     i = i + 1    
6. return -1
```

$t_{LS}([2, 4, 6, 8], 6) = 10$

What is the runtime of  
 $LS(A, x)$ ?

if the first index where  $x$  is  
found is  $k$   
i.e.,  $A[k] == x$

$$\begin{aligned} t_{LS}(A, x) &= 1 + 3(k+1) \\ &= 3k + 4 \end{aligned}$$

# Today's Outline

- ❖ Formal definition of  $O$ ,  $\Omega$ ,  $\Theta$

# FORMAL DEFINITIONS OF 0, Ω, Θ

---

# Recap $O(n^2)$

Set of functions that **grow no faster** than  $n^2$

- ❖ count the number of steps
- ❖ constant factors don't matter
- ❖ only highest-order term matter

The following functions are in  $O(n^2)$

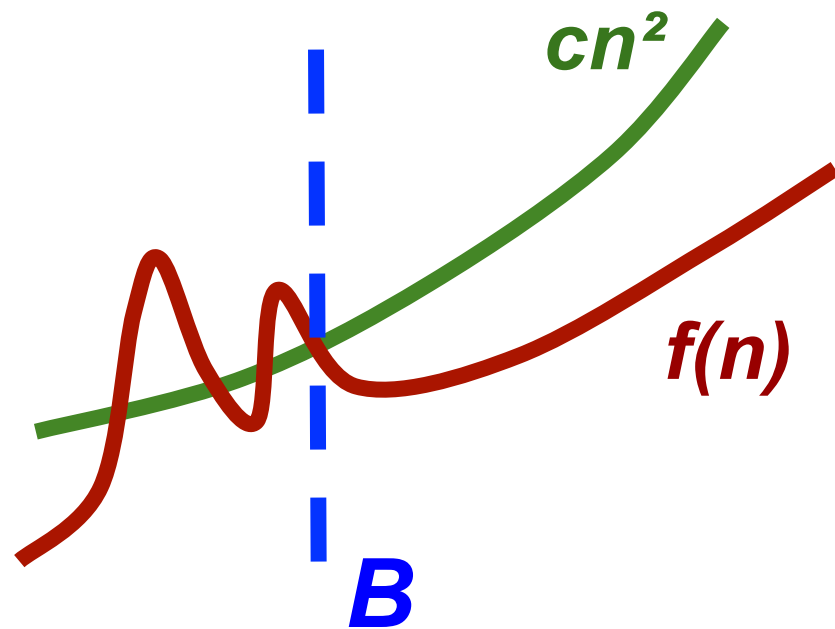
$$n^2 \quad 2n^2 + 3n \quad \frac{n^2}{165} + 1130n + 3.14159$$

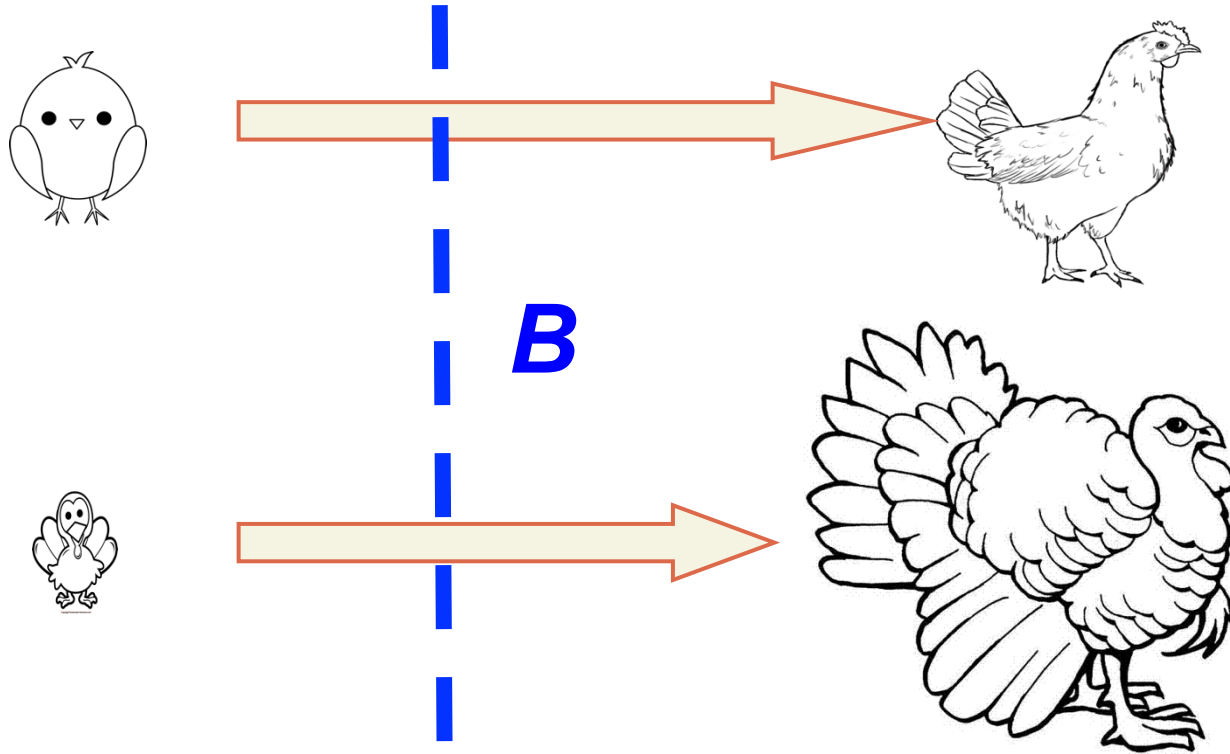
# Formal definition of $O(n^2)$

a function  $f(n)$  is in  $O(n^2)$  iff

$\exists c \in \mathbb{R}^+, \exists B \in \mathbb{N}$ , such that  $\forall n \in \mathbb{N}, n \geq B \Rightarrow f(n) \leq cn^2$

Beyond **breakpoint  $B$** ,  
 **$f(n)$**  is upper-bounded  
by  **$cn^2$** , where  **$c$**  is  
some wisely chosen  
constant multiplier.





“chicken size” <sup>$B$</sup>  is in  $O$ (“turkey size”)

A chicken **grows slower** than a turkey in the sense that, after a certain **breakpoint**, a chicken will always be smaller than a turkey.

# Formal Definition $O(n^2)$

a function  $f(n)$  is in  $O(n^2)$  iff

$\exists c \in \mathbb{R}^+, \exists B \in \mathbb{N}$ , such that  $\forall n \in \mathbb{N}, n \geq B \Rightarrow f(n) \leq cn^2$

**Simple example: prove  $700n^2 \in O(n^2)$**

Pick  $c = 711$ , or any real number  $\geq 700$

Pick  $B = 0$ , or any natural number  $\geq 0$

then  $\forall n \in \mathbb{N}, n \geq 0 \Rightarrow 700n^2 \leq 711n^2$

then  $\exists c \in \mathbb{R}^+, \exists B \in \mathbb{N}, \forall n \in \mathbb{N}, n \geq B \Rightarrow 700n^2 \leq cn^2$

then  $700n^2 \in O(n^2)$



# Formal Definition $\Omega(n^2)$

a function  $f(n)$  is in  $O(n^2)$  iff

**upper-bound**

$\exists c \in \mathbb{R}^+, \exists B \in \mathbb{N}$ , such that  $\forall n \in \mathbb{N}, n \geq B \Rightarrow f(n) \leq \underline{cn^2}$

a function  $f(n)$  is in  $\Omega(n^2)$  iff

**lower-bound**

$\exists c \in \mathbb{R}^+, \exists B \in \mathbb{N}$ , such that  $\forall n \in \mathbb{N}, n \geq B \Rightarrow f(n) \geq \underline{cn^2}$

**$O(n^2)$** : set of functions that **grow no faster than  $n^2$**

**$\Omega(n^2)$** : set of functions that **grow no slower than  $n^2$**

**$\Theta(n^2)$** : set of functions that are in **both  $O(n^2)$  and  $\Omega(n^2)$**

(functions growing **as fast as  $n^2$** )

## Growth rate ranking of some common functions

$$f(n) = n^n$$

$$f(n) = 2^n$$

$$f(n) = n^3$$

$$f(n) = n^2$$

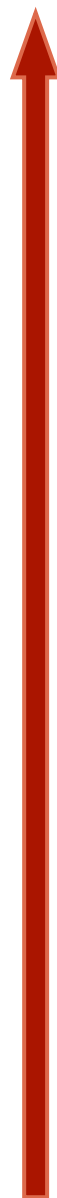
$$f(n) = n \log n$$

$$f(n) = n$$

$$f(n) = \sqrt{n}$$

$$f(n) = \log n$$

$$f(n) = 1$$



**grow fast**

**grow slowly**

# Examples

$$7n \in \mathcal{O}(n^2)$$

$$7n \notin \Omega(n^2)$$

$$7n^3 \notin \mathcal{O}(n^2)$$

$$7n^3 \in \Omega(n^2)$$

$$7n^2 \in \mathcal{O}(n^2)$$

$$7n^2 \in \Omega(n^2)$$

$$7n^2 \in \Theta(n^2)$$

# Next Week

- Worst-case analysis of two algorithms