

Semantic Slicing of Software Version Histories

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ASE 2015 / Lincoln, NE

Motivation



Feb, 2015



release [1.3.8]

...



make 'groovy.sandbox.blacklist' append-only

avoid NullPointerException if optional Groovy jar is removed

make groovy sandbox method blacklist dynamically additive

...



prepare for next development iteration (1.3.7-SNAPSHOT)



updated docs to use v1.3.6 as current

Nov, 2014



release [1.3.6]

30 authors
67 commits
87 files changed

Motivation



Motivation

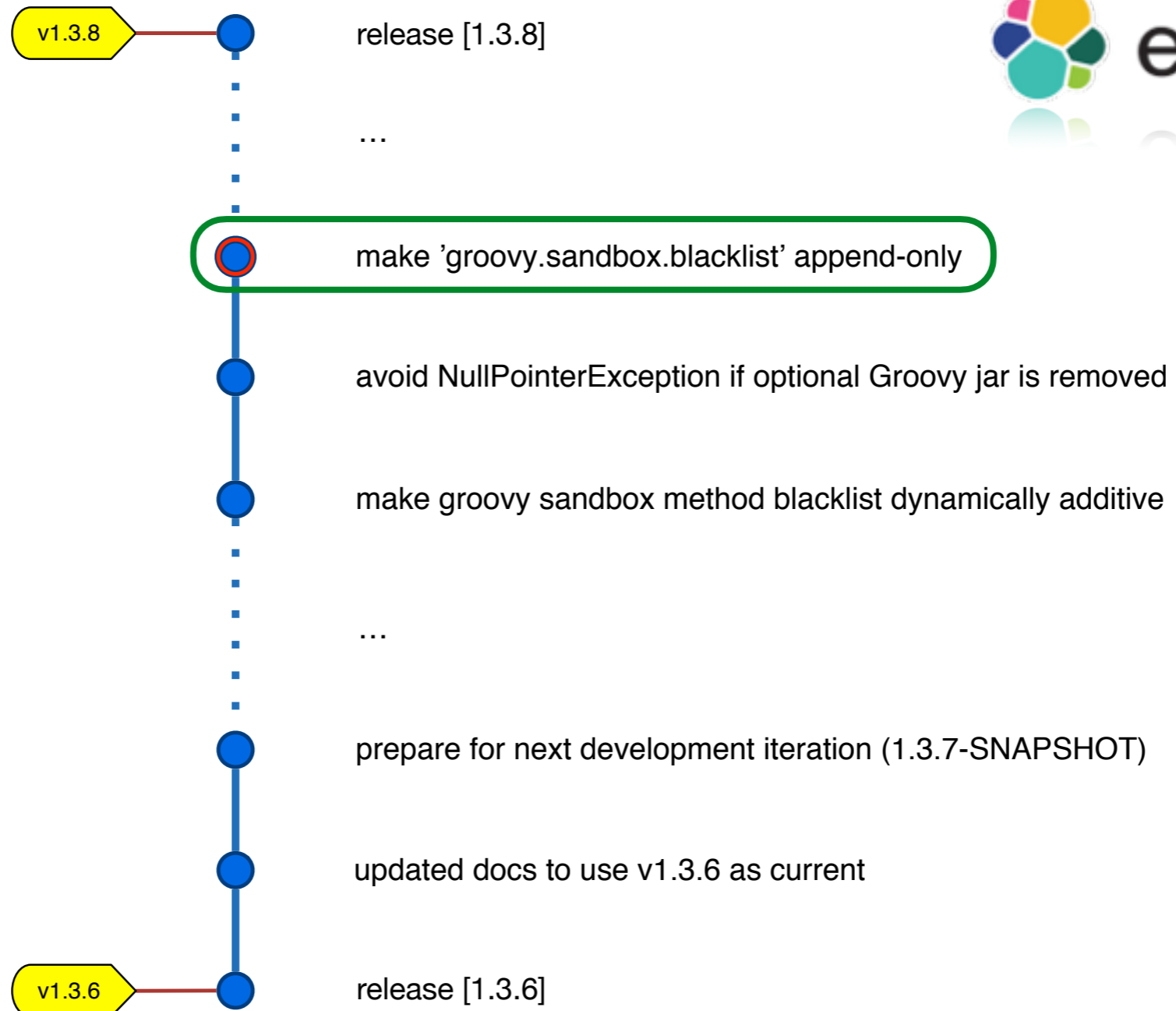


Feb, 2015

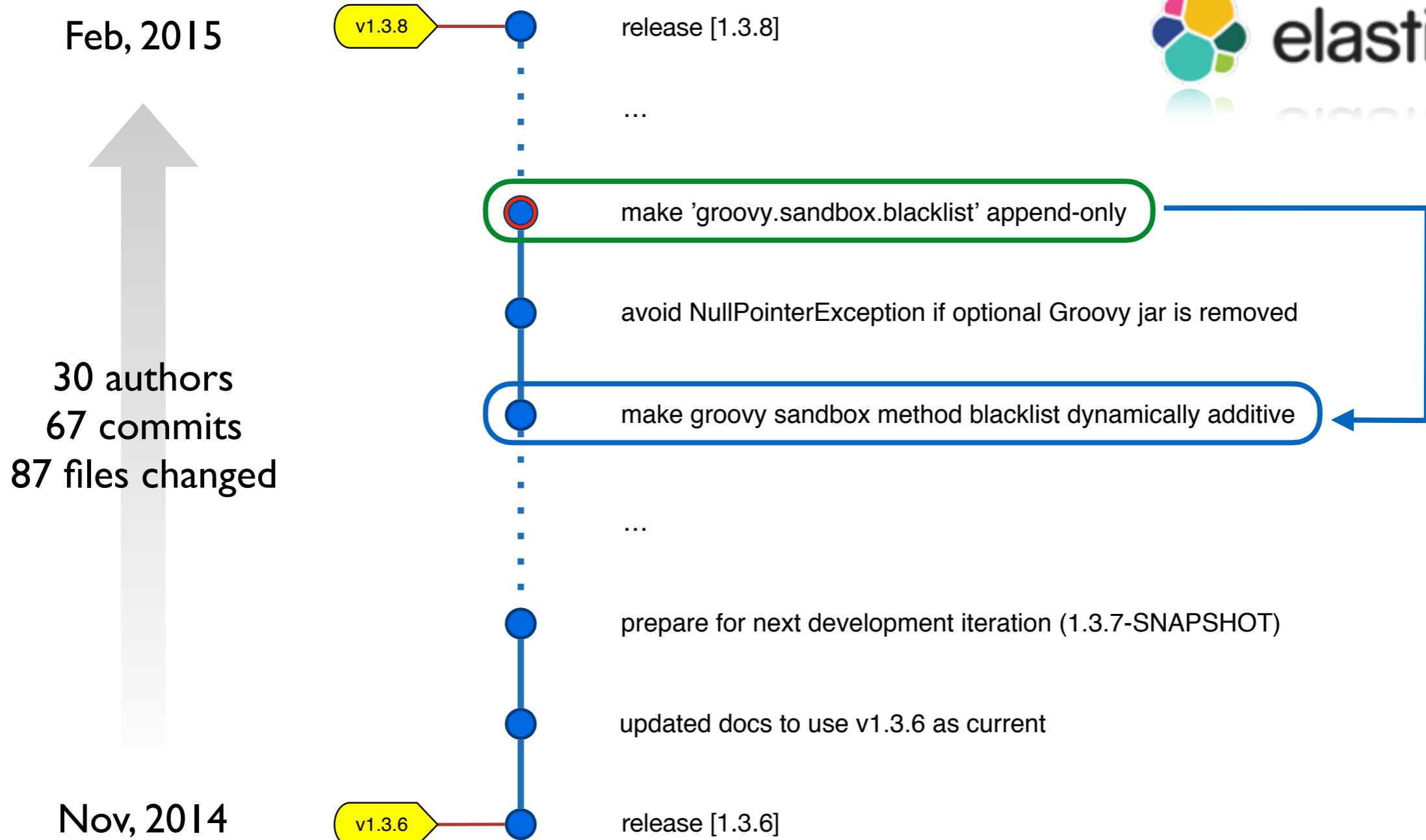


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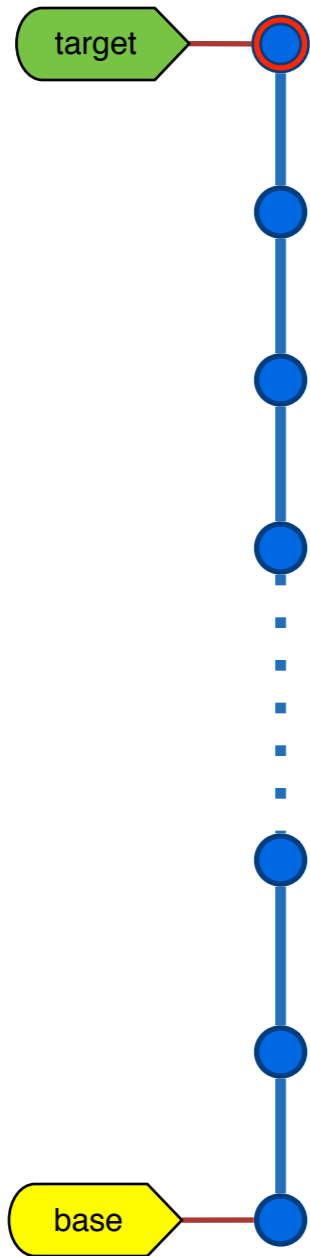
Nov, 2014



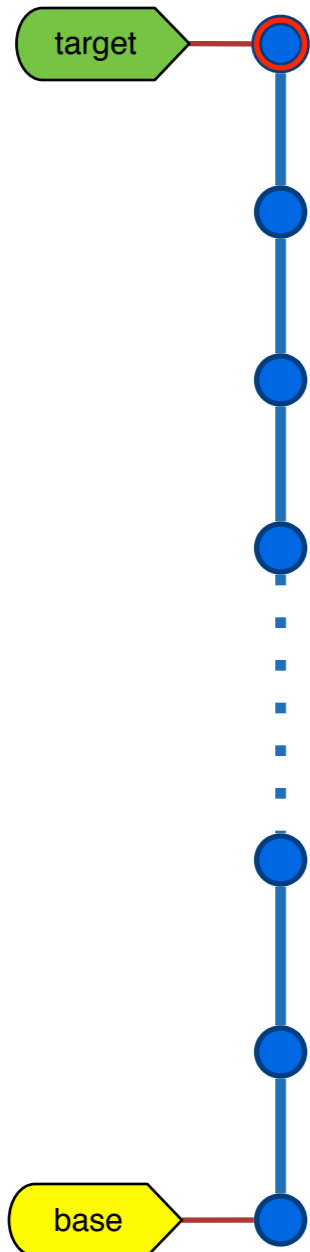
Motivation



Why is it so hard?



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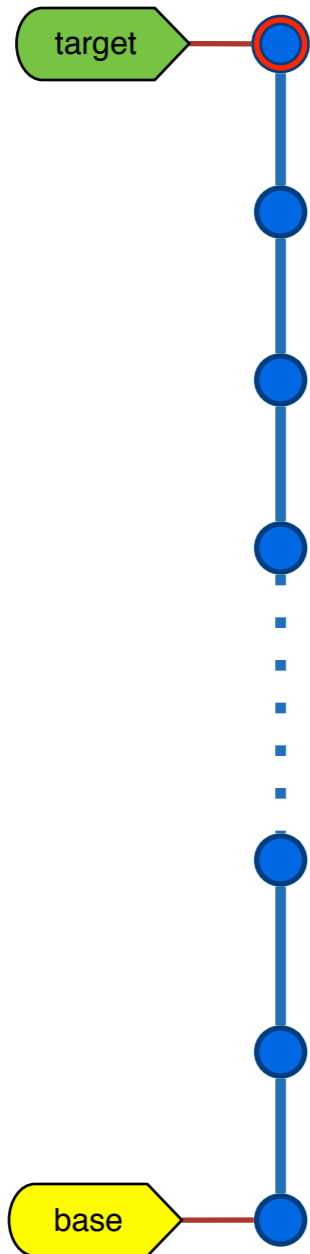


Options?



1. Pick target commits
2. Pick the whole history
3. Manually identify necessary commits

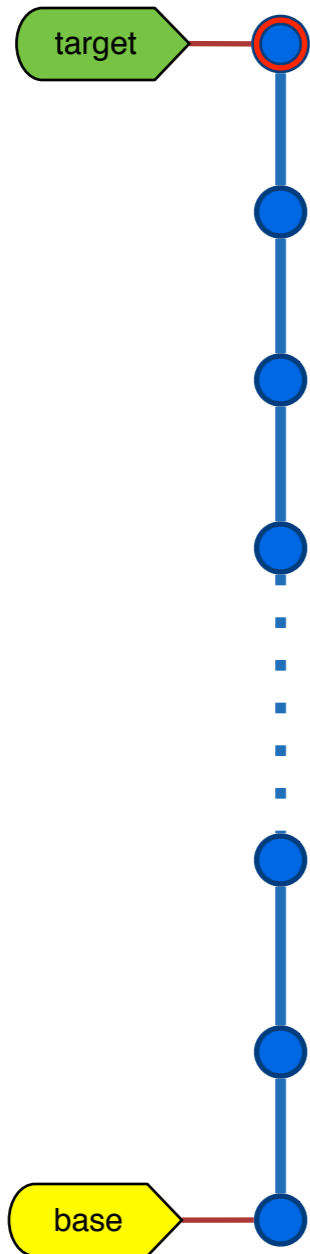
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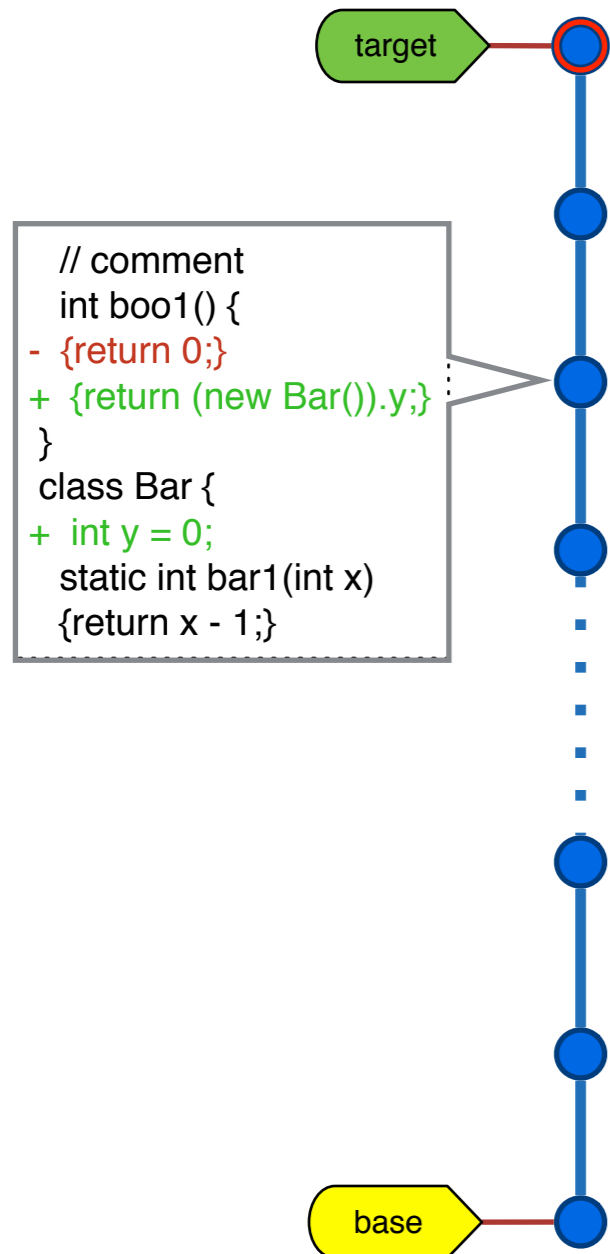
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Why is it so hard?



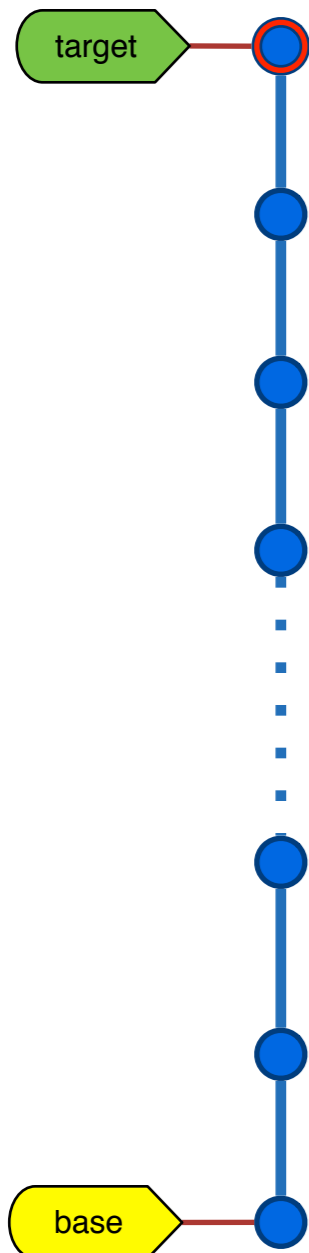
Options?

- ~~X~~ 1. Pick target commits
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Existing version control tools:

- Code treated as plain texts
- Do not understand the semantics
- User provided semantic/logical grouping is inaccurate!

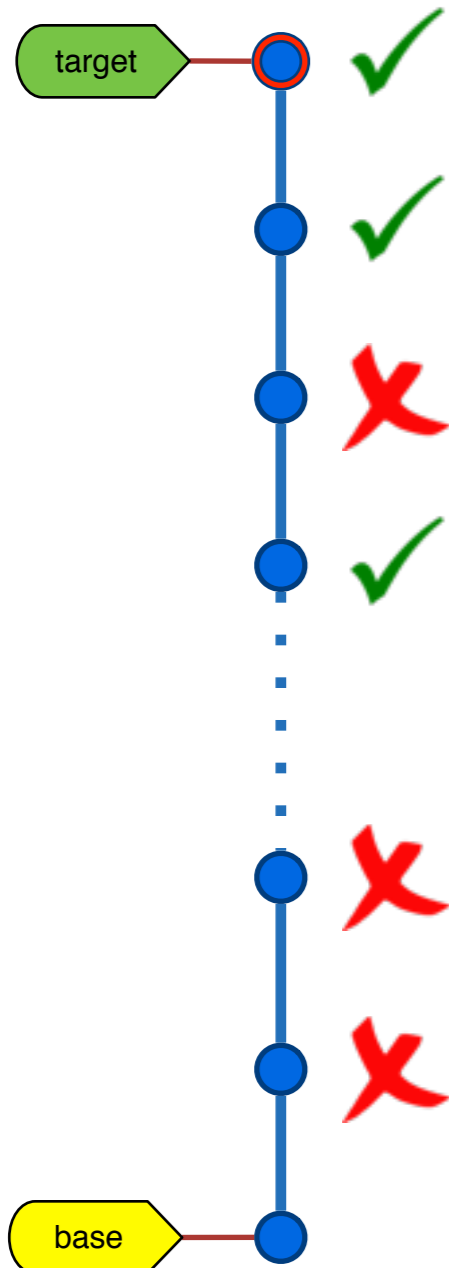
What can we do?



Exploit existing artifacts:

- Strictly structured data
- Well-defined language syntax and semantics
- Carefully designed test suites

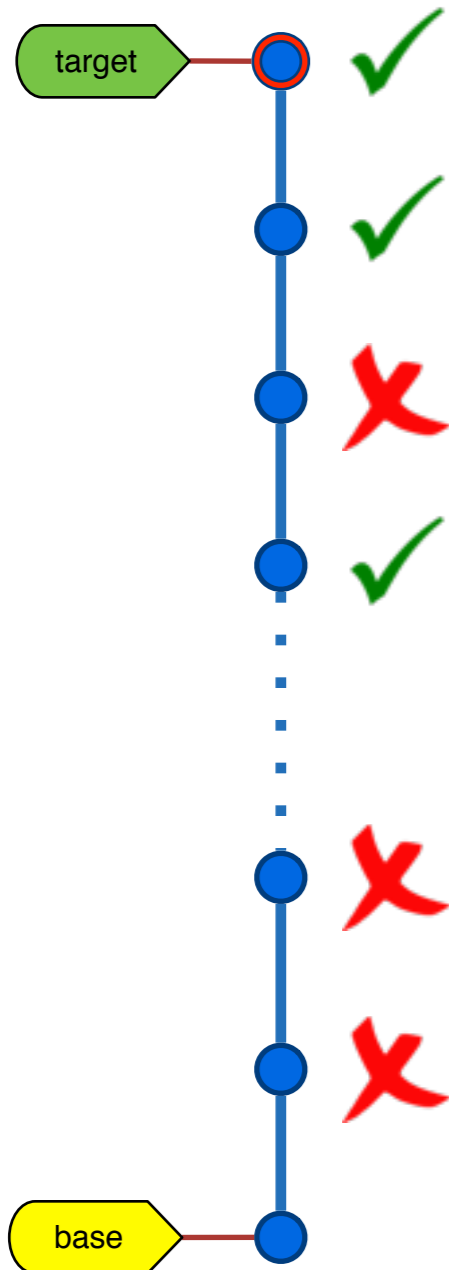
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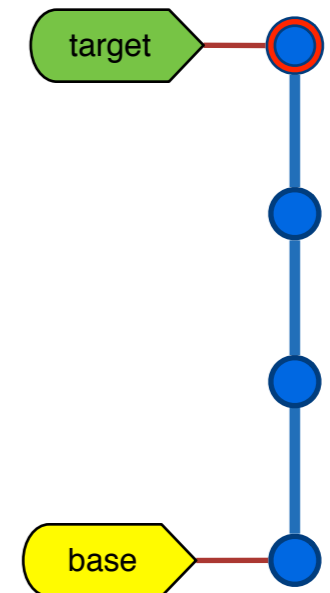
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Solution: Semantic Slicing



Exploit existing artifacts:

- Strictly structured data
- Well-defined language syntax and semantics
- Carefully designed test suites



History:
sequence of commits
+
Criterion:
set of tests

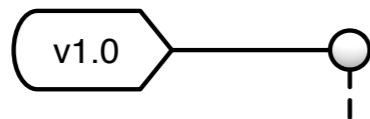


Sub-history:
well-formed: compiles
&
semantic preserving:
passing tests

Outline

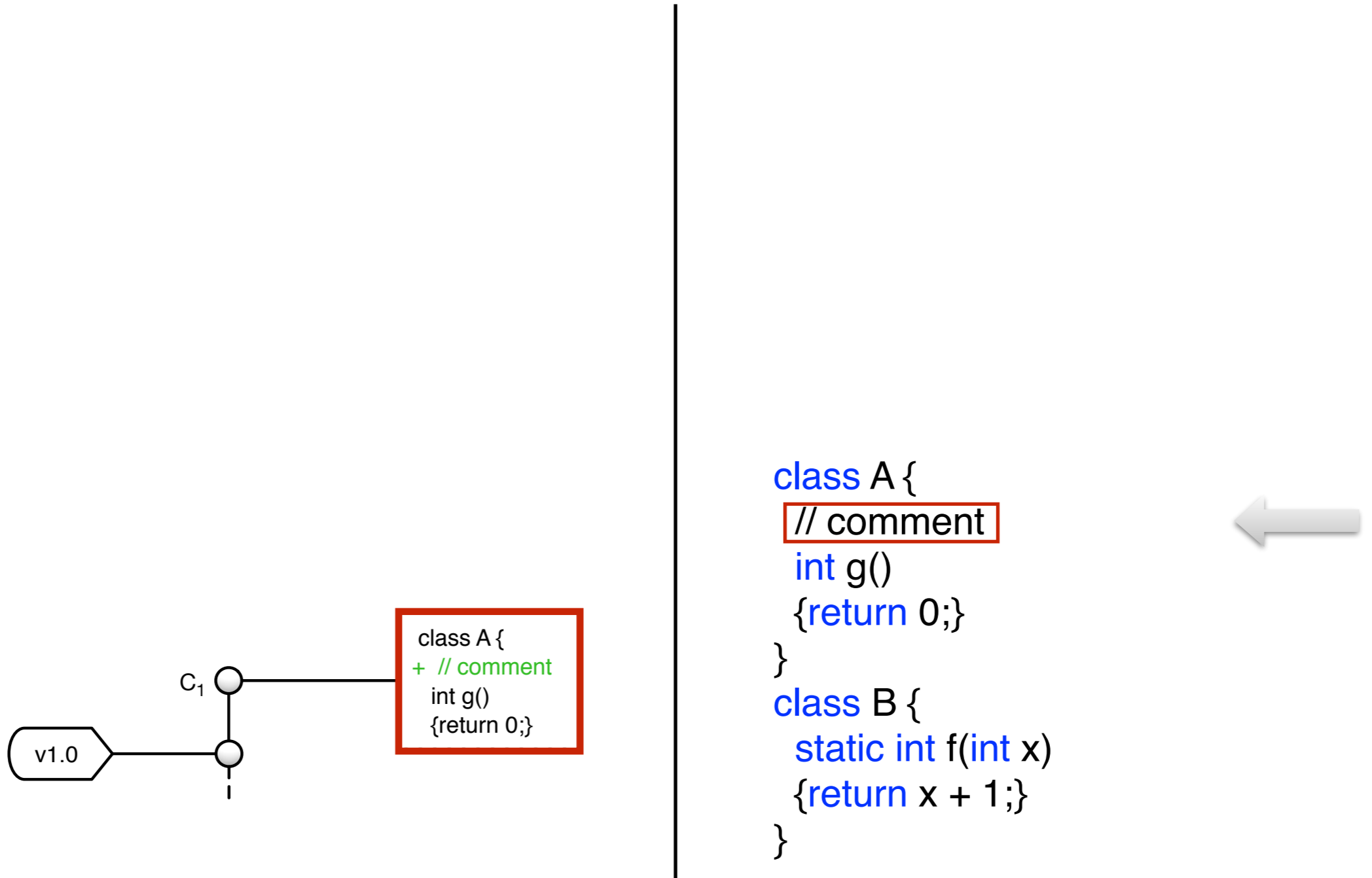
1. Introduction
2. Dependency Hierarchy
3. CSlicer Algorithm
4. Evaluation
5. Related Work & Conclusion

Running Example



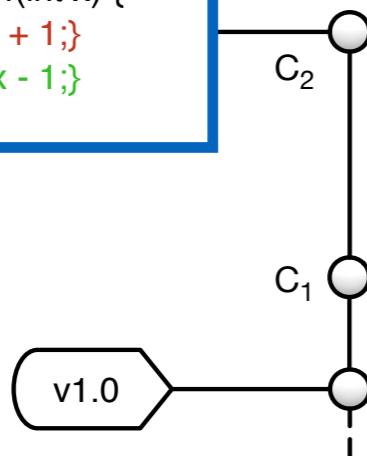
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class A {  
    int g()  
    {return 0;}  
}  
class B {  
    static int f(int x)  
    {return x + 1;}  
}
```

Running Example



Running Example

```
class A {  
  static int f(int x) {  
- {return x + 1;}  
+ {return x - 1;}  
}
```

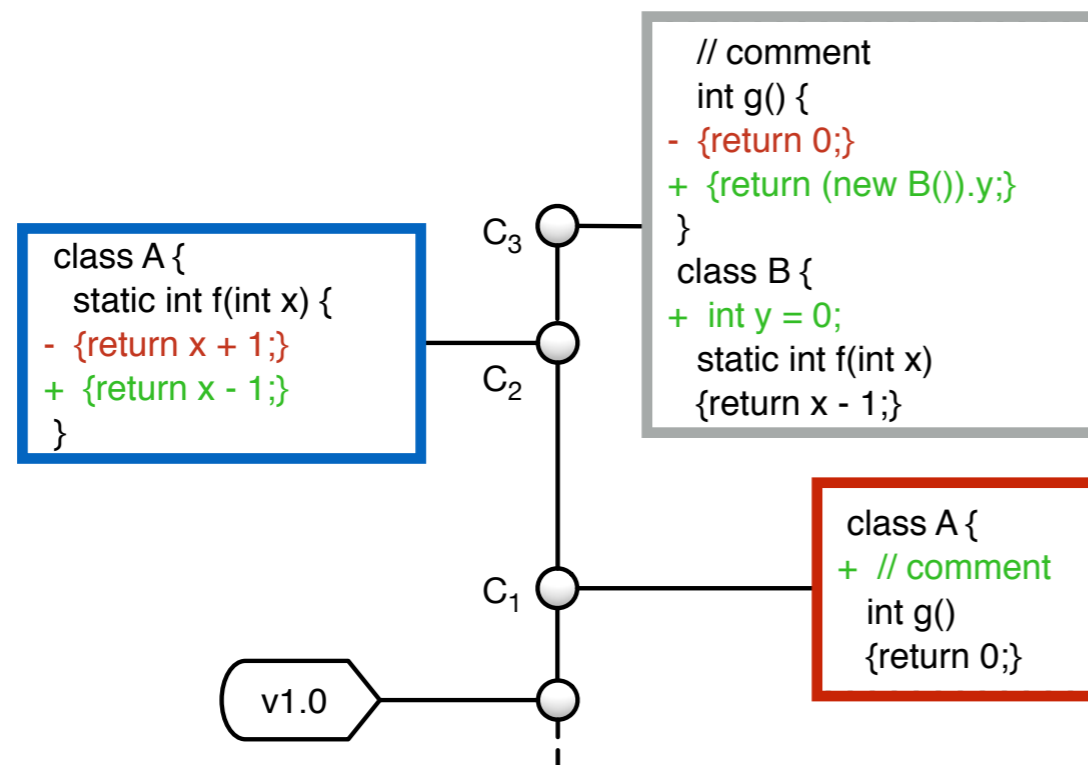


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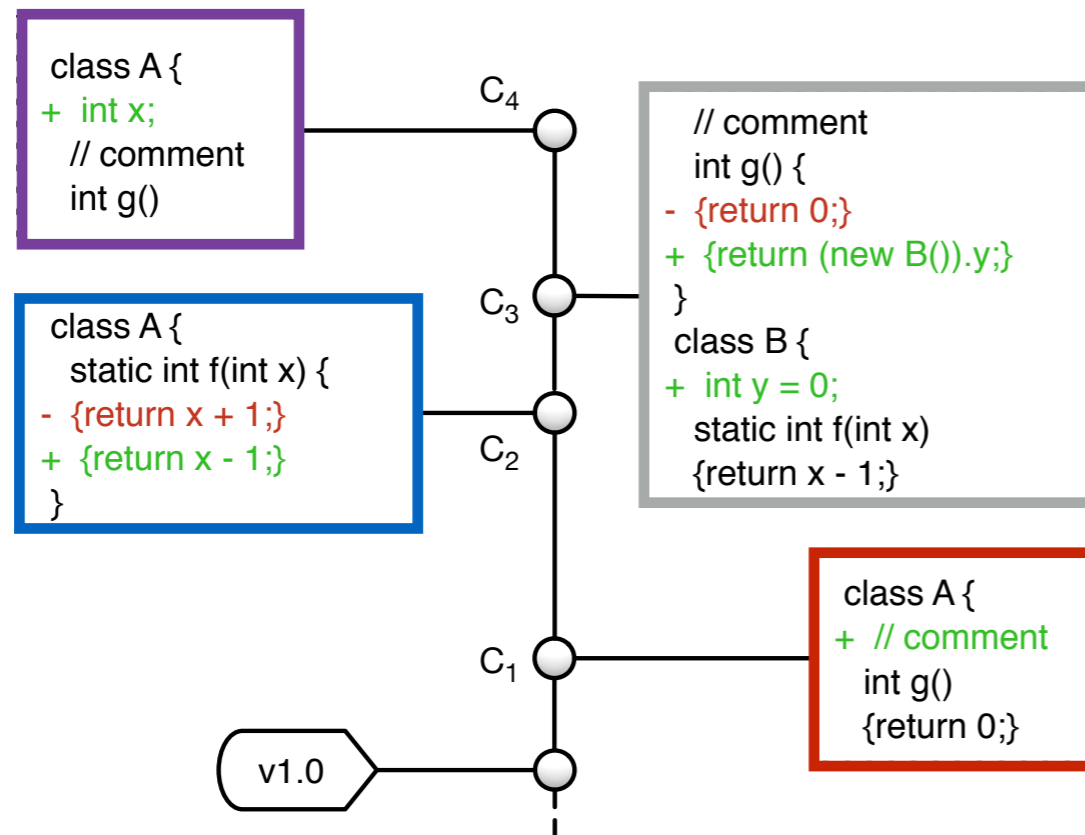
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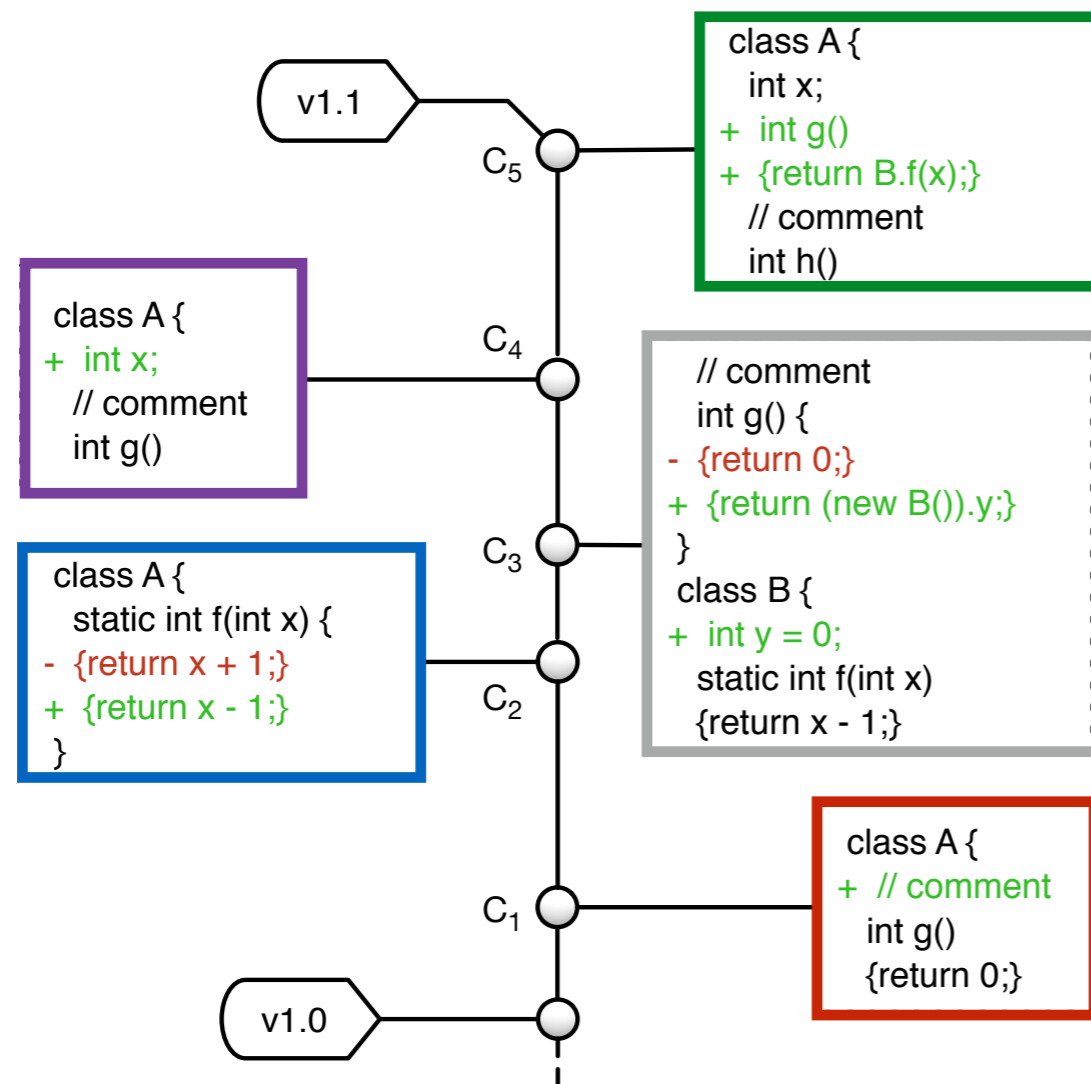


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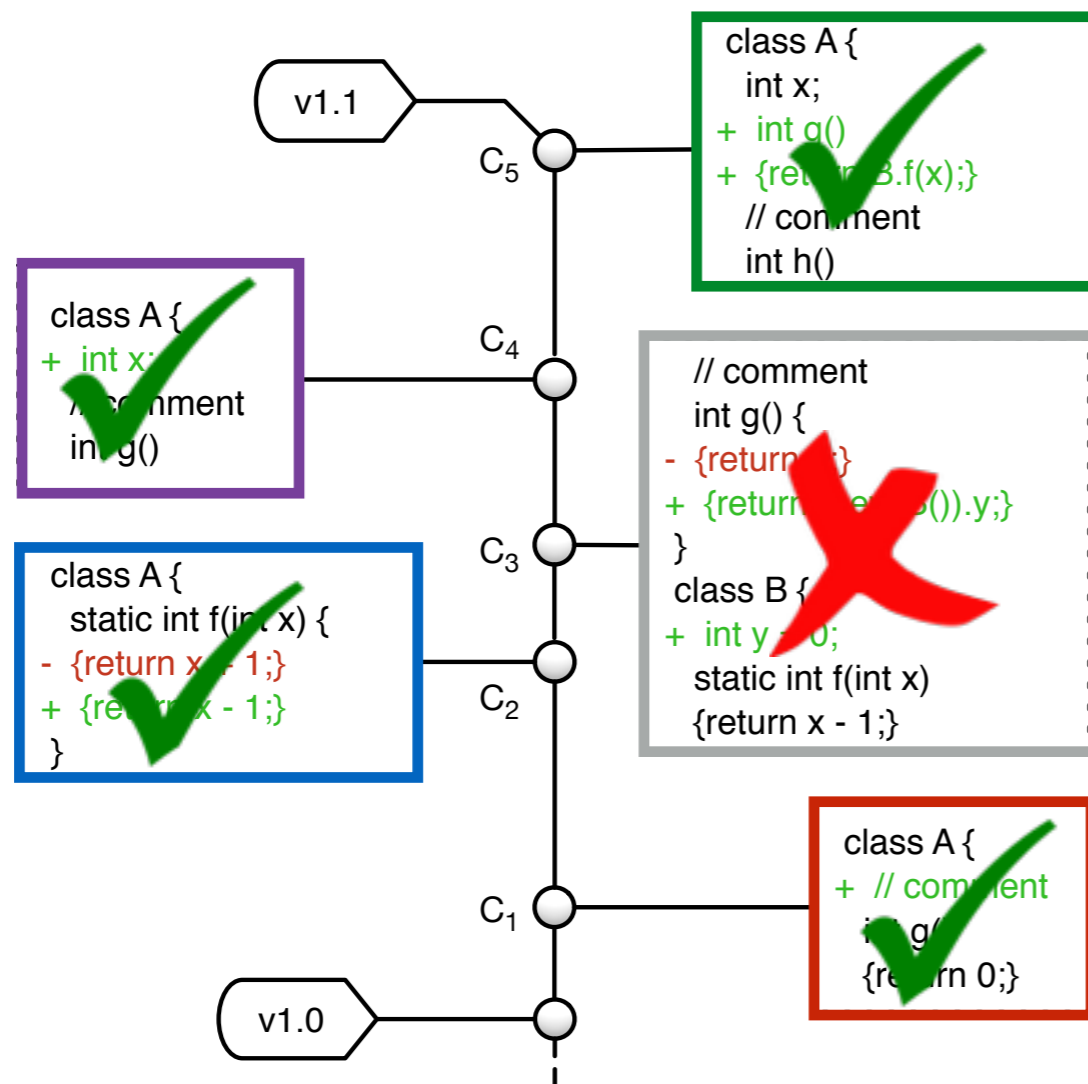
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class TestA {
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  {assertEquals(-1, a.g());}
}
```

Test case:
a.g() == -1

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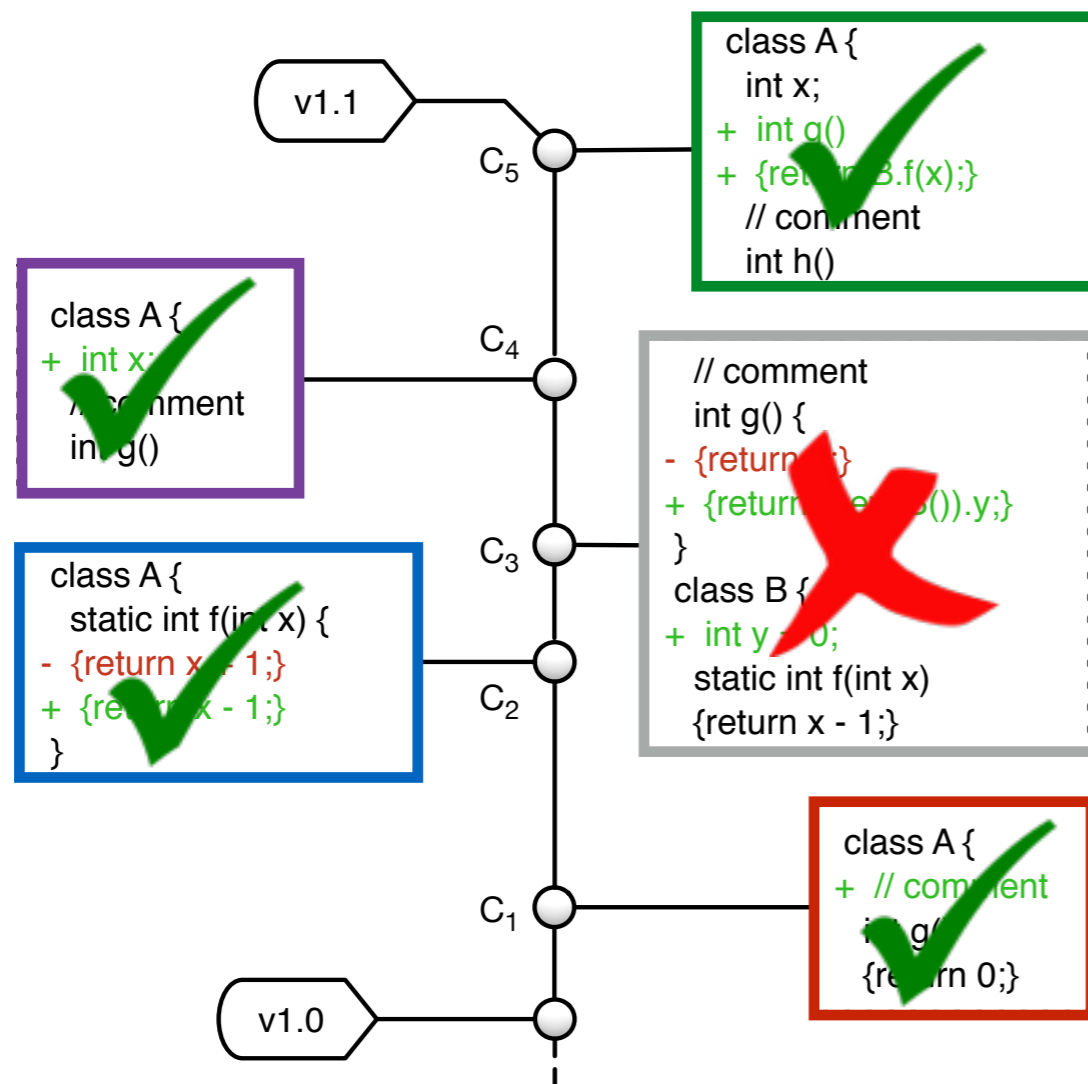


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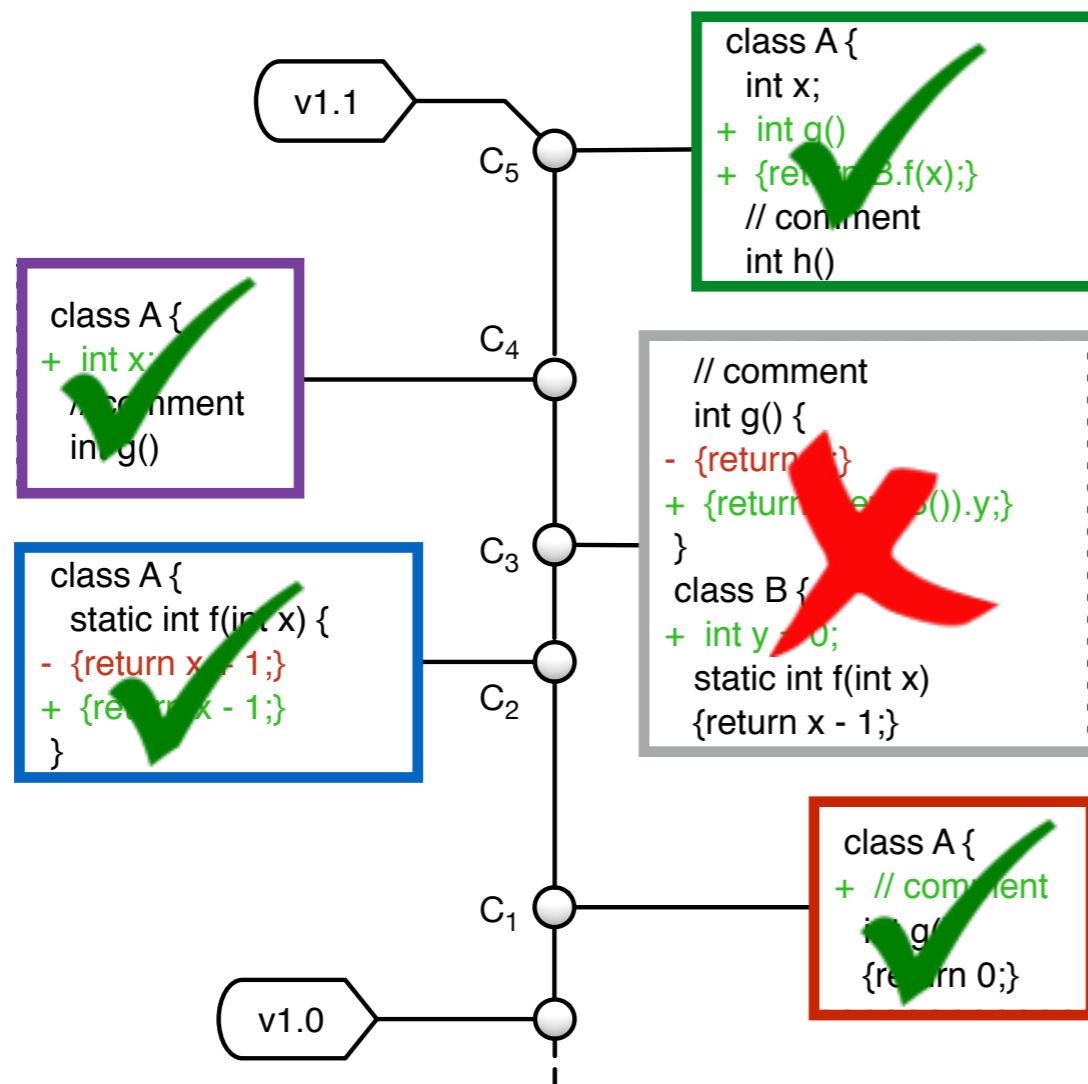


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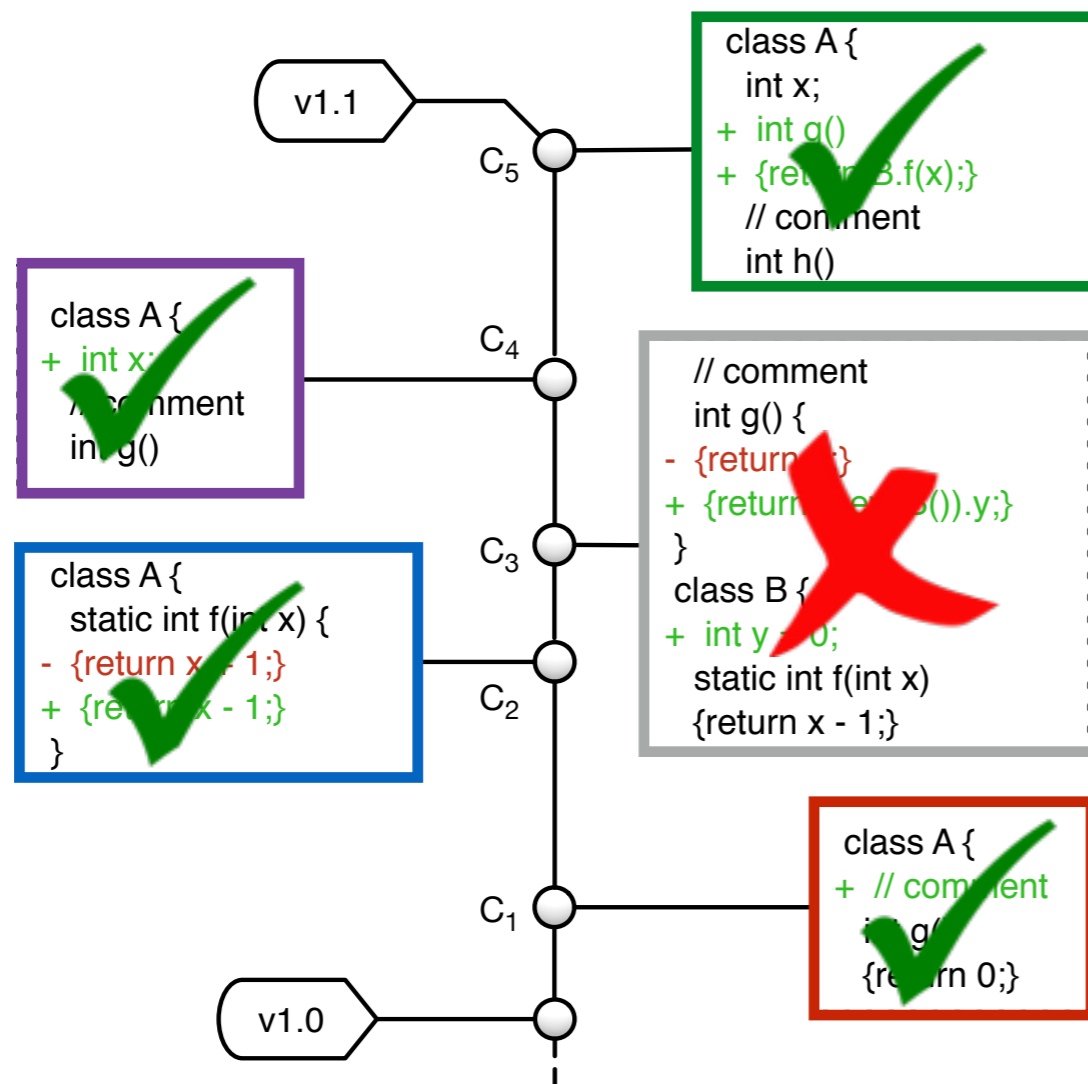


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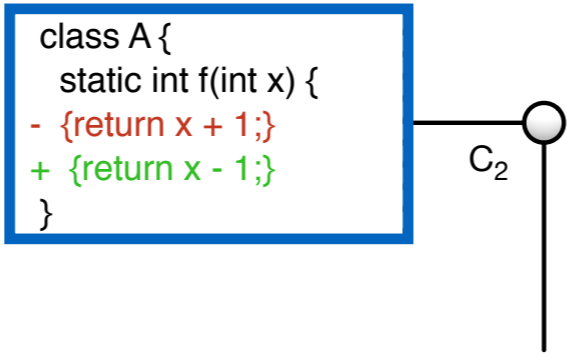
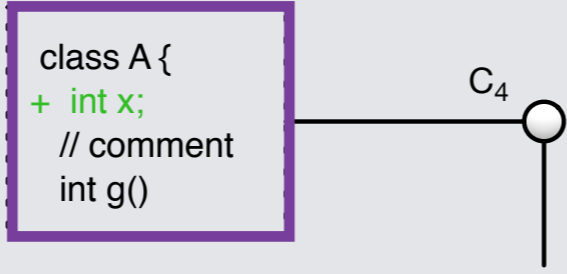
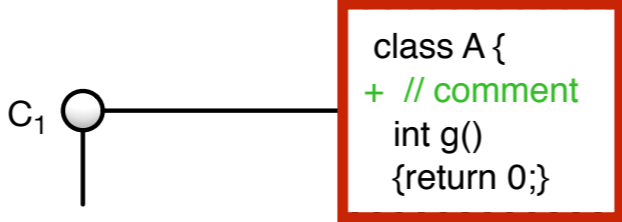
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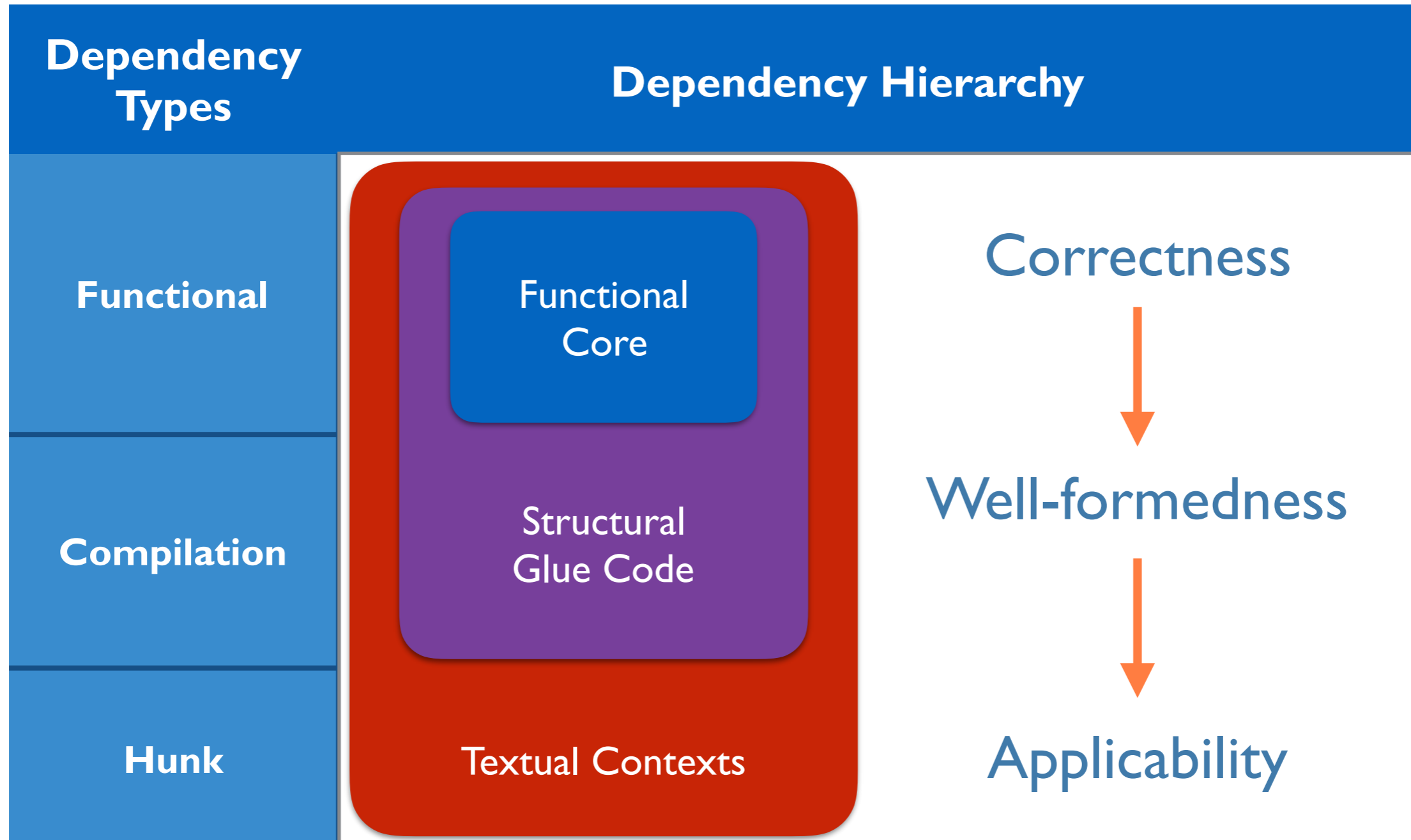
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Dependency Hierarchy

Dependency Types	Examples	Definitions
<p>Functional</p>	 <pre>class A { static int f(int x) { - {return x + 1;} + {return x - 1;} } }</pre>	<p>required for maintaining the semantic behaviours (e.g., pass the same tests)</p>
<p>Compilation</p>	 <pre>class A { + int x; // comment int g() }</pre>	<p>required for maintaining the wellformedness of the program (e.g., free from compilation errors)</p>
<p>Hunk</p>	 <pre>class A { + // comment int g() {return 0;} }</pre>	<p>specific to text-based version control systems (e.g., Git)</p>

Dependency Hierarchy



Outline

1. Introduction
2. Dependency Hierarchy
- 3. CSlicer Algorithm**
- 4. Evaluation**
- 5. Related Work & Conclusion**

CSlicer Overview

Input:

- $H = p_0 \dots p_k$ *well-formed*
- $T = \{t_1, \dots, t_m\}$ *tests for p_k*

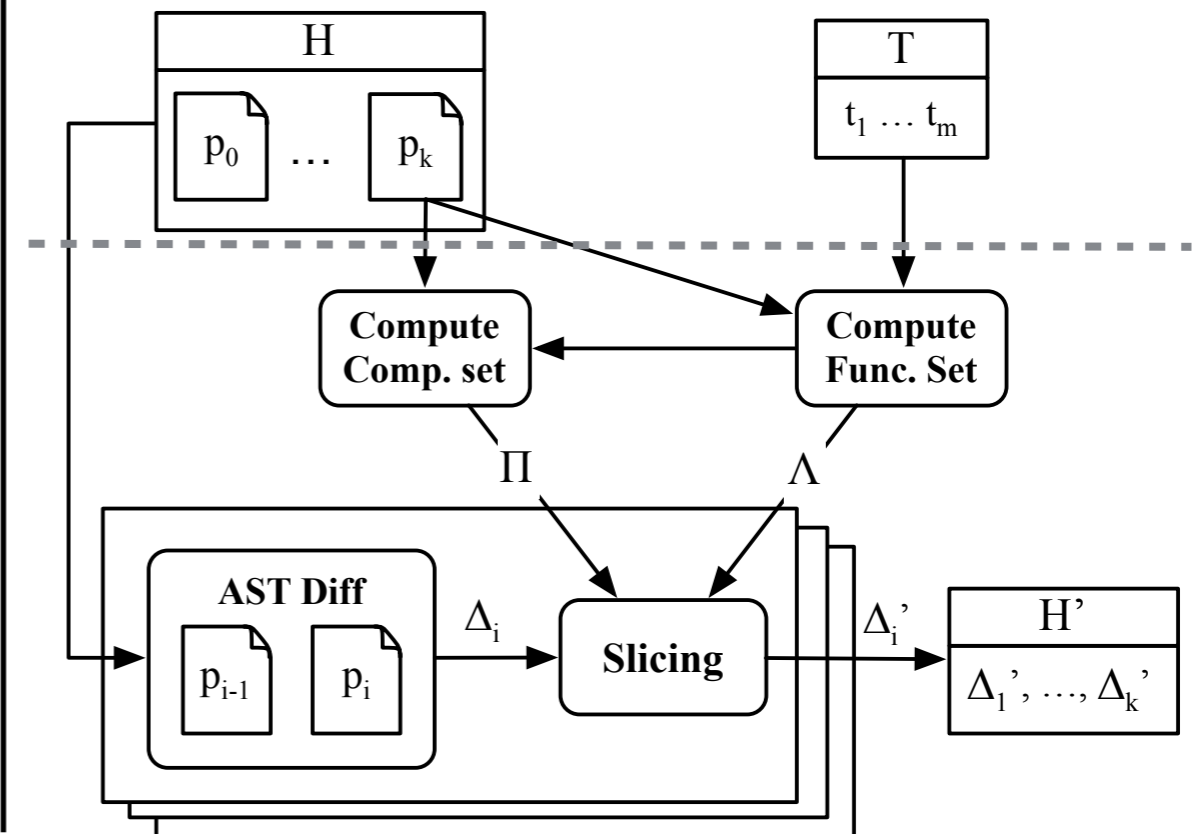
Slicing core:

- **FUNC** set: Λ
- **COMP** set: Π
- **Slicer**(Λ, Π, Δ_i) = Δ_i'

Output:

- $H' = \langle \Delta_1', \dots, \Delta_k' \rangle$

slice



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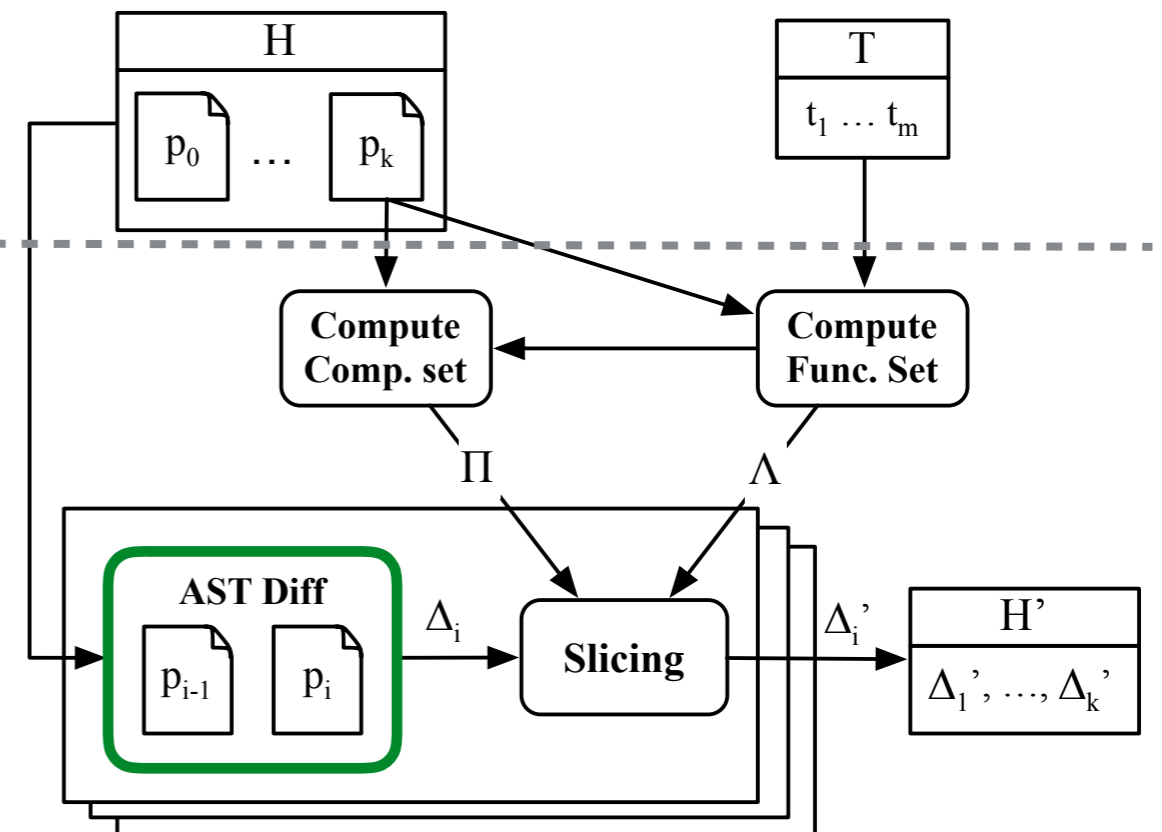
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I. AST differencing



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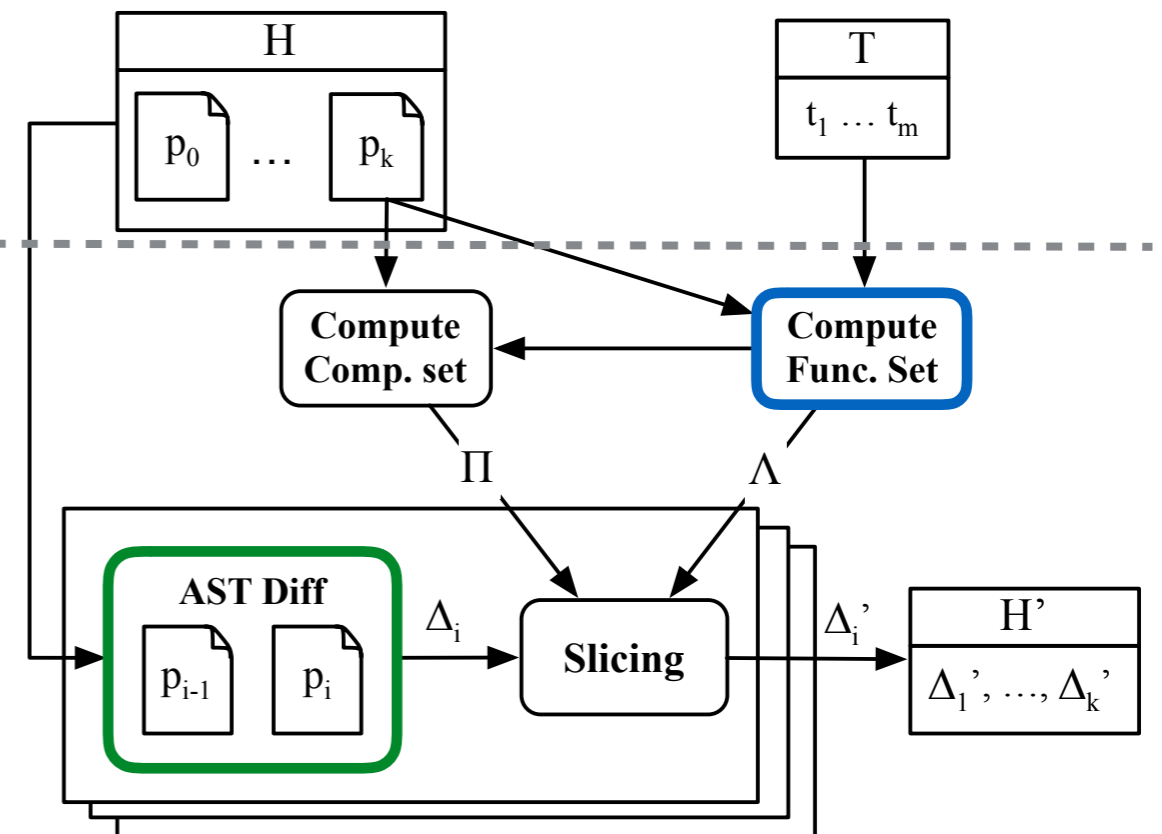
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1. AST differencing
2. Compute Functional set



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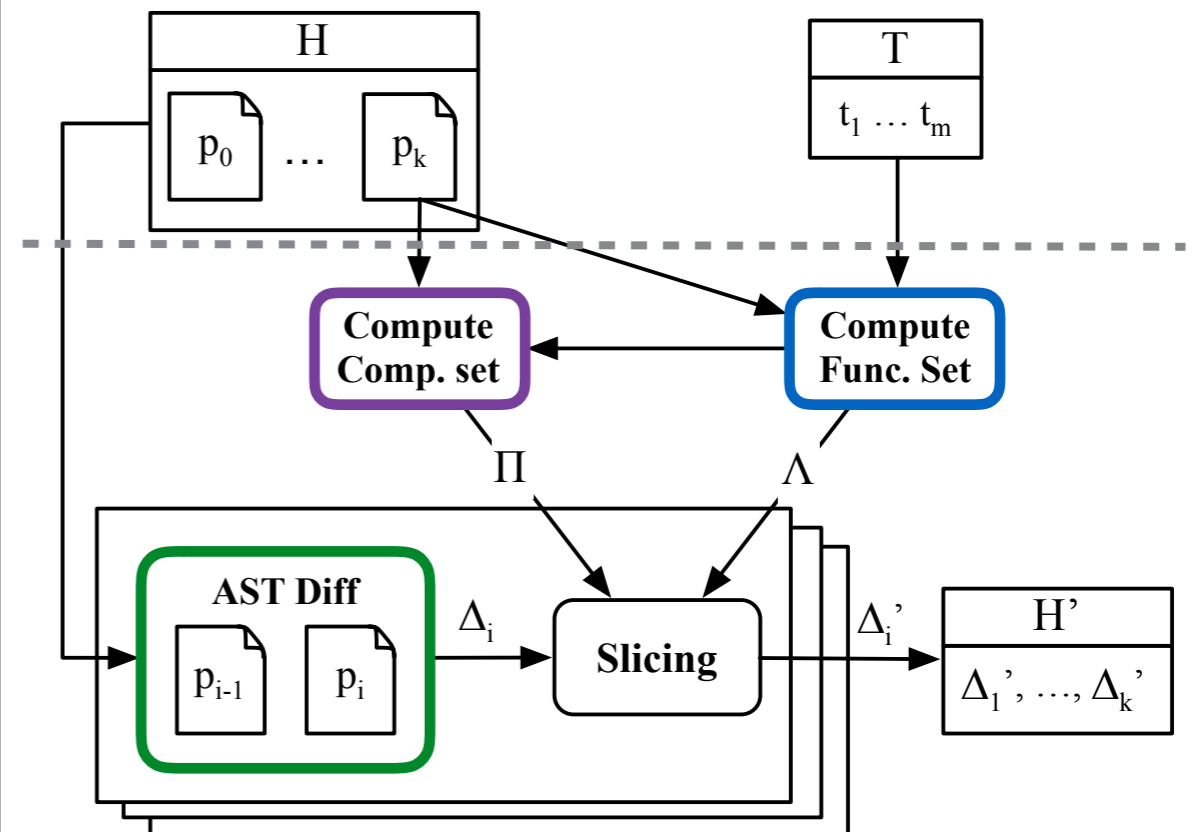
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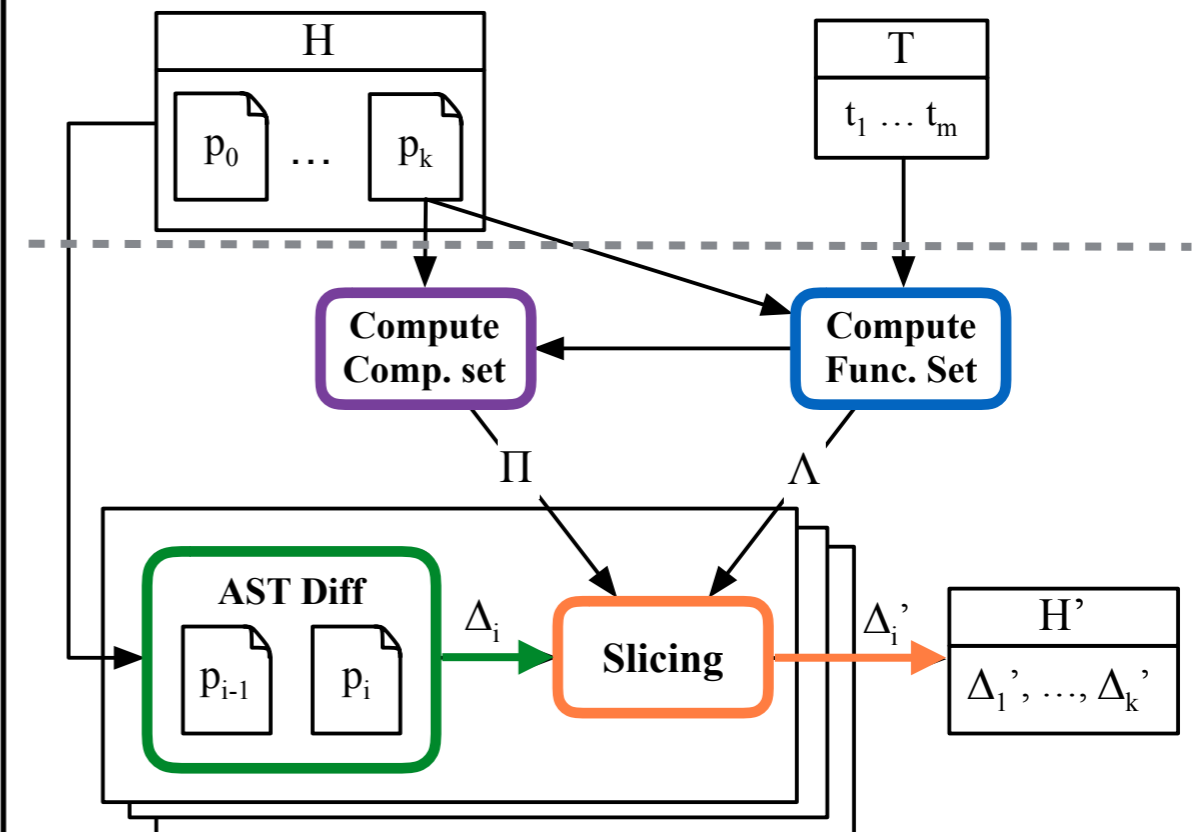
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1. AST differencing
2. Compute Functional set
3. Compute Compilation set
4. Changeset Slicing

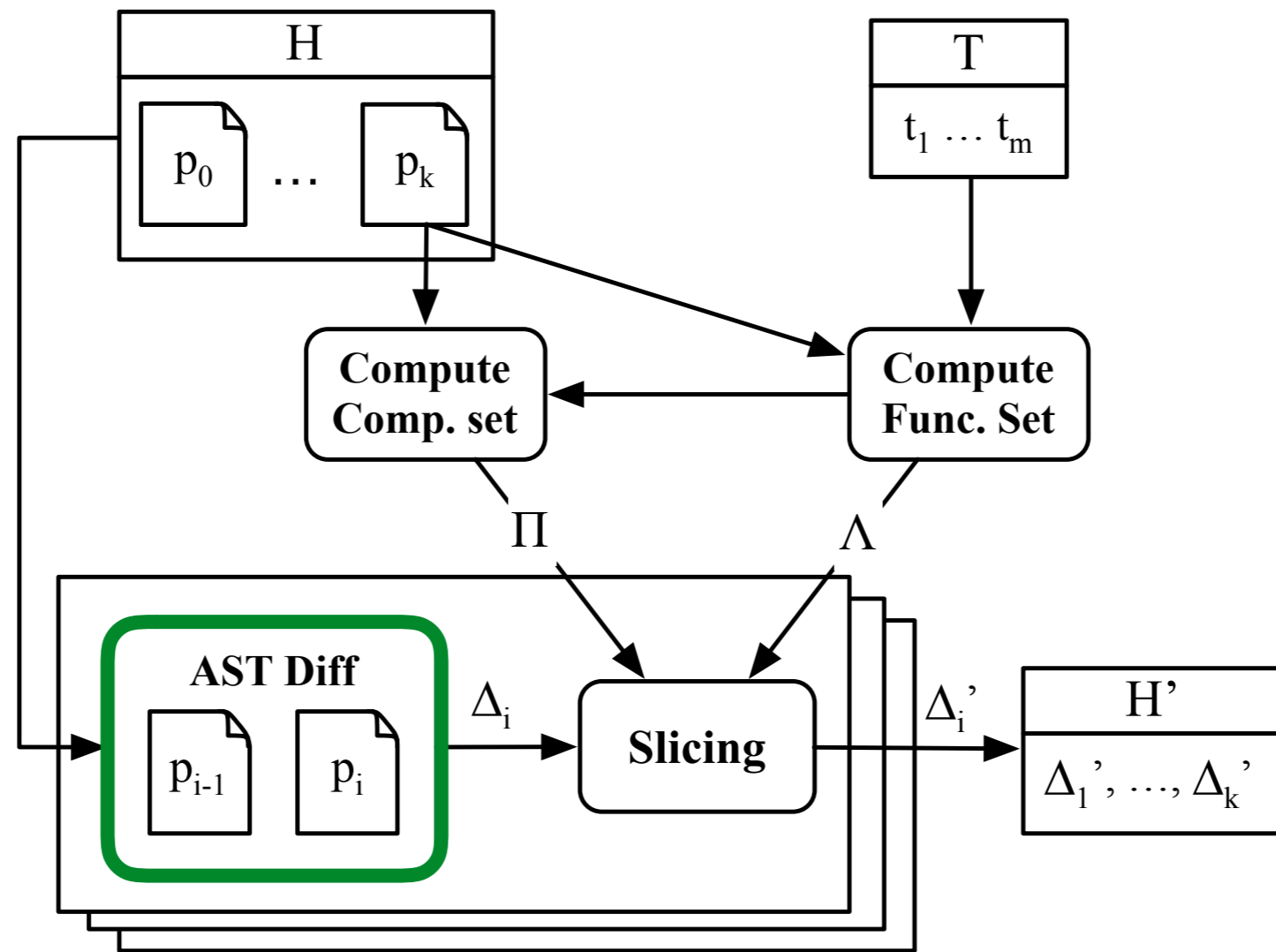


Language Model

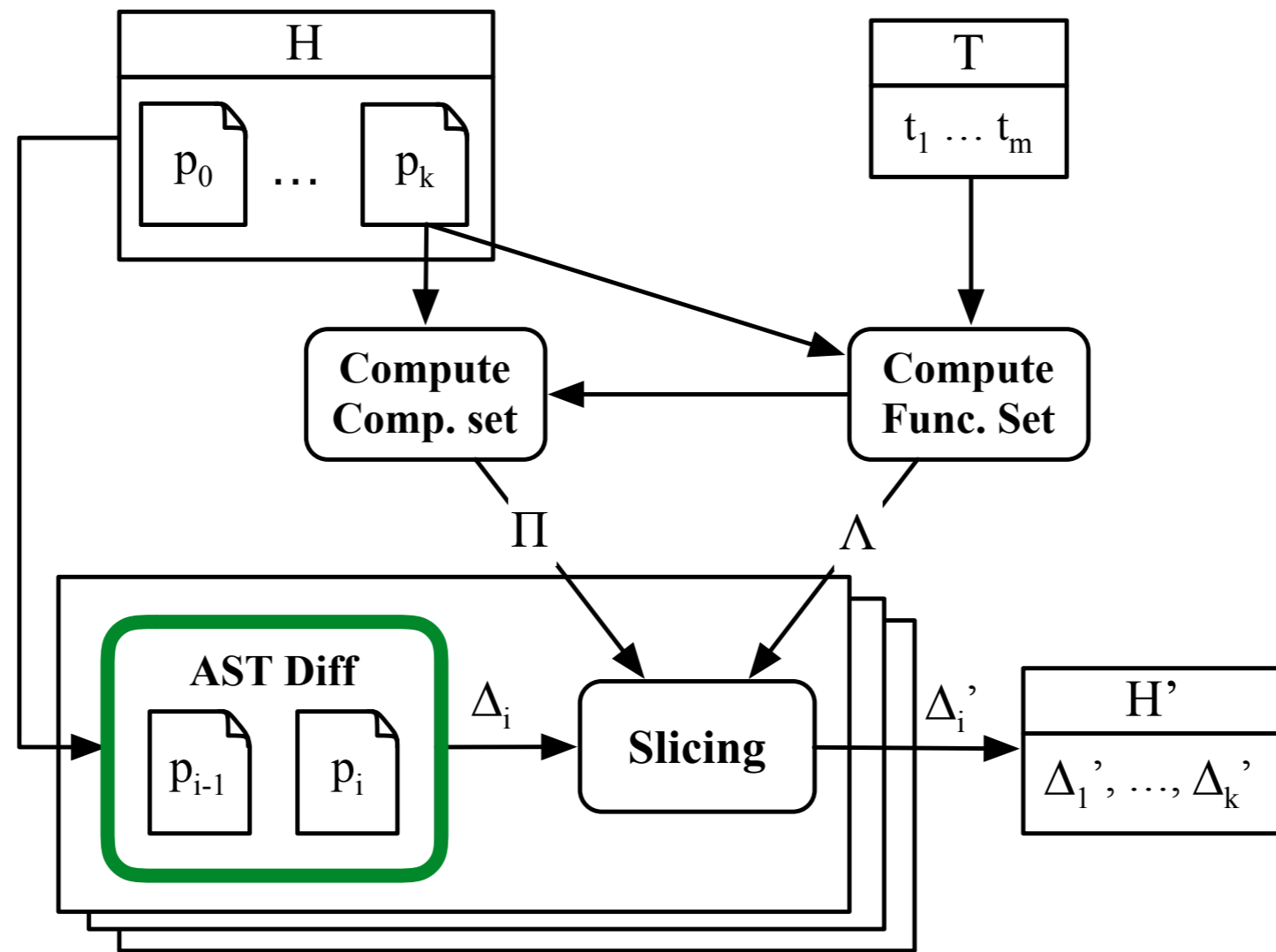
Simplified language model:

- Featherweight Java [[Igarashi et al., ACM TOPLAS'01](#)]
- Core object-oriented features and type system
- No reflection, abstract class, etc.
- Advanced Java features can be handled as algorithmic extensions

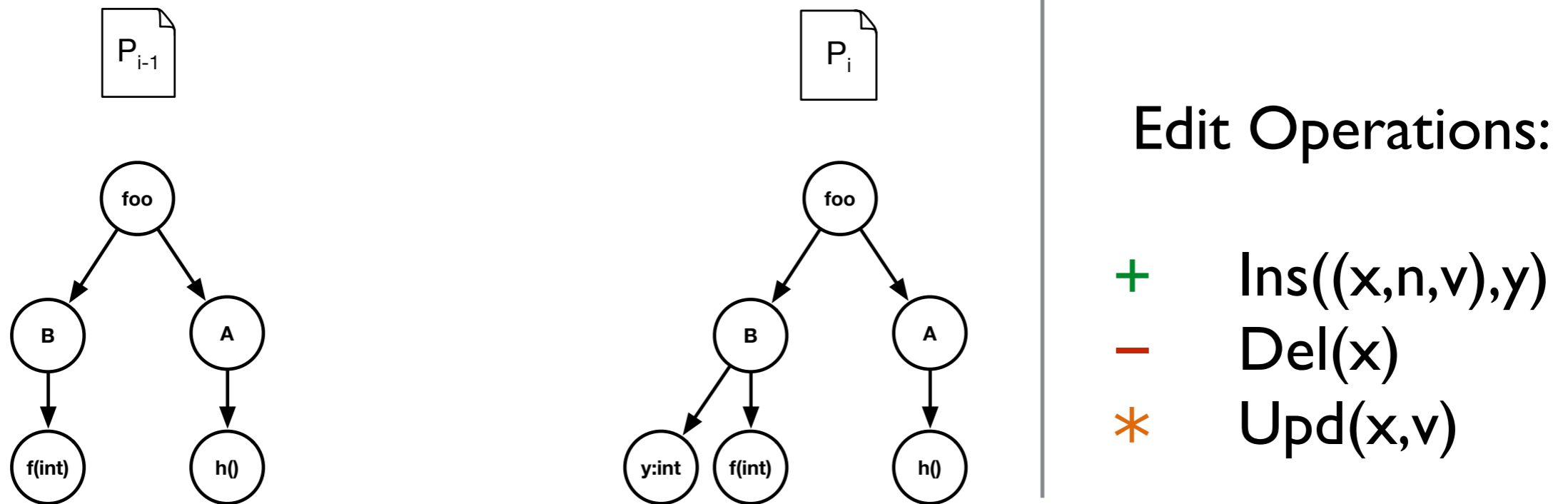
AST Differencing



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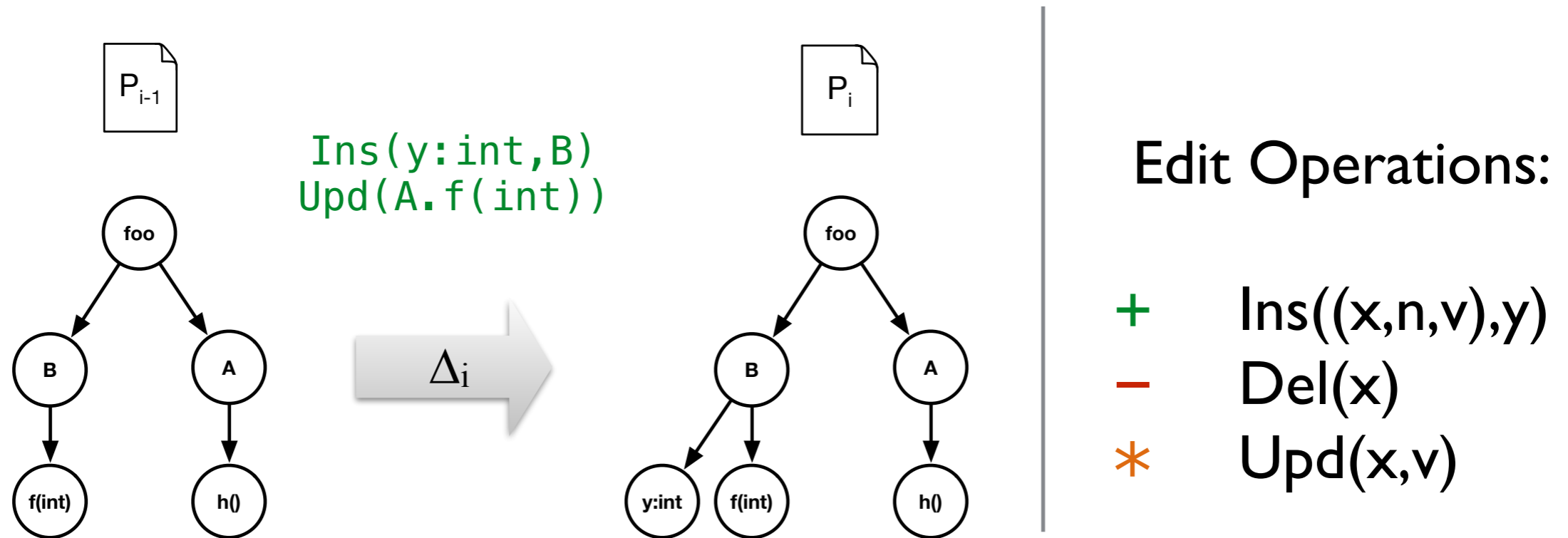
AST Differencing



Compare two abstract syntax trees:

- Ignore cosmetic changes; match on unique names
- Focus on structural nodes (class, method, and field)
- Structural differencing [Fluri et al., IEEE TSE'07]

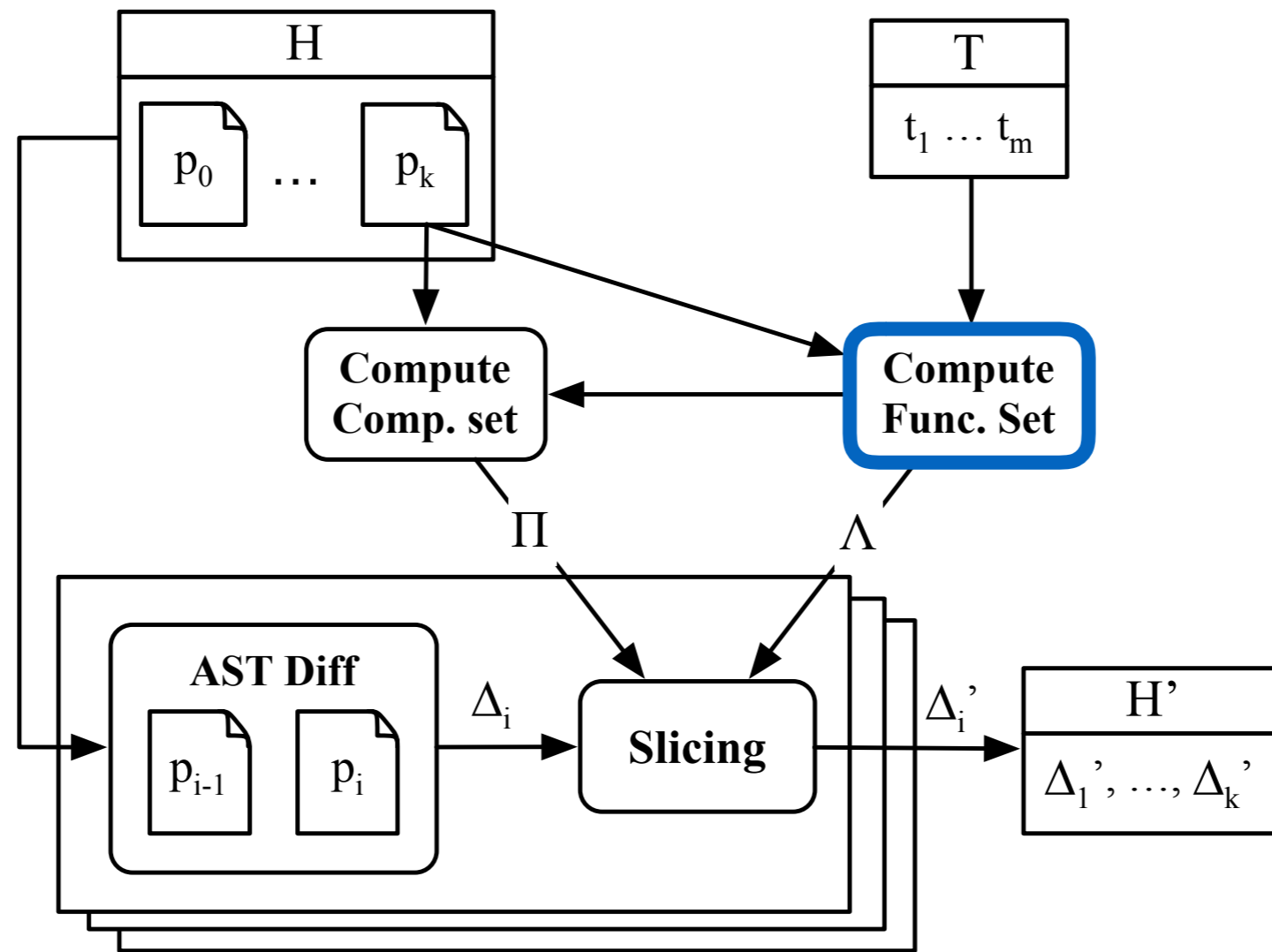
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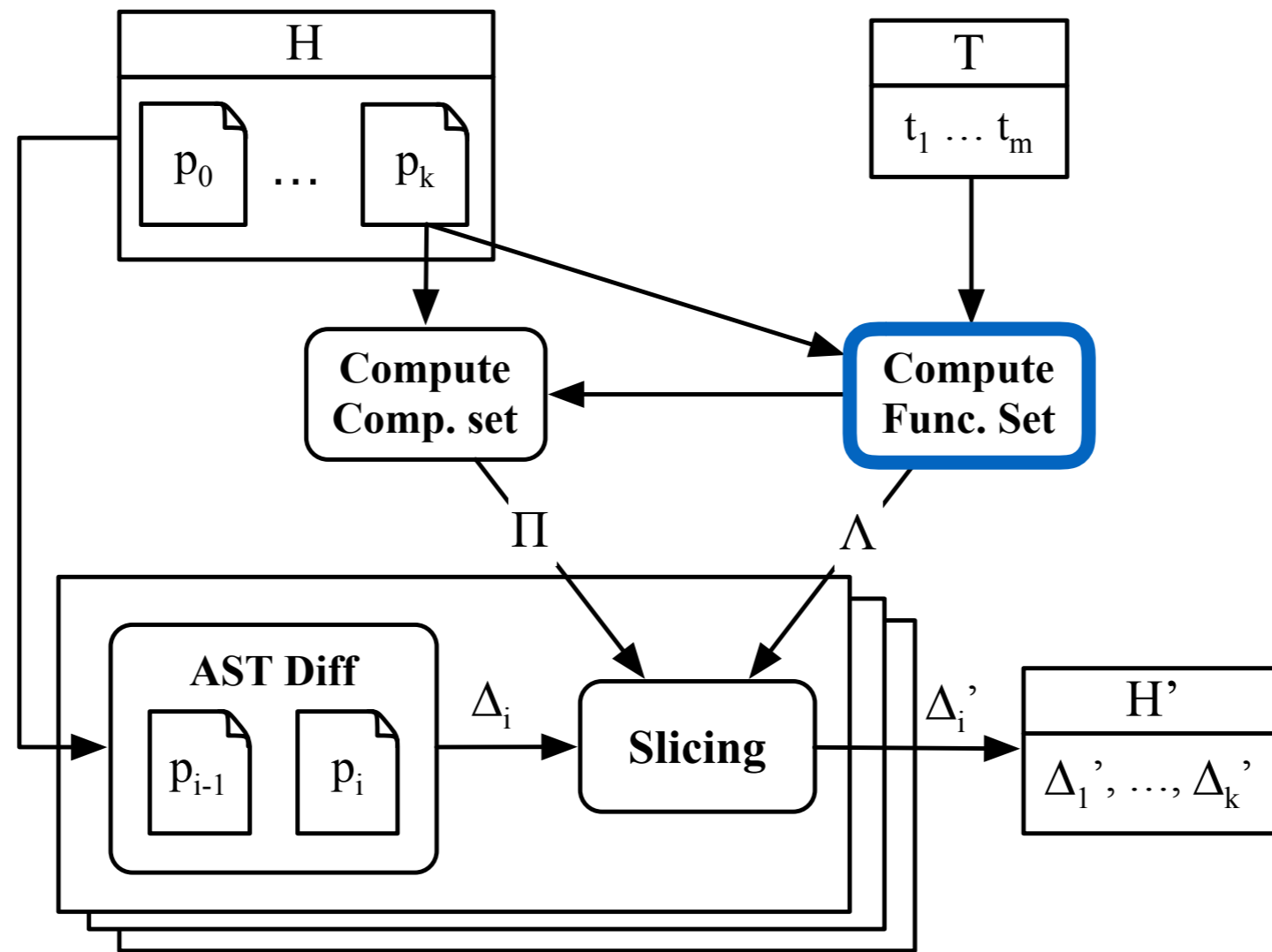
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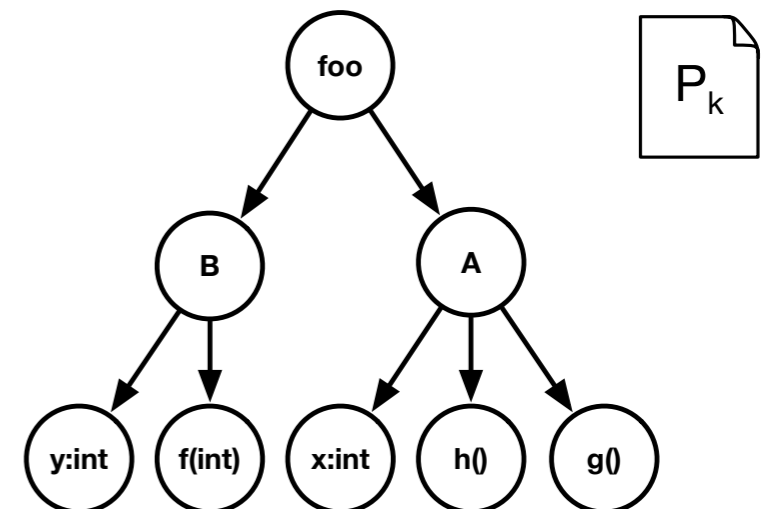


Compute Functional Set

Functional Set:

- Nodes directly traversed during test execution
- Dynamic analysis
- Ensure functional correctness

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class A {  
  int x;  
  int g()  
  {return B.f(x);}  
  // comment  
  int h()  
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  int y = 0;  
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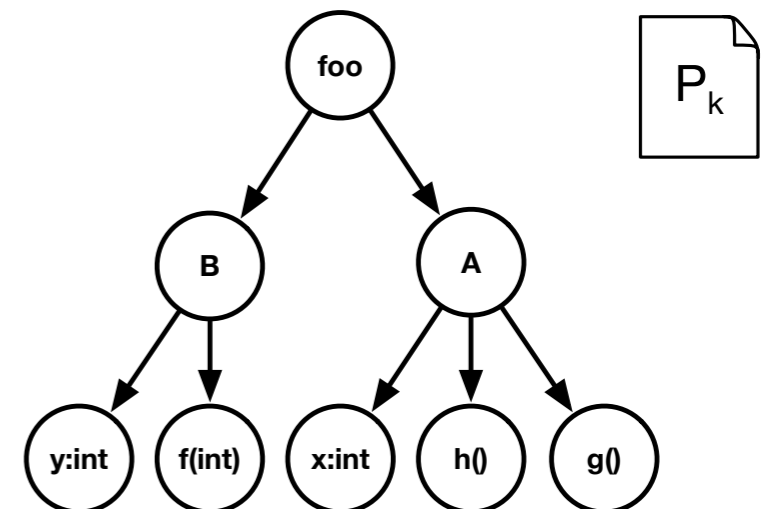
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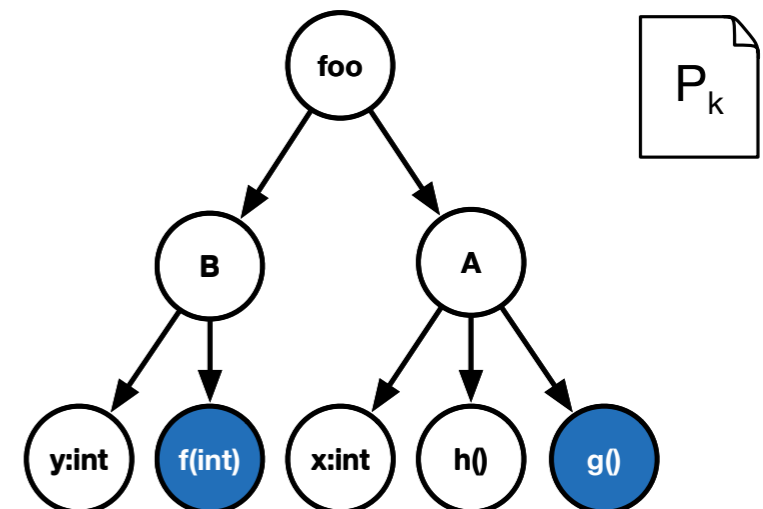
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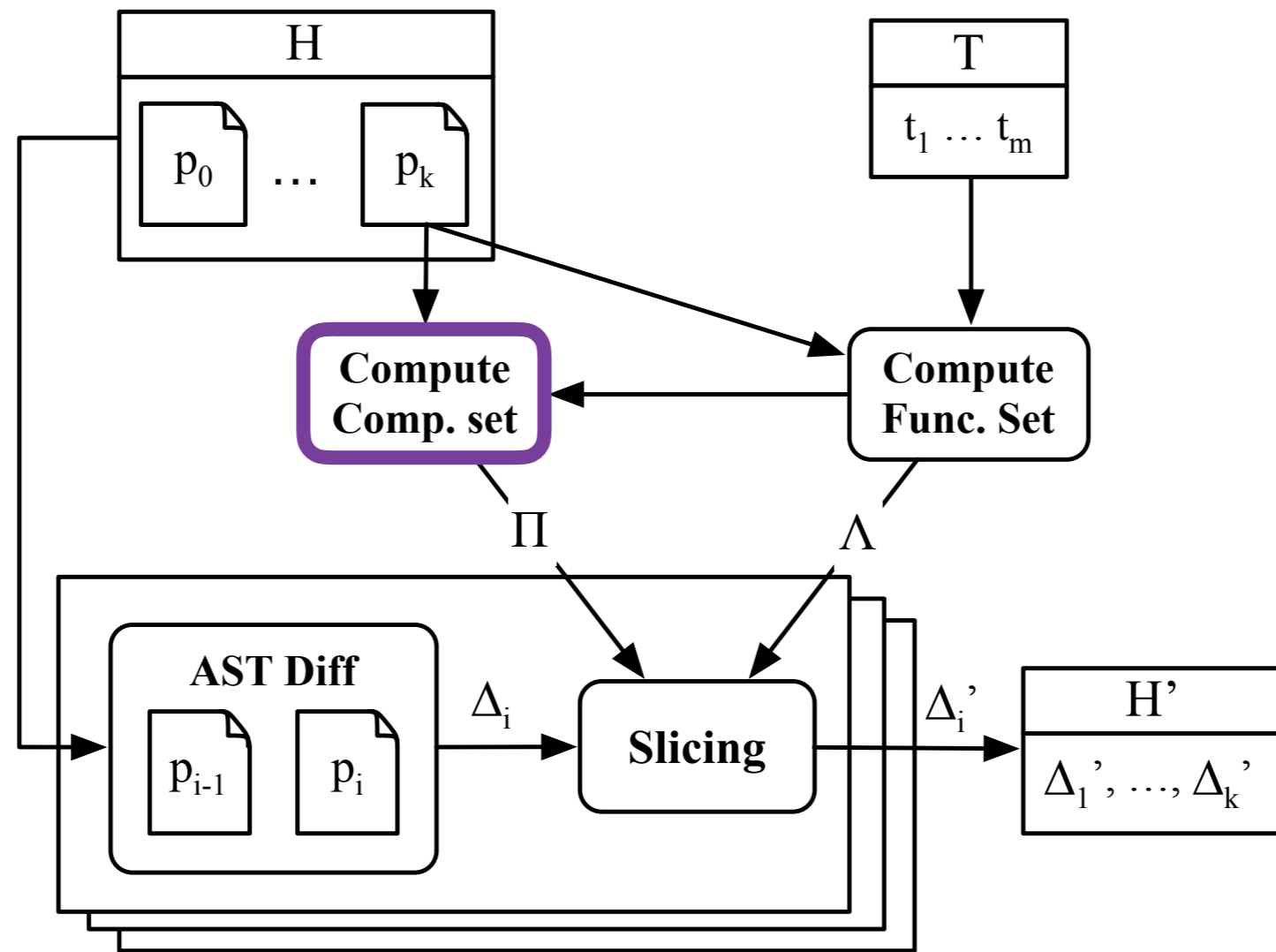
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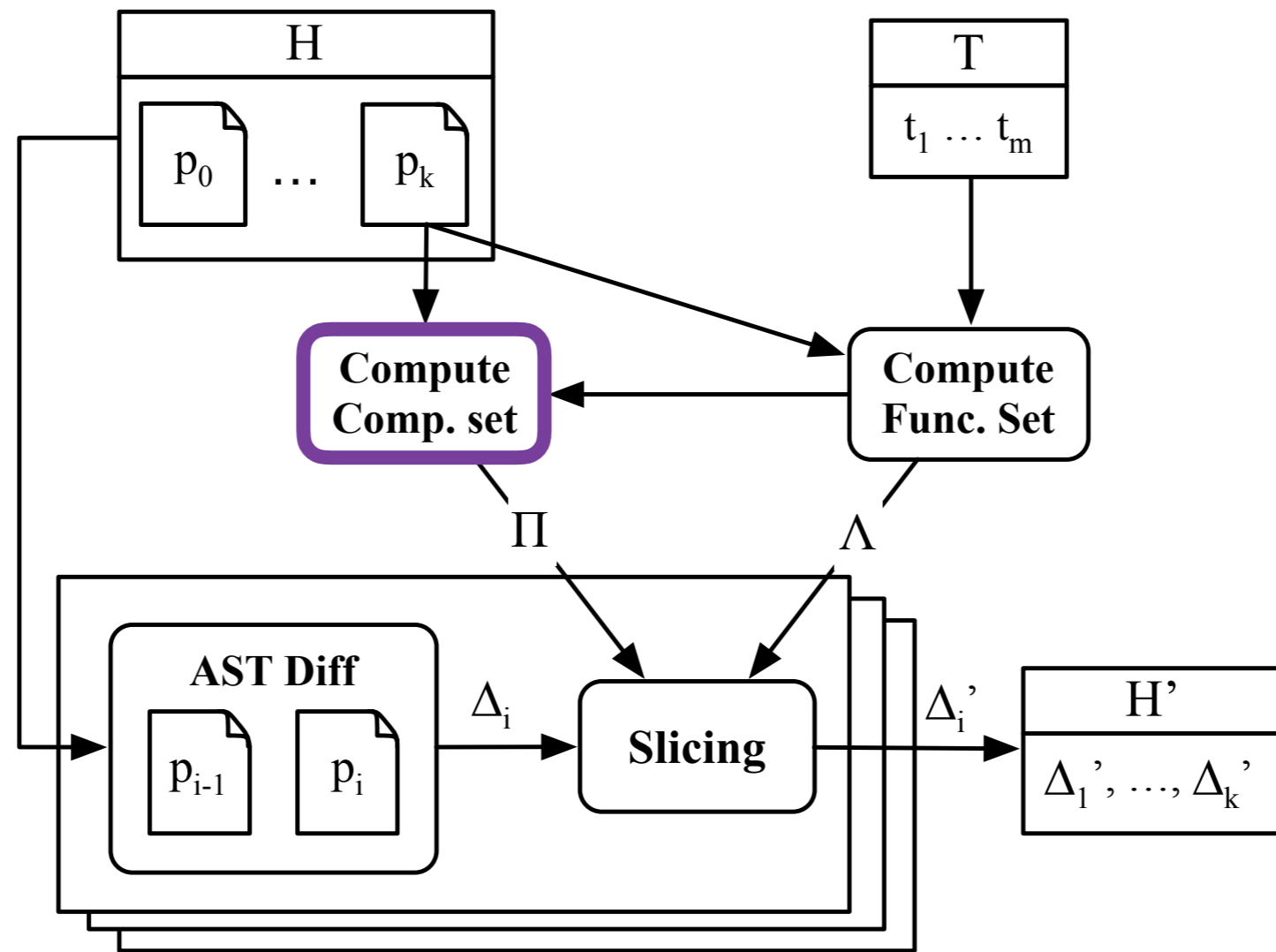
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Compute Compilation Set



Compute Compilation Set



Compute Compilation Set

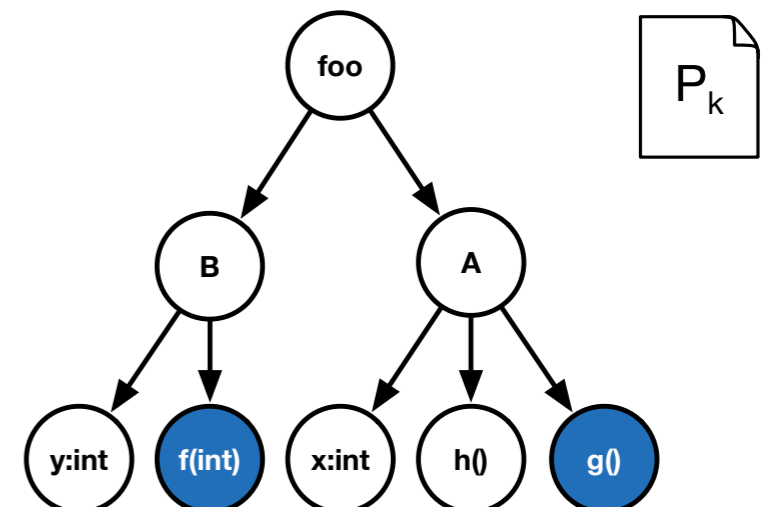
Compilation Set:

- Nodes referenced by the functional set
- Static analysis
- Ensure type safety

Inference Rules:

- Enclosing classes should exist
- Accessed fields should exist
- etc.

```
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  int x;  
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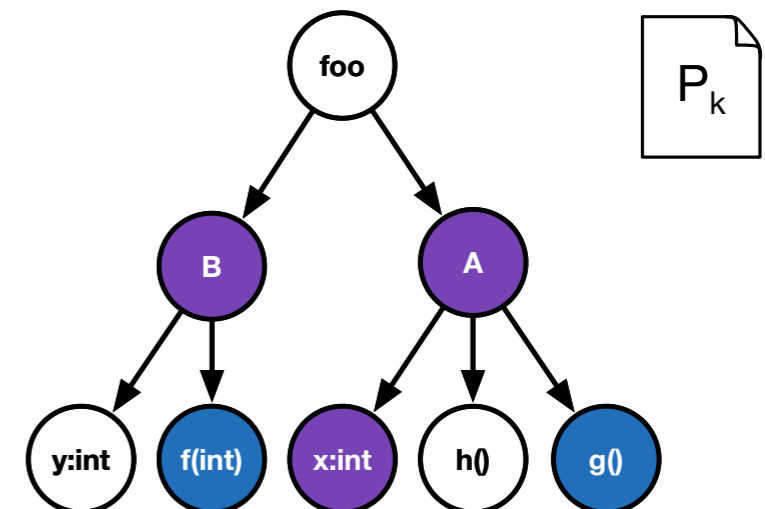
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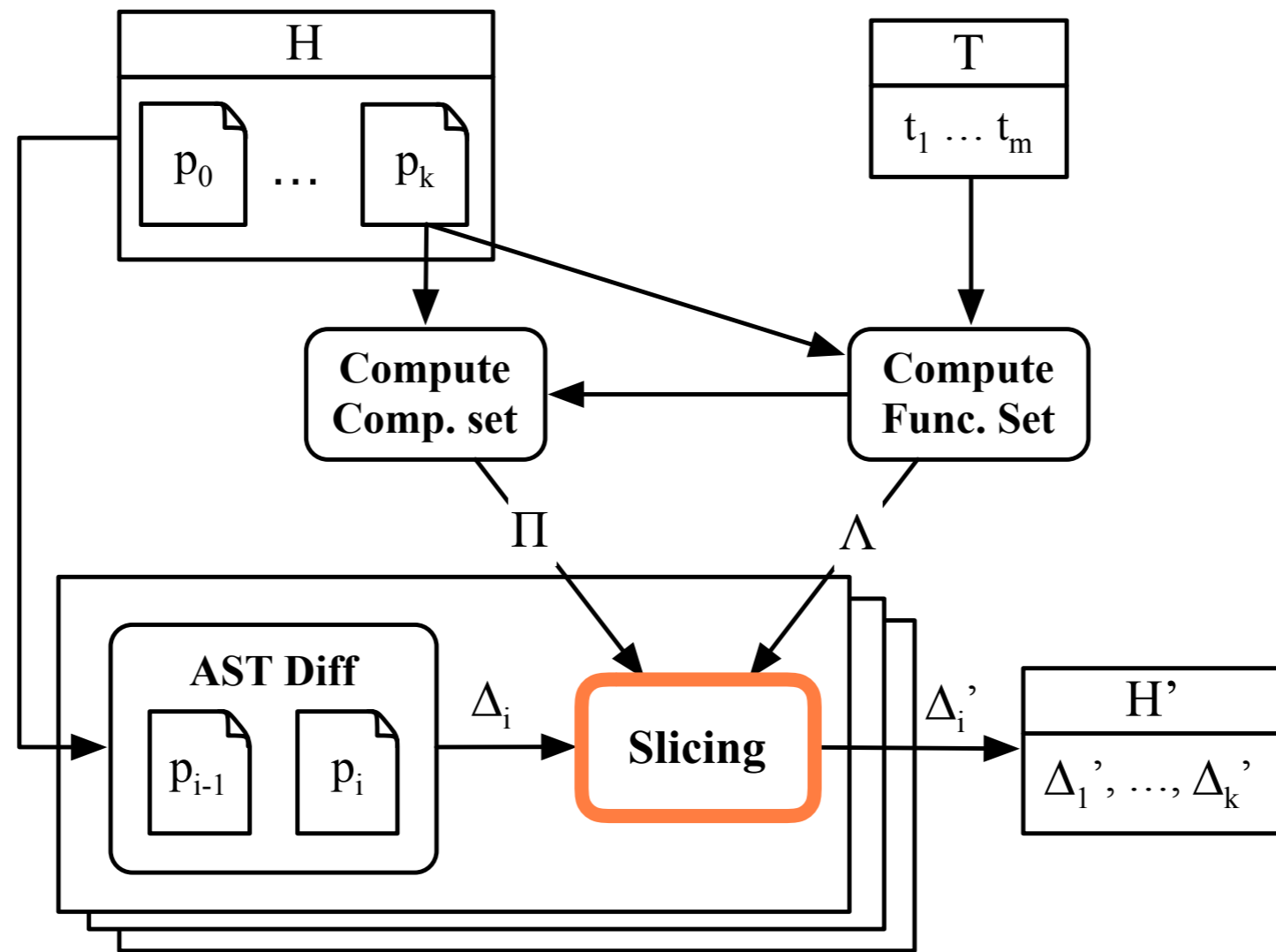
Compute Compilation Set

$$\begin{array}{c}
 \frac{C <: D \quad C \in \Pi}{D \in \Pi} \text{ [L.1]} \quad \frac{f : C \in \Pi}{C \in \Pi} \text{ [L.2]} \quad \frac{C(\overline{D \ f})\{\text{super}(\overline{f}); \overline{\text{this.f}} = \overline{f};\} \in \Pi}{C \in \Pi \quad \overline{D} \in \Pi \quad \overline{f} \in \Pi} \text{ [K1]} \\
 \\
 \frac{C \ m(\overline{D \ x})\{\text{return } e; \} \in \Pi}{C \in \Pi \quad \overline{D} \in \Pi} \text{ [M1]} \quad \frac{\dots\{\text{return } e.f; \} \in \Pi}{f \in \Pi} \text{ [E1]} \quad \frac{\dots\{\text{return } e.m(\overline{e}); \} \in \Pi}{m \in \Pi} \text{ [E2]} \\
 \\
 \frac{\dots\{\text{return new } C(\overline{e}); \} \in \Pi}{C \in \Pi} \text{ [E3]} \quad \frac{\dots\{\text{return } (C)e; \} \in \Pi}{C \in \Pi} \text{ [E4]} \quad \frac{x \in \Pi}{\text{PARENT}(x) \in \Pi} \text{ [P1]} \quad \frac{x \in \Lambda}{x \in \Pi} \text{ [T1]}
 \end{array}$$

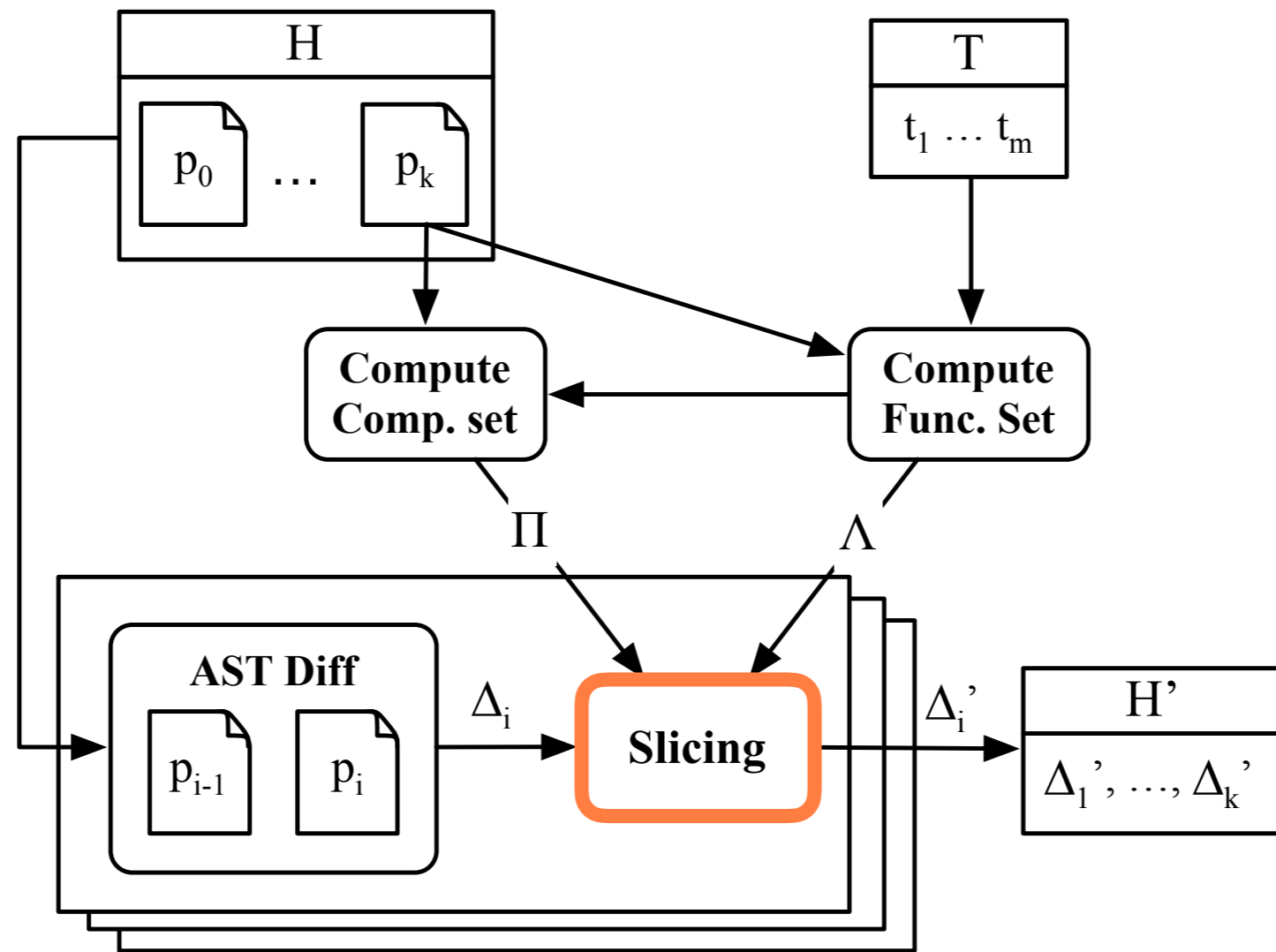
Inference Rules:

- Based on [\[Kastner & Apel, ASE'08\]](#)
- Tailored for method-field level granularity
- Complete for our language model

Changeset Slicing



Changeset Slicing



Changeset Slicing

	// comment	B.f(int)	B.y	A.h()	A.x	A.g()
C5						+
C4					+	
C3			+	*		
C2		*				
C1	+					

Functional

Compilation

+ Ins

- Del

* Upd

Change Matrix: maps *atomic changes* to commits

- Cells are marked by change types
- Atomic changes are color coded

Changeset Slicing

Functional

Compilation

+ Ins - Del * Upd

General Slicing Rules:

- Keep **blue** cells
- Keep **purple** +, -
- Drop white - unless affecting method lookup

	δ_1	δ_2	δ_3	δ_4	δ_5
C5				*	
C4	*	*		+	
C3			-		+
C2				-	
C1	+				

Changeset Slicing



Side-effects (Git):

- Keeping original commit
- Dependencies between white cells
- Detection and resolution

	δ_1	δ_2	δ_3	δ_4	δ_5
C5				*	
C4	*	*		+	
C3			-		+
C2				-	
C1	+				

Changeset Slicing



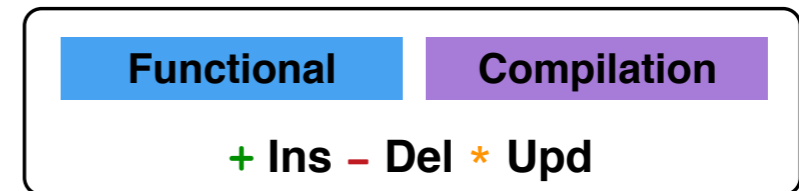
Side-effects (Git):

- Keeping original commit
- Dependencies between white cells
- Detection and resolution



	δ_1	δ_2	δ_3	δ_4	δ_5
C5				*	
C4	*	*		+	
C3			-		+
C2				-	
C1	+				

Changeset Slicing



Side-effects (Git):

- Keeping original commit
- Dependencies between white cells
- Detection and resolution



	δ_1	δ_2	δ_3	δ_4	δ_5
C5				*	
C4	*	*		+	
C3			-		+
C2				-	
C1	+				

Diagram details: A purple arrow points from the orange '*' in C5, δ_4 to the orange '*' in C4, δ_2 . A yellow box highlights the orange '*' in C4, δ_2 . A red 'X' is over the green '+' in C3, δ_5 .

Changeset Slicing

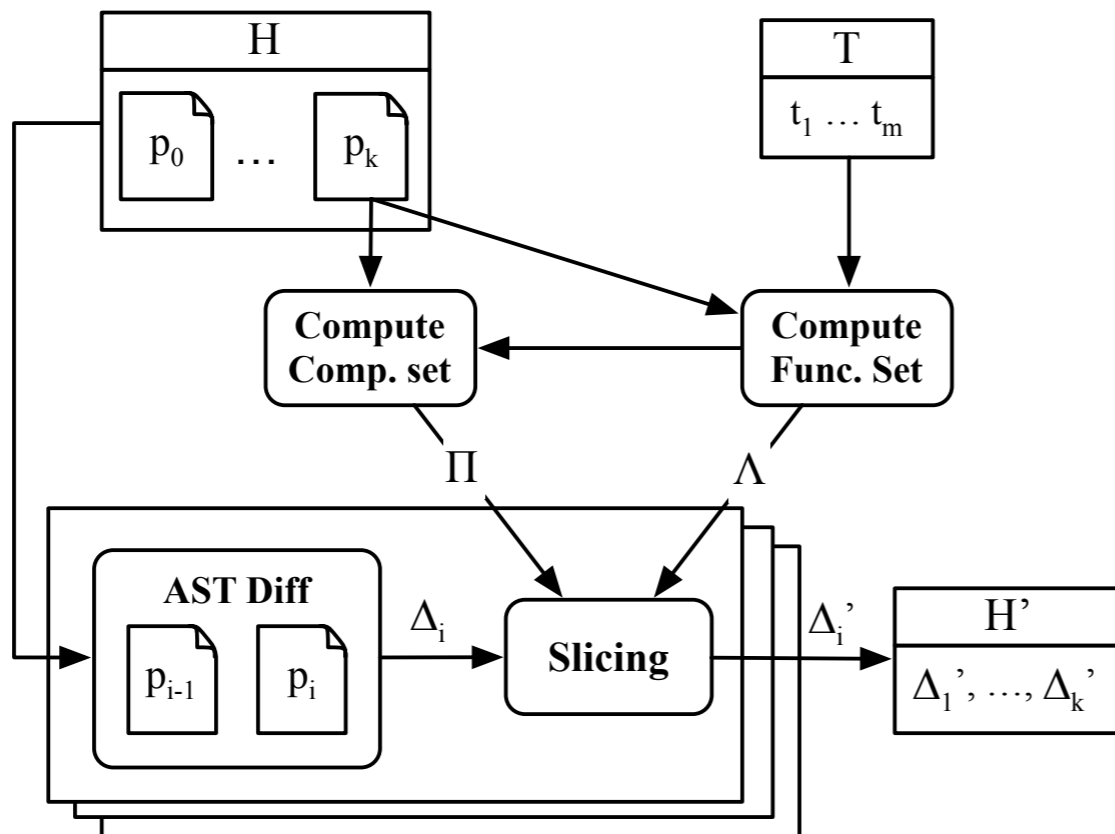
Side-effects (Git):

- Keeping original commit
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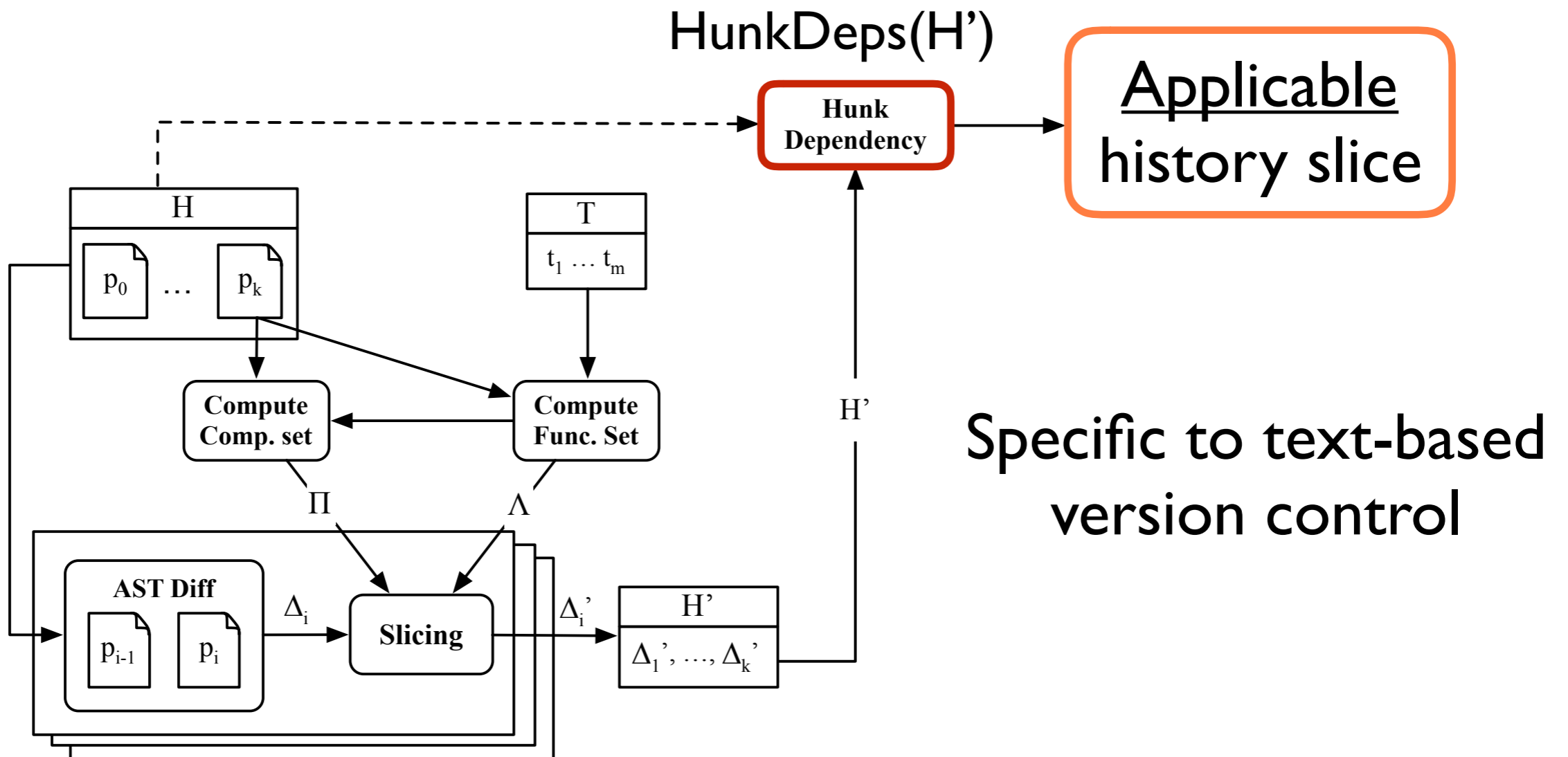


	δ_1	δ_2	δ_3	δ_4	δ_5
CN		X			
C5				*	
C4	*	X		+	
C3			-		+
C2				-	
C1	+				

Hunk Dependency



Hunk Dependency



Outline

1. Introduction
2. Dependency Hierarchy
3. CSlicer Algorithm
- 4. Evaluation**
- 5. Related Work & Conclusion**

Evaluation

Research questions

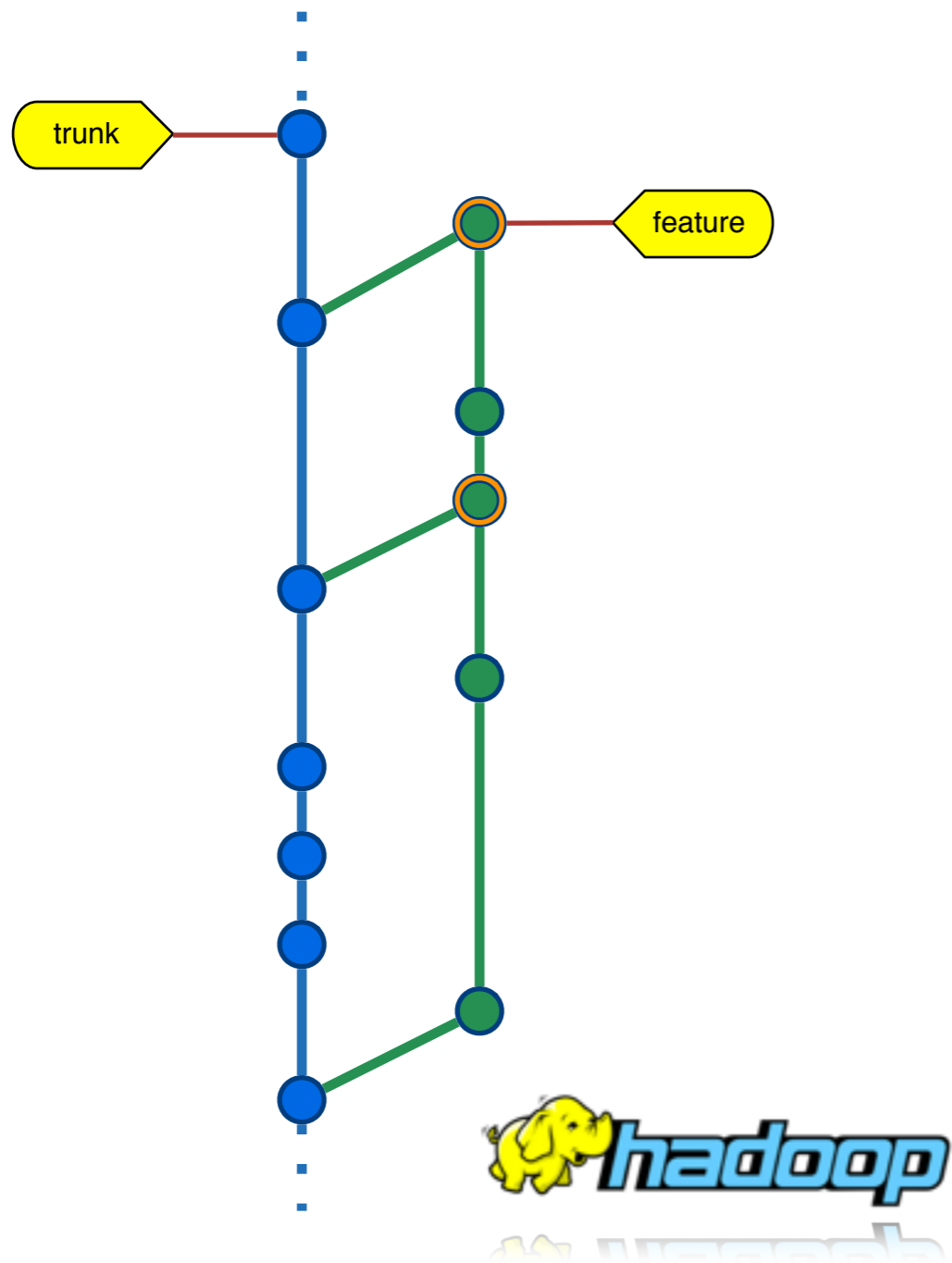
- Accuracy: do we find what we want?
- Effectiveness: reduction rate?
- Efficiency: performance w.r.t. project scale & history length?

Subjects

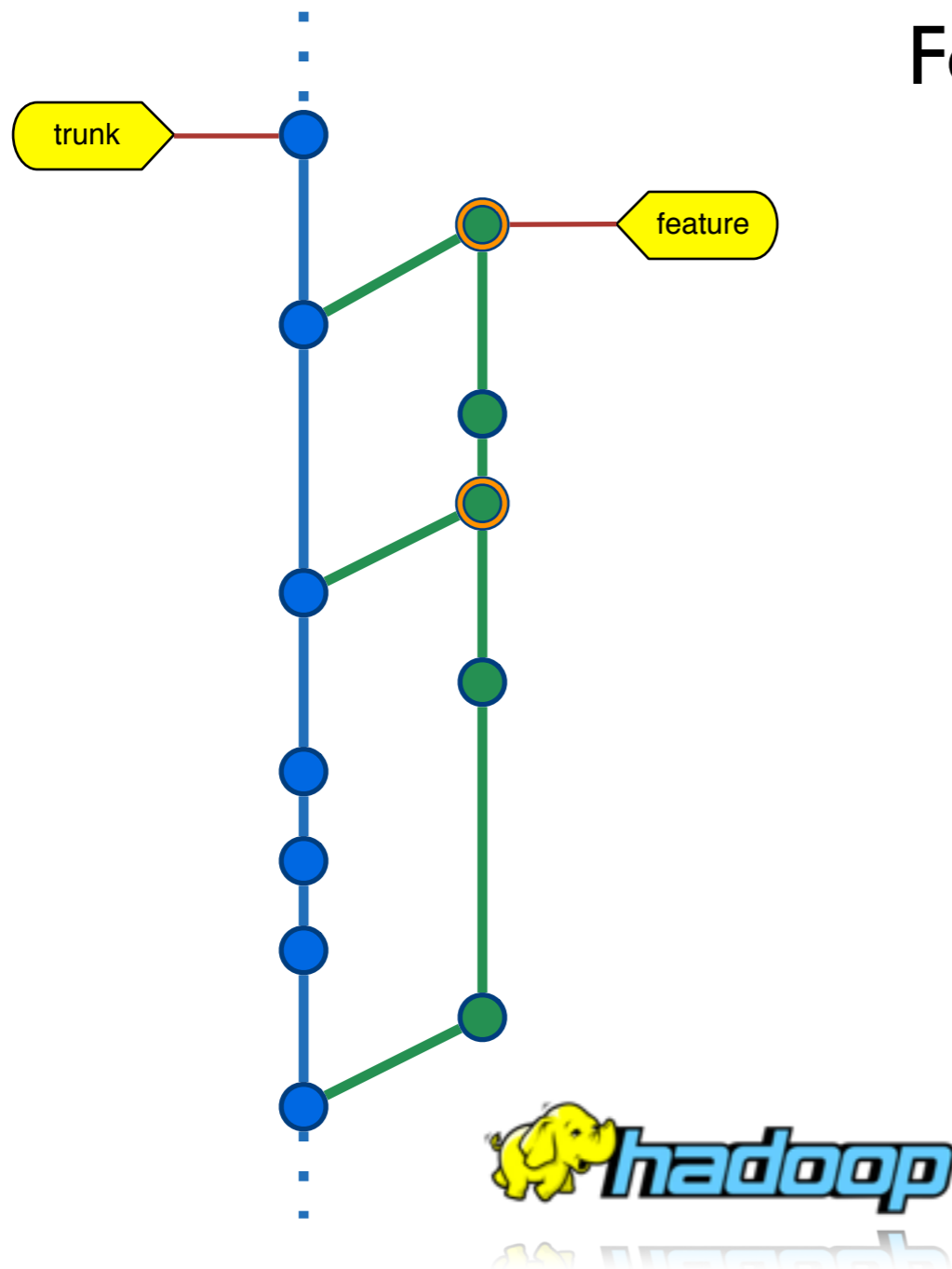
- Advanced Java features not tested: abstract class, reflection, etc.
- Non-Java changes are included by default

Project	# Java Files	LOC	# Authors
Hadoop	5,861	1,291K	169
Elasticsearch	3,865	616K	649
Maven	1,048	142K	78
CSlicer	141	18K	2

Accuracy



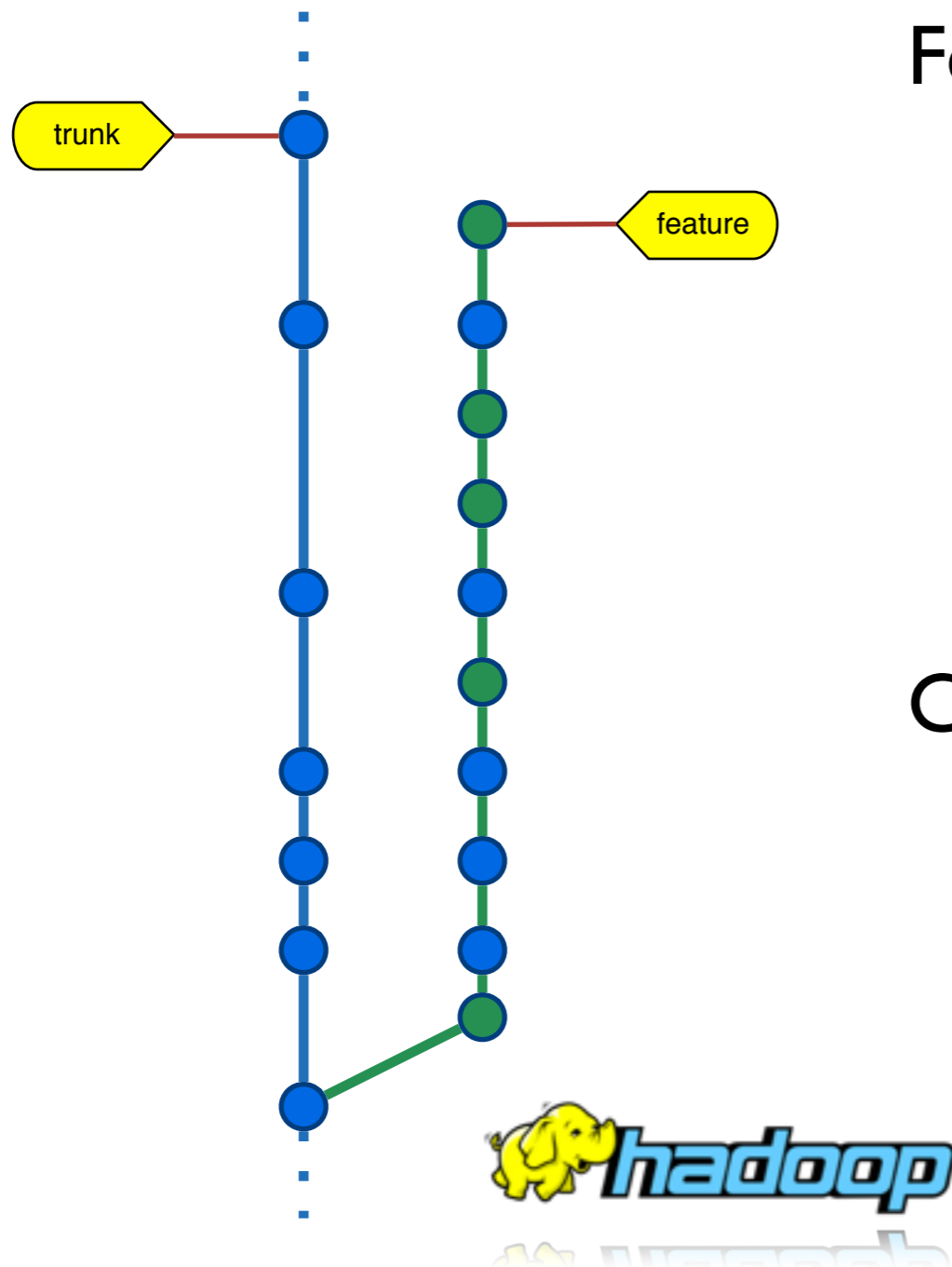
Accuracy



Feature branch

- Merges with the main branch periodically
- 42 feature commits + 47 merges
- 58 accompanied test cases

Accuracy



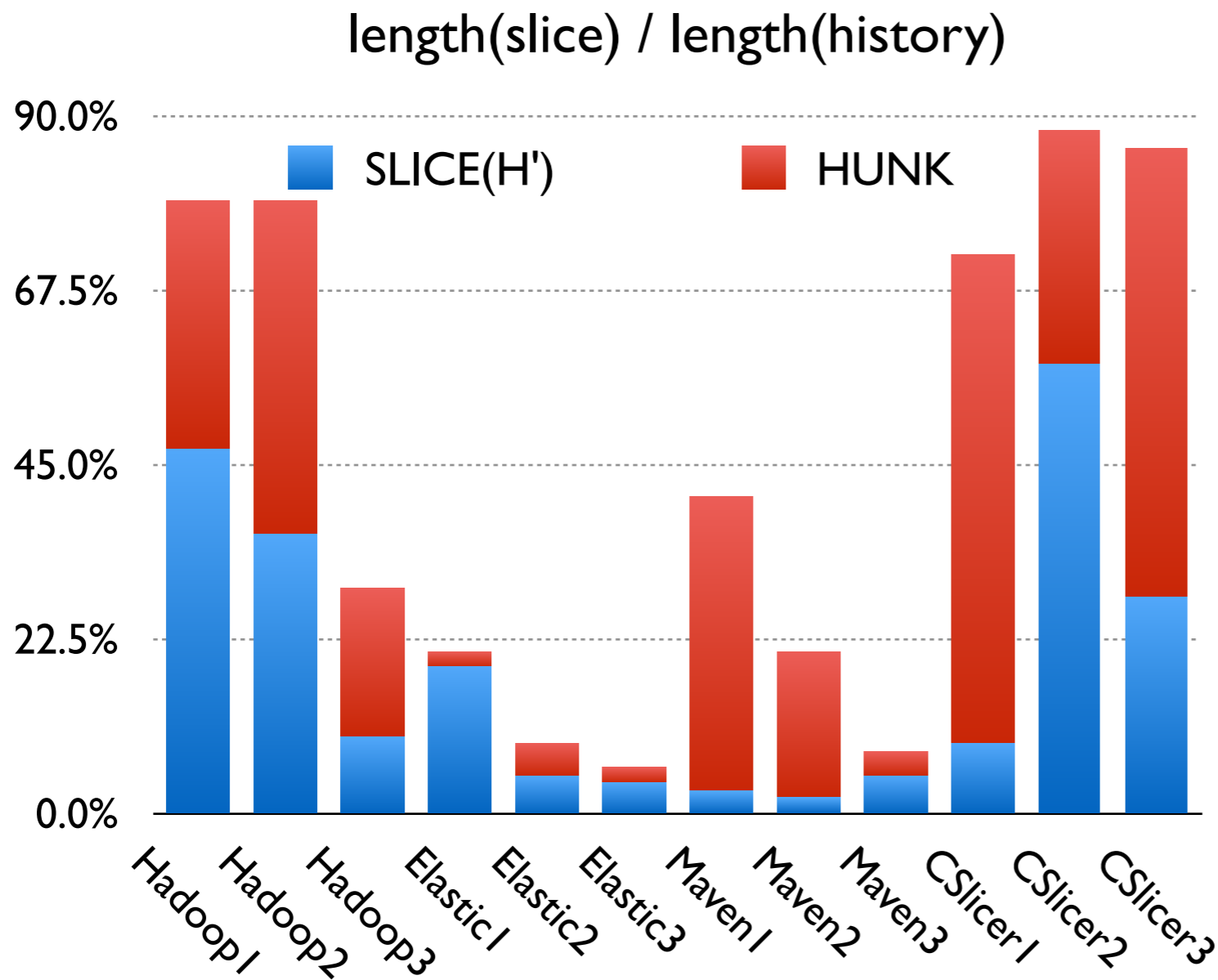
Feature branch

- Merges with the main branch periodically
- 42 feature commits + 47 merges
- 58 accompanied test cases

Case Study:

- Separate feature changes
- Identified 65 out of 267 commits related to the feature
- 41 matches original

Effectiveness



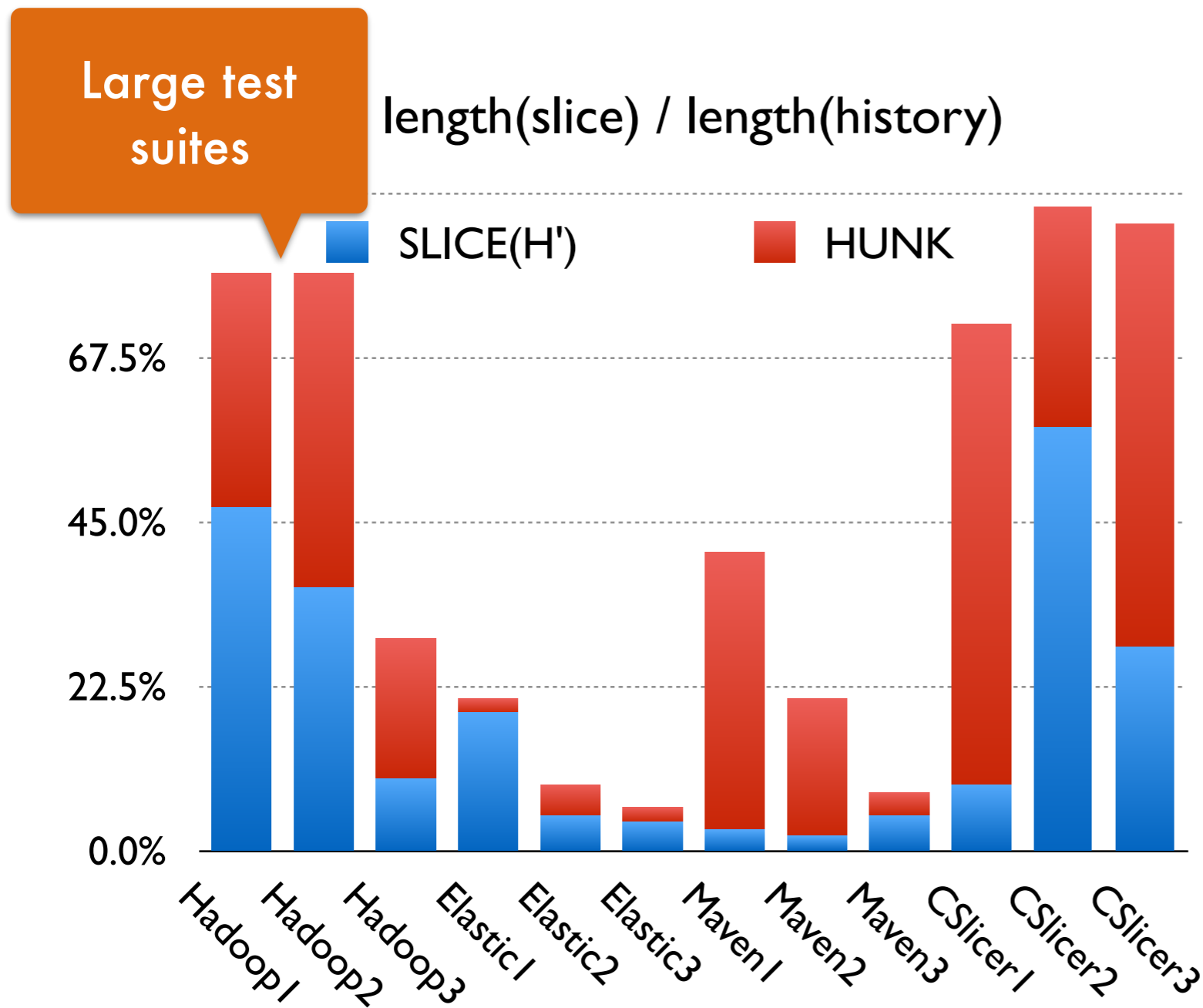
Average Reduction:

~80%!

Reduction depends on:

1. tests complexity
2. committing styles

Effectiveness



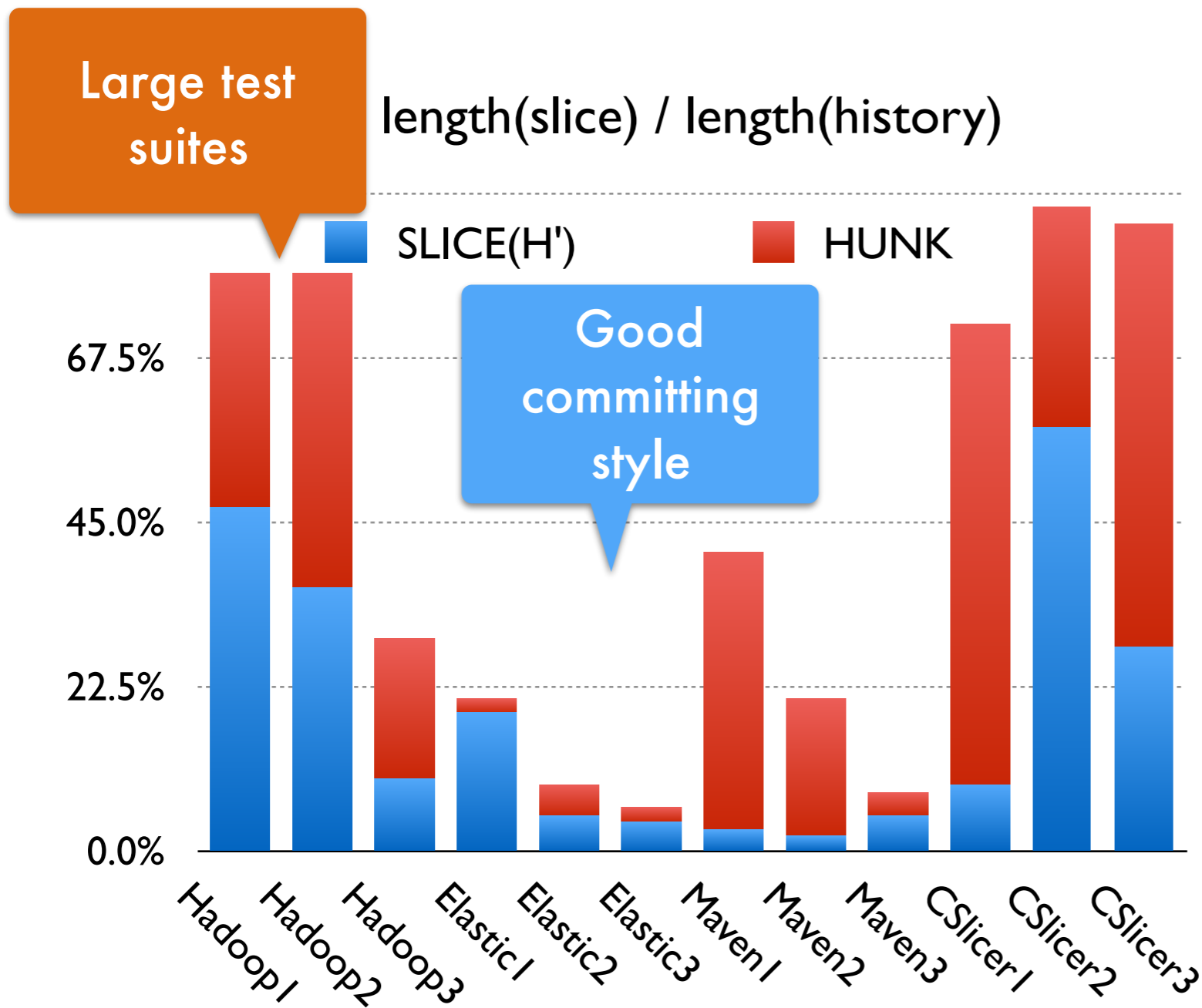
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Effectiveness



Average Reduction:

~80%!

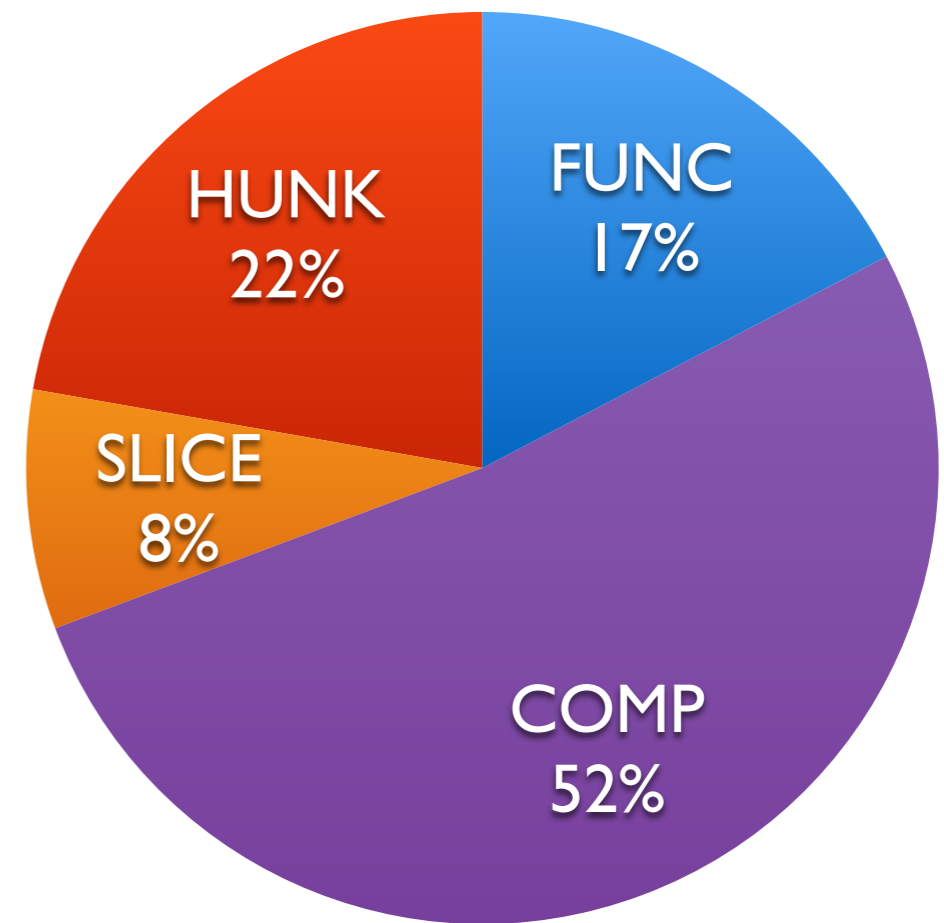
Reduction depends on:

1. tests complexity
2. committing styles

Performance

- Total CSlicer time: 2 ~ 65 s
- Major part spent in functional & compilation set computation
- History length has little effects on performance for large projects

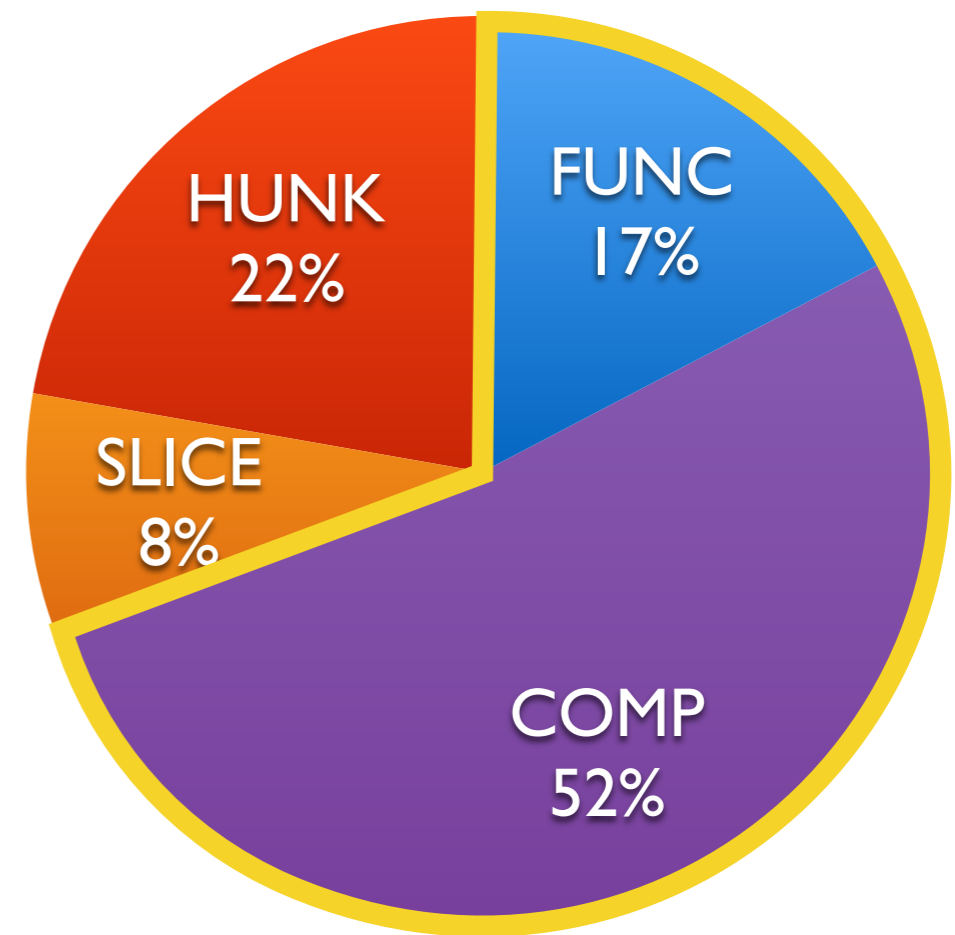
CSlicer time breakdown



Performance

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- Major part spent in functional & compilation set computation
- History length has little effects on performance for large projects

CSlicer time breakdown



Outline

1. Introduction
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5. **Related Work & Conclusion**

Related Work

Change Representation

- Code change classification [Falleri et al., ASE'14; Chawathe, SIGMOD'96]
- History granularity transformation [Muslu et al., ASE'15]

Change Impact Analysis

- Compute affected regression tests [Ren et al., OPPSLA'04]
- Fault localization [Zhang et al., ICSM'01]

Conclusion & Future Work

CSlicer: history semantic slicing

- Filling the gap between texts and semantics
- Adapted to existing version control tools
- Many interesting applications: history comprehension; functionality transferring ...

What's next?

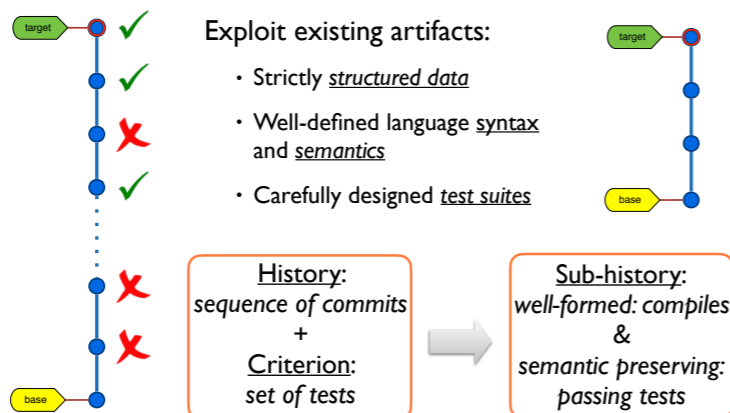
bitbucket.org/liyistc/gitslice

- Handle distributed histories
- Slice integration — the “paste” step



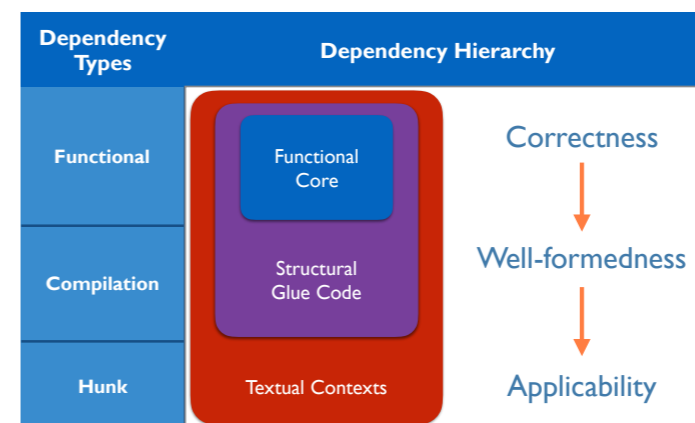
Questions?

Semantic Slicing



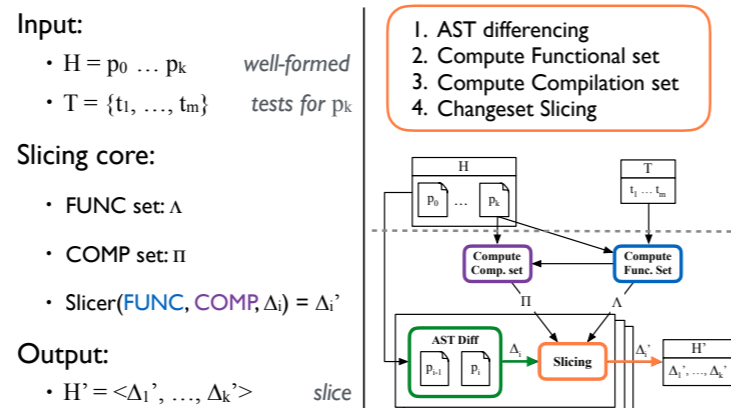
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Dependency Hierarchy



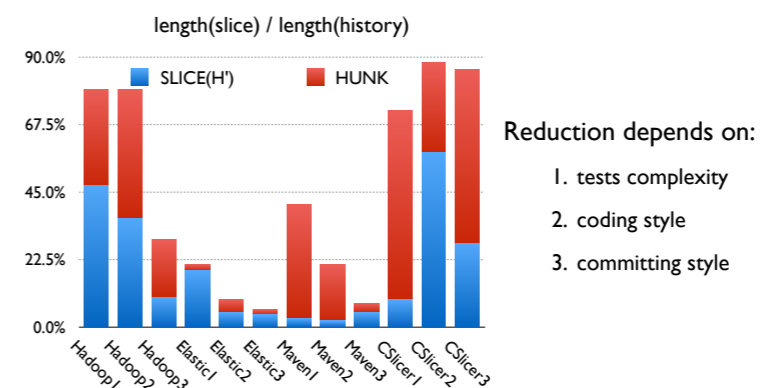
7

CSlicer Overview



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Experiments



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