

Efficient Algorithms on Cocomparability Graphs Via Vertex Ordering Characterizations

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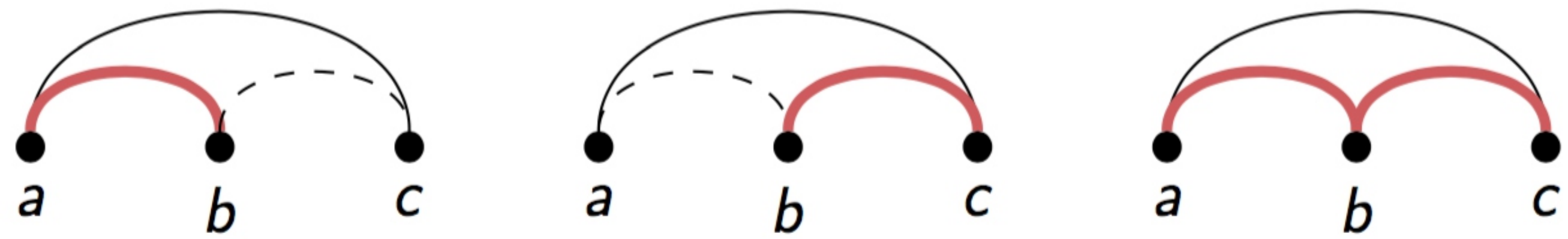
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Joint work with Ekkehard Köhler and Michel Habib

Cocomparability Graphs

- The complements of comparability graphs.
- Many-to-one mapping from posets to cocomparability graphs.
- Cocomparability graphs are a superclass to:
 - Co-graphs, trapezoid, permutation, interval graphs
- An ordering σ of V is a *cocomparability (cocomp) ordering* if for every triple $a < b < c$ in σ :



- G is a *cocomp graph* iff it admits a *cocomp ordering*.

The Algorithms

Max Weighted Independent Set (MWIS)

Input: A cocomparability graph $G=(V, E)$:

- Compute a cocomparability ordering σ .
- Scan σ left to right: compute an ordering τ , where vertices are inserted in τ in increasing order of their (updated) weight.
- Scan τ right to left to greedily collect a maximum weight independent set.

Max Weighted Induced Matching (MWIM)

Input: Given a cocomparability graph $G=(V, E)$:

- Compute a cocomparability ordering σ .
- Using either Rule \bullet or Rule \star , compute τ .
- Use the MWIS algorithm on τ to compute a maximum weight induced matching on G .

Induced Matching

- *Induced matching*: A matching where every pair of edges is at distance at least 2 in G .
- Induced Matching is NP-complete on bipartite graphs, even if the graphs have degree 3, or are planar.
- Given an ordering σ of $V(G)$, construct an ordering on the vertices of $L(G)$, the line graph of G , such that for every pair $e_i = (a,b), e_j = (u,v)$:

Rule (\bullet): $e_i \prec_{\bullet} e_j \iff a \preceq_{\sigma} u$ and $b \preceq_{\sigma} v$

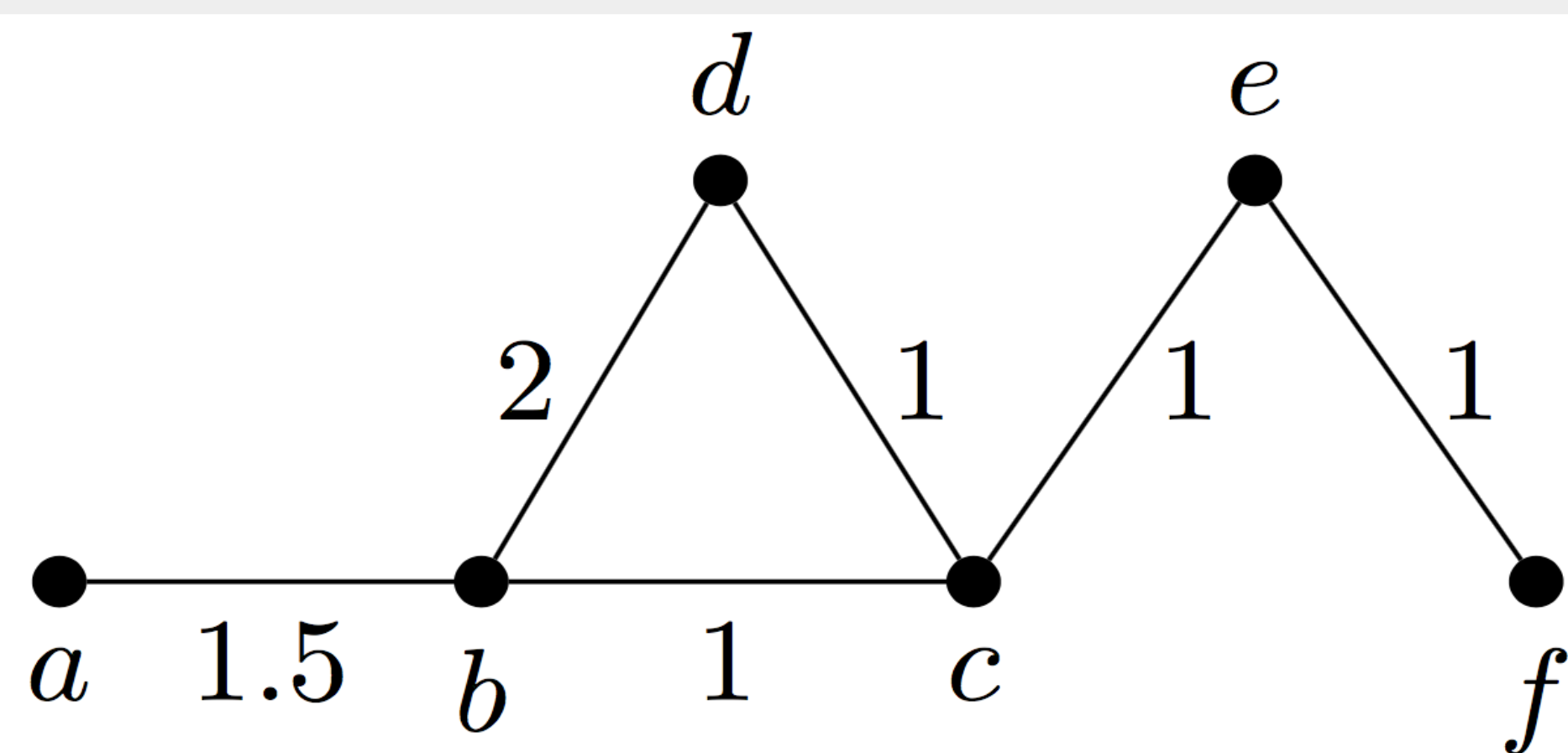
Rule (\star): $e_i \prec_{\star} e_j \iff \begin{cases} a \prec_{\sigma} u & \text{if } a \neq u \\ a = u \text{ and } b \prec_{\sigma} v & \text{o.w.} \end{cases}$

The Results:

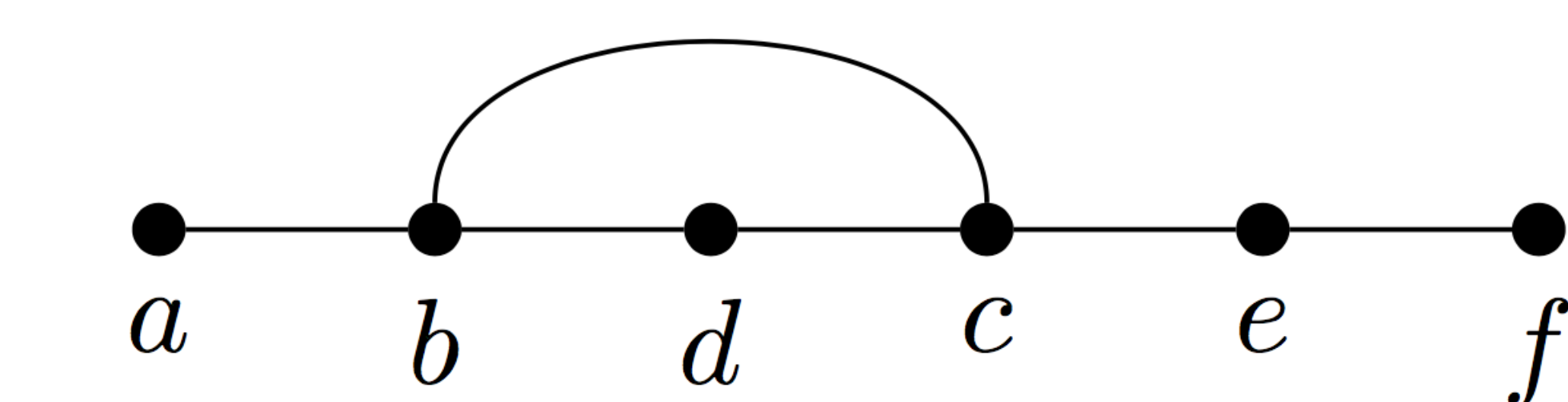
Let $S = \{\text{interval, split, threshold, cocomparability}\}$

- Given a vertex ordering of G in S , rules \bullet and \star compute vertex orderings of $L^2(G)$.
- Every graph class in S is closed under $L^2(\bullet)$.
- Maximum weighted independent set (MWIS) can be computed in $O(m+n)$ time on **cocomparability** graphs.
- Maximum weighted induced matching (MWIM) can be computed in $O(mn)$ time on **cocomparability** graphs.
- Both the MWIS and MWIM algorithms are robust.

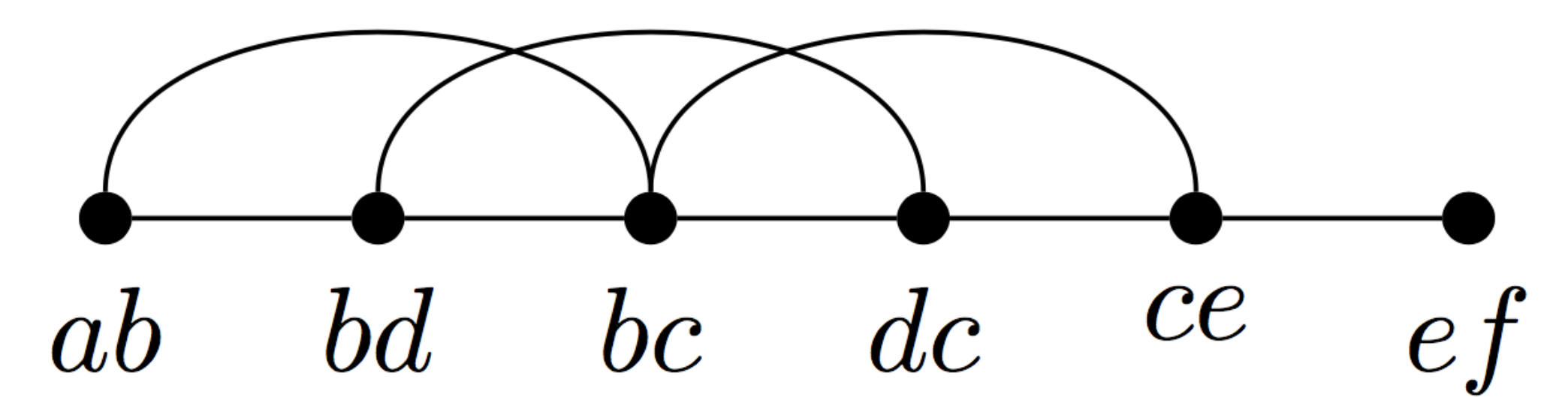
Example



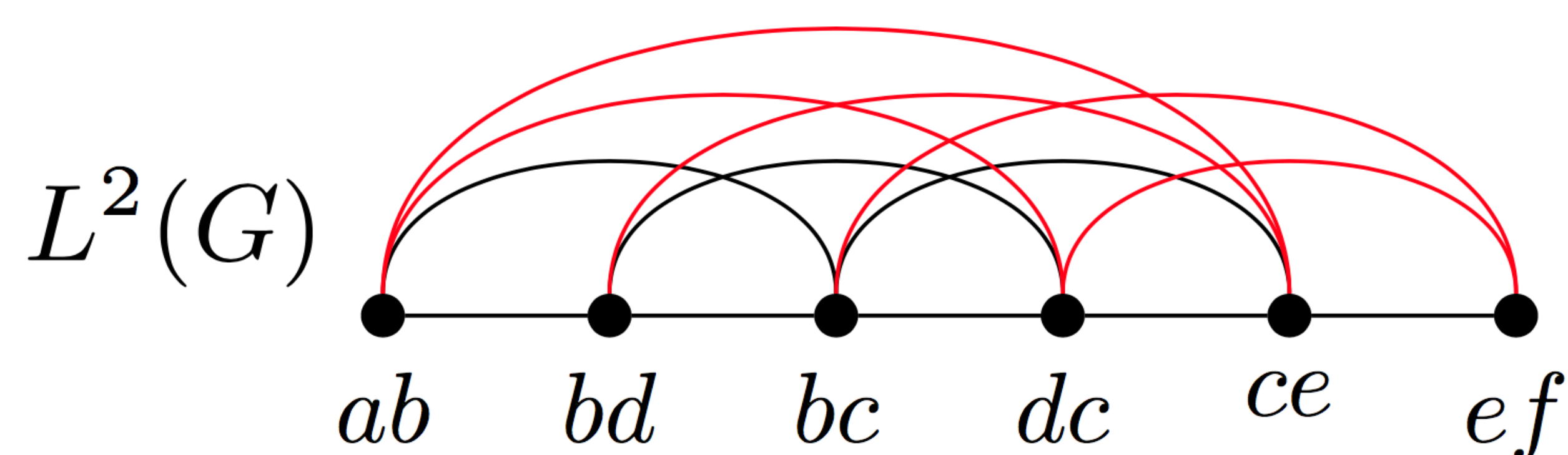
$G = (V, E, w)$



σ , a cocomparability ordering of G



$L(G)$



$w(\cdot) : 1.5 \quad 2 \quad 1 \quad 1 \quad 1 \quad 1$

e_i	u	S_{e_i}	$w(S_{e_i})$	τ_i
$e_1 = ab$	-	$\{e_1\}$	1.5	e_1
$e_2 = bd$	-	$\{e_2\}$	2	e_1, e_2
$e_3 = bc$	-	$\{e_3\}$	1	e_3, e_1, e_2
$e_4 = cd$	-	$\{e_4\}$	1	e_3, e_4, e_1, e_2
$e_5 = ce$	-	$\{e_5\}$	1	e_3, e_4, e_5, e_1, e_2
$e_6 = ef$	bd	$\{e_2, e_6\}$	3	$e_3, e_4, e_5, e_1, e_2, e_6$

Future Work

- Colouring Algorithm via VOCs.
- Certifying algorithms.
- Stepping outside perfection: AT-free graphs are closed under $L^2(\bullet)$ as well.

References

- Köhler & Mouatadid, *Linear Time MWIS on Cocomparability Graphs*. IPL
- Brandstädt & Hoàng, *Maximum induced matchings for chordal graphs in linear time*. Algorithmica
- Cameron, *Induced matchings in intersection graphs*. Discrete Math
- Habib & Mouatadid: *Efficient MIM Algorithms via VOCs*. Submitted