

# Linear Time Maximum Weight Independent Set On Cocomparability Graphs

Ekkehard Köhler<sup>◦</sup>, Lalla Mouatadid<sup>\*</sup>

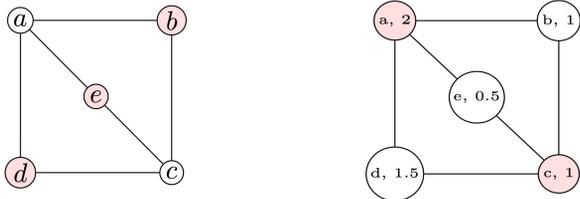
<sup>◦</sup>Department of Mathematics, Brandenburg University of Technology, Germany

<sup>\*</sup>Department of Computer Science, University of Toronto

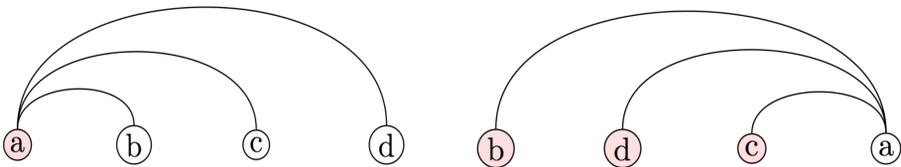
## Motivation

- Given a graph  $G(V, E)$ , where  $V$  is the set of vertices, and  $E$  the set of edges, consider the following problem: What is the largest subset of  $V$  where every two vertices are pairwise nonadjacent?

### The Maximum (Cardinality/ Weight) Independent Set Problem



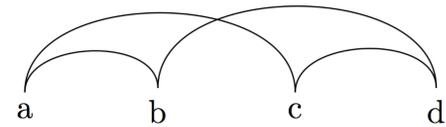
- Why do we care?
  - Scheduling. Biology. Coding Theory ...
- Very easy to trick a greedy algorithm:



- The (Weighted) Maximum Independent Set, (W)MIS, problem is NP-hard on arbitrary graphs.

## Cocomparability Graphs

- A cocomparability graph is the complement of a comparability graph (i.e. induced by a partial order).
- Vertex Ordering Characterization:**
  - $G$  is a cocomparability graph iff  $V$  admits an ordering  $(v_i)$  where every triple  $a < b < c$  with  $ac \in E$ ,  $ab \in E$  or  $bc \in E$  or both. For example:



- Such ordering is called an umbrella free ordering.
- $O(m+n)$  to compute – McConnell & Spinrad [1].

- Cocomparability graphs are a superclass to:
  - Trapezoid and Permutation graphs
  - Cographs
  - Interval graphs

► Solving a problem on cocomparability graphs yields a solution to all these graph classes !

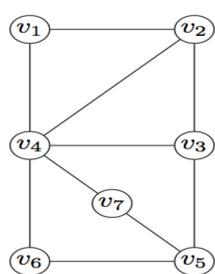
## The Algorithm

- Given a cocomparability graph  $G(V, E)$ :
- Compute a valid cocomparability ordering  $\sigma$ .
- Scan  $\sigma$  from left to right to compute a new ordering  $\tau$  of  $V$ , where vertices are inserted in  $\tau$  in nondecreasing order of their (updated) weight.
- Scan  $\tau$  from right to left to greedily collect a maximum weight independent set.

## Proof of Correctness

- Associate with every vertex a set  $S(v_i)$ , then at every iteration  $i$ :
- For every vertex  $v_i$ ,  $S(v_i)$  is an independent set.
- Every  $S(v_i)$  is a maximum weighted independent set containing  $v_i$  in  $G[v_i, \dots, v_1]$ .
- Let  $z_i$  be the rightmost vertex of  $\tau_i$ , then  $S(z_i)$  is a maximum weighted independent set in  $G[v_i, \dots, v_1]$ .

## Example



$\sigma$ :  $v_1 \quad v_2 \quad v_3 \quad v_4 \quad v_5 \quad v_6 \quad v_7$   
 $w(v)$ :  $1 \quad 0.5 \quad 1-\epsilon \quad 3\epsilon \quad 4 \quad 2-\epsilon \quad 2$

$\tau$ :  $v_1$   
 $\tilde{w}(v)$ :  $1$

$\tau$ :  $v_2 \quad v_1$   
 $\tilde{w}(v)$ :  $0.5 \quad 1$

$\tau$ :  $v_2 \quad v_1 \quad v_3$   
 $\tilde{w}(v)$ :  $0.5 \quad 1 \quad 2-\epsilon$

$\tau$ :  $v_4 \quad v_2 \quad v_1 \quad v_3$   
 $\tilde{w}(v)$ :  $3\epsilon \quad 0.5 \quad 1 \quad 2-\epsilon$

$\tau$ :  $v_4 \quad v_2 \quad v_1 \quad v_3 \quad v_5$   
 $\tilde{w}(v)$ :  $3\epsilon \quad 0.5 \quad 1 \quad 2-\epsilon \quad 5$

$\tau$ :  $v_4 \quad v_2 \quad v_1 \quad v_3 \quad v_6 \quad v_5$   
 $\tilde{w}(v)$ :  $3\epsilon \quad 0.5 \quad 1 \quad 2-\epsilon \quad 4-2\epsilon \quad 5$

$\tau$ :  $v_4 \quad v_2 \quad v_1 \quad v_3 \quad v_6 \quad v_5 \quad v_7$   
 $\tilde{w}(v)$ :  $3\epsilon \quad 0.5 \quad 1 \quad 2-\epsilon \quad 4-2\epsilon \quad 5 \quad 6-2\epsilon$

$IS = \{v_7, v_6, v_3, v_1\}$   
 $W = 6-2\epsilon$

## Future Work

- Certify the algorithm. There exists a certifying algorithm for the unweighted case that computes a minimum clique cover of equal cardinality [2].
- Extend the algorithm to the  $k$ -colourable subgraph problem.

[1] Ross M McConnell and Jeremy P Spinrad. Modular decomposition and transitive orientation. *Discrete Mathematics*. 1999

[2] Derek G Corneil, Jérémie Dusart, Michel Habib, and Ekkehard Köhler. On the power of graph searching for cocomparability graphs. *In preparation*.