Threads

lightweight processes

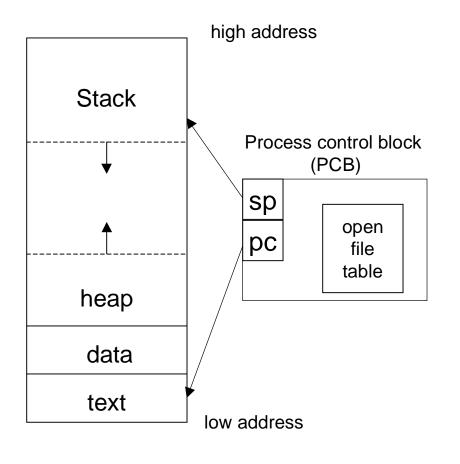
Motivation

- Processes are expensive to create.
- It takes quite a bit of time to switch between processes
- Communication between processes must be done through an external kernel structure

- files, pipes, shared memory

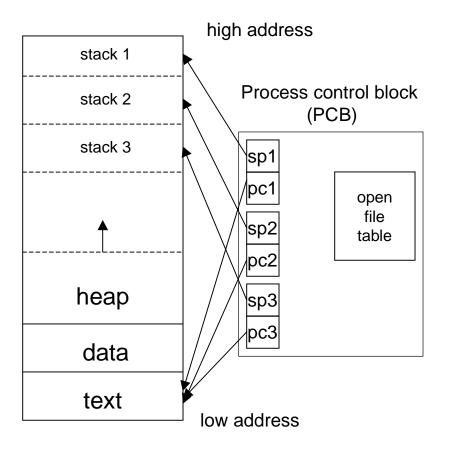
- Synchronizing between processes is cumbersome.
- Is there another model that will solve these problems?

Processes



- Each process has its own
 - program counter
 - stack
 - stack pointer
 - address space
- Processes may share
 - open files
 - pipes

Threads



- Each thread has its own
 - program counter
 - stack
 - stack pointer
- Threads share
 - address space
 - variables
 - code
 - open files

Advantages

- Communication between threads is cheap
 - they can share variables!
- Threads are "lightweight"
 - faster to create
 - faster to switch between
- Synchronization avoids kernel

Threaded design

- Several common models for threaded programs exist:
 - Manager/worker: a single manager thread assigns work to other threads, the workers. The manager typically handles all input and parcels out work to the workers.
 - Pipeline: a task is broken into a series of suboperations, each of which is handled in series, but concurrently, by a different thread. Is like an automobile assembly line.
 - Peer: similar to the manager/worker model, but after the main thread creates other threads, it participates in the work.

Pthreads

 POSIX threads (pthreads) is the most commonly used thread package on Unix/Linux

pthread_create

- tid uniquely identifies a thread within a process and is returned by the function
- attr sets attributes such as priority, initial stack size
 - can be specified as NULL to get defaults
- func the function to call to start the thread
 - accepts one void * argument, returns void *
- arg is the argument to func
- returns 0 if successful, a positive error code if not
- does not set errno but returns compatible error codes
- can use strerror() to print error messages

pthread_join

- tid the tid of the thread to wait for
 cannot wait for any thread (as in wait())
- status, if not NULL returns the void * returned by the thread when it terminates.
- a thread can terminate by
 - returning from func
 - the main() function exiting or exit() called
 - pthread_exit()
 - pthread_cancel()

More functions

- void pthread_exit(void *status)
 - a second way to exit, returns status explicitly
 - status must not point to an object local to the thread, as these disappear when the thread terminates.
- int pthread_detach(pthread_tid);
 - if a thread is detached its termination cannot be tracked with pthread_join()
 - it becomes a daemon thread
- pthread_t pthread_self(void)
 - returns the thread ID of the thread which called it
 - often see pthread_detach(pthread_self())

Passing Arguments to Threads

- We can pass any variable (including a structure or array) to our thread function.
- It assumes the thread function knows what type it is.
- This example is **bad** if the main thread alters fd later.

Solution

- Use malloc() to create memory for the variable
 - initialize variable's value
 - pass pointer to new memory via pthread_create()
 - thread function releases memory when done.

```
• Example:
```

```
typedef struct myArg {
```

```
int fd;
```

```
char name[25];
```

```
} MyArg;
```

```
int result;
pthread_t thread_ID;
```

Example (cont'd)

MyArg *p = (MyArg *)malloc(sizeof(MyArg)); p->fd = fd; /* assumes fd is defined */ strncpy(p->name, "CSC209", 7); result = pthread_create(&threadID, NULL, myThreadFcn, (void *)p); void *myThreadFcn(void *p) { MyArg *theArg = (MyArg *) p; write(theArg->fd, theArg->name, 7); close(theArg->fd); free(theArg); return NULL;

Thread-safe functions

- Not all functions can be called from threads
 - many use global/static variables
 - new versions of UNIX have thread-safe
 replacements like strtok_r()
- Safe:
 - ctime_r(), gmtime_r(), localtime_r(), rand_r(), strtok_r()
- Not Safe:
 - ctime(), gmtime(), localtime(), rand(), strtok(), gethostxxx()
- Could use semaphores to protect access but will generally result in poor performance.

Pthread Mutexes (Semaphores)

int pthread_mutex_lock(pthread_mutex_t *mp); int pthread_mutex_trylock(pthread_mutex_t *mp); int pthread_mutex_unlock(pthread_mutex_t *mp); int pthread_mutex_destroy(pthread_mutex_t *mp);

- easier to use than semget() and semop()
- only the thread that locks a mutex can unlock it
- mutexes often declared as globals

Example

```
pthread_mutex_t myMutex;
int status;
```

```
status = pthread_mutex_init(&myMutex, NULL);
if(status != 0)
```

```
printf("Error: %s \n", strerror(status));
pthread_mutex_lock(&myMutex);
/* critical section here */
pthread_mutex_unlock(&myMutex);
status = pthread_mutex_destroy(&myMutex);
if(status != 0)
    printf("Error: %s\n", strerror(status));
```