

Search Combinators

Tom Schrijvers



with Guido Tack, Pieter Wuille, Horst Samulowitz, Peter Stuckey

Search heuristics
are crucial.

Support for Search?

General Purpose
Programming Language



Solver-Provided
Options



“everything is possible,
nothing is easy”

“everything is easy,
nothing is possible”

State-of-the-Art Modularity: Prolog

- ✓ `label1(Vars1), label2(Vars2)`
- ✓ `label1(Vars1) ; label2(Vars2)`
- ✓ `once(label1(Vars1))`
- ✓ `once((label1(Vars1),label2(Vars2)))`

Lack of Modularity

✗ `lds(label1(vars1))`

Lack of Modularity

- ✗ `Ids(label1(Vars1))`
- ✗ `Ids((label1(Vars1) , label2(Vars2)))`

Lack of Modularity

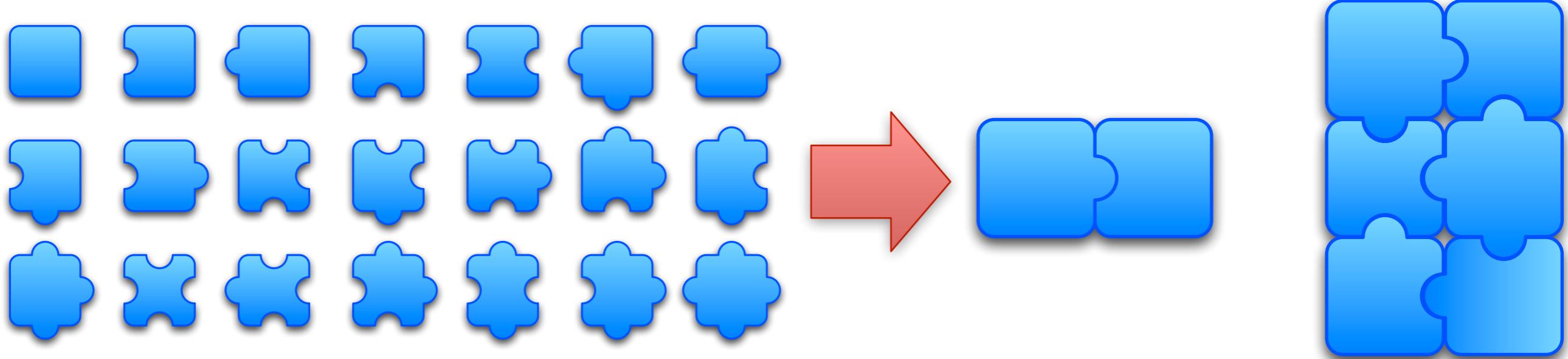
- ✗ `Ids(label1(Vars1))`
- ✗ `Ids((label1(Vars1) , label2(Vars2)))`
- ✓ `Ids_label1(Vars1)`
- ✓ `Ids_label1_label2(Vars1,Vars2)`

Can we do better?

- ✓ Lots of **expressivity** and flexibility
- ✓ Lots of **productivity** through high-level specifications
- ✓ Modular reuse of search specifications

Yes: Search Combinators

High-level **modular** building blocks



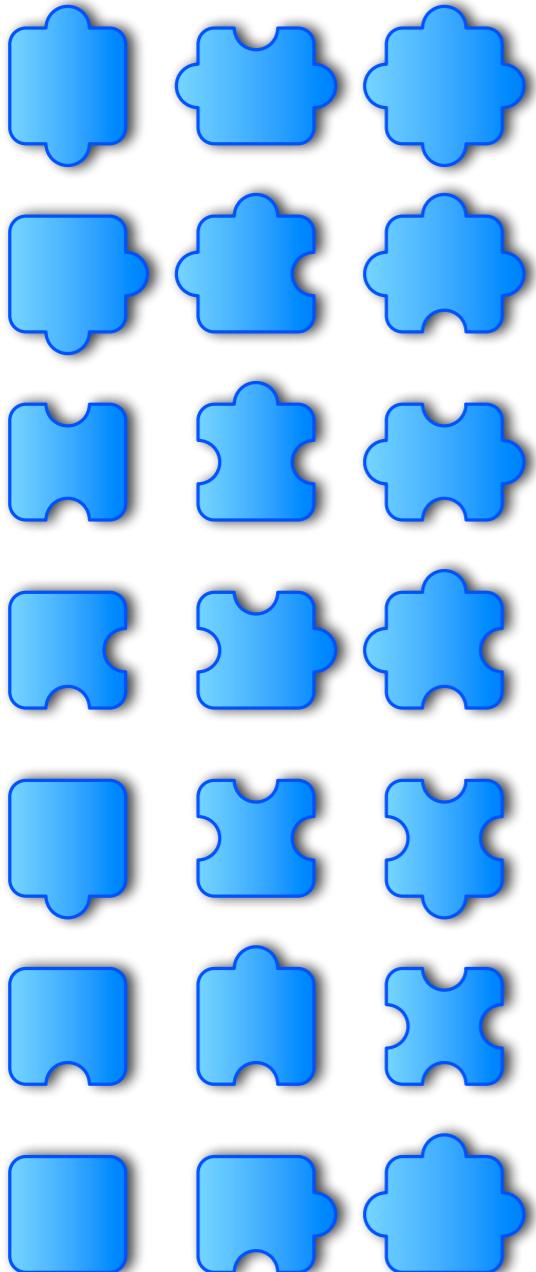
“Everything is possible and easy”

Base Search

```
s  ≡  
int_search(vars,var_sel,val_sel)
```

- provided by the solver
- augment with combinator

Combinators



prune

let(v, e, s)

assign(v, e)

post(c, s)

if(c, s_1, s_2)

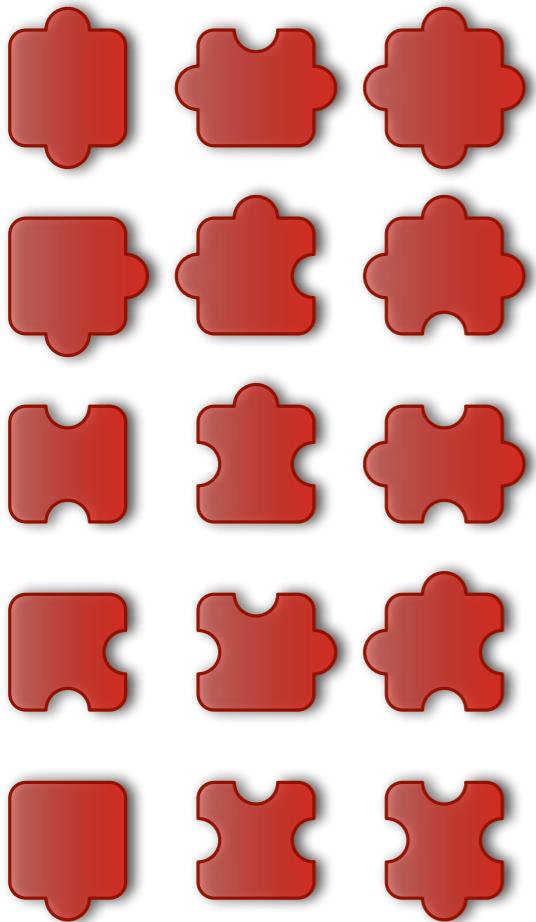
and([s_1, s_2, \dots, s_n])

or([s_1, s_2, \dots, s_n])

portfolio([s_1, s_2, \dots, s_n])

restart(c, s)

Statistics Combinators



$\text{depth}(v, s)$

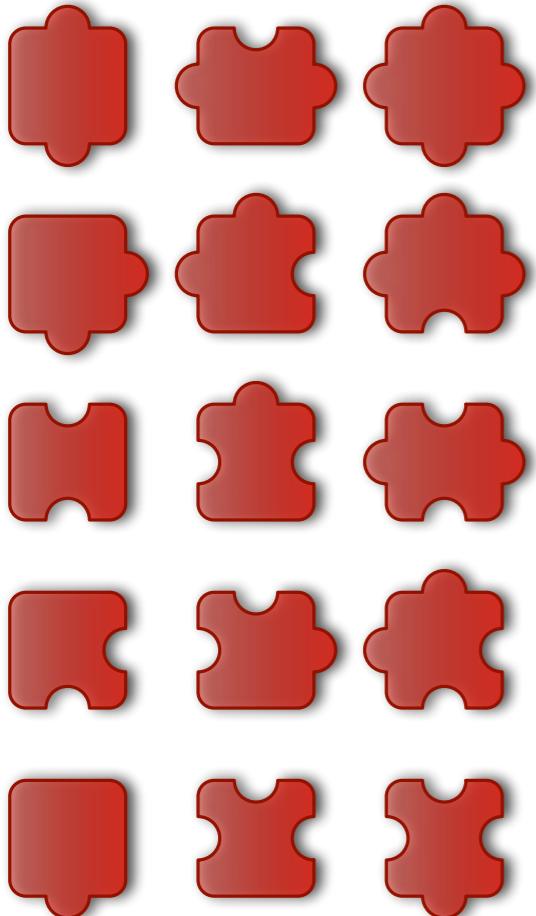
$\text{discrepancy}(v, s)$

$\text{nodes}(v, s)$

$\text{failures}(v, s)$

$\text{time}(v, s)$

Statistics Combinators



`depth(v, s)`

`discrepancy(v, s)`

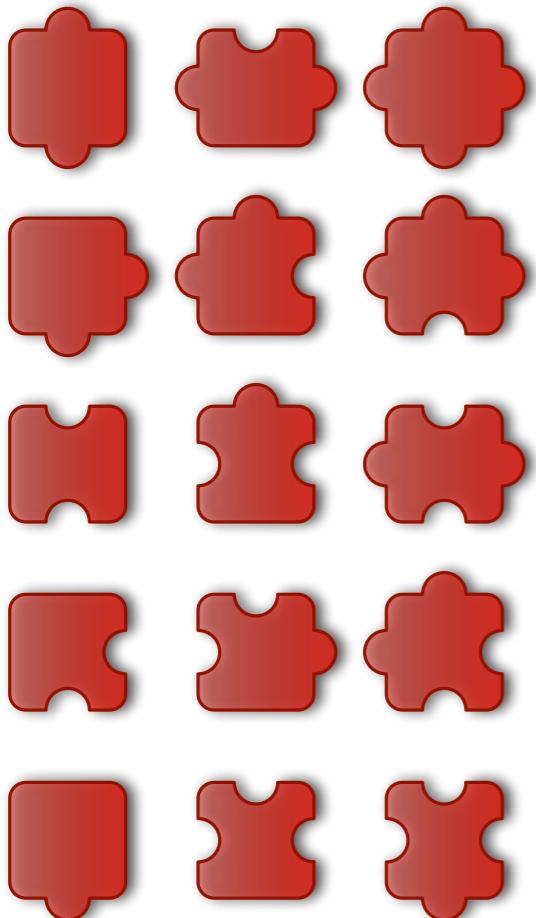
`nodes(v, s)`

`failures(v, s)`

`time(v, s)`

```
if(v < 5, depth(v, s1), s2)
```

Statistics Combinators



`depth(v, s)`

`discrepancy(v, s)`

`nodes(v, s)`

`failures(v, s)`

`time(v, s)`

```
if(depth < 5, s1, s2)
```

Reusable Abstractions

```
limit(c, s) ≡ if(c, s, prune)
```

```
for(v, l, u, s) ≡ ...
```

```
lds(s) ≡  
  for(n, 0, ∞,  
    limit(discrepancy ≤ n, s)  
)
```

Reuse Examples

- ✓ `lds(int_search(vars, ...))`
- ✓ `lds(and([int_search(vars1, ...), int_search(vars2, ...)]))`
- ✓ `lds(or([int_search(vars1, ...), int_search(vars2, ...)]))`
- ...

More Abstractions

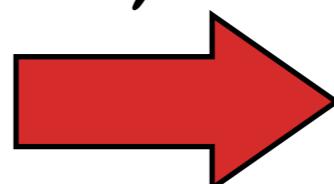
bab(*obj, s*)

restart_bab(*obj, s*)

dicho(*obj, s, lb, ub*)

id(*s*)

hot_start(*c, s₁, s₂*)



see paper

Radiotherapy Planning

```
bab(k,  
    and([int_search(N,...)]  
        ++  
        [once(int_search(rowi,...))  
         | i in 1..n  
        ]  
    )  
)
```

Modular Syntax

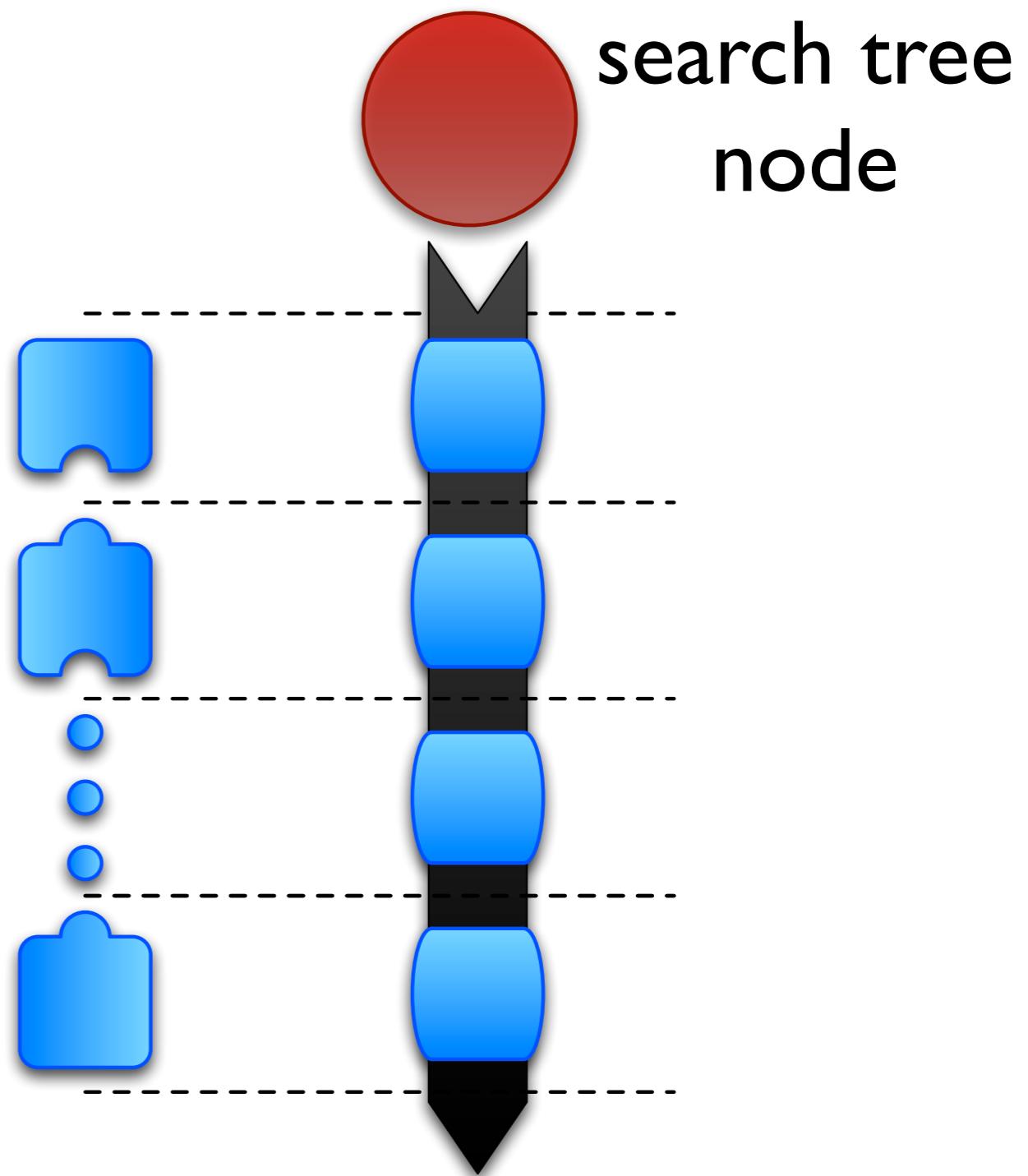
vs.

Modular Semantics

Modular Semantics

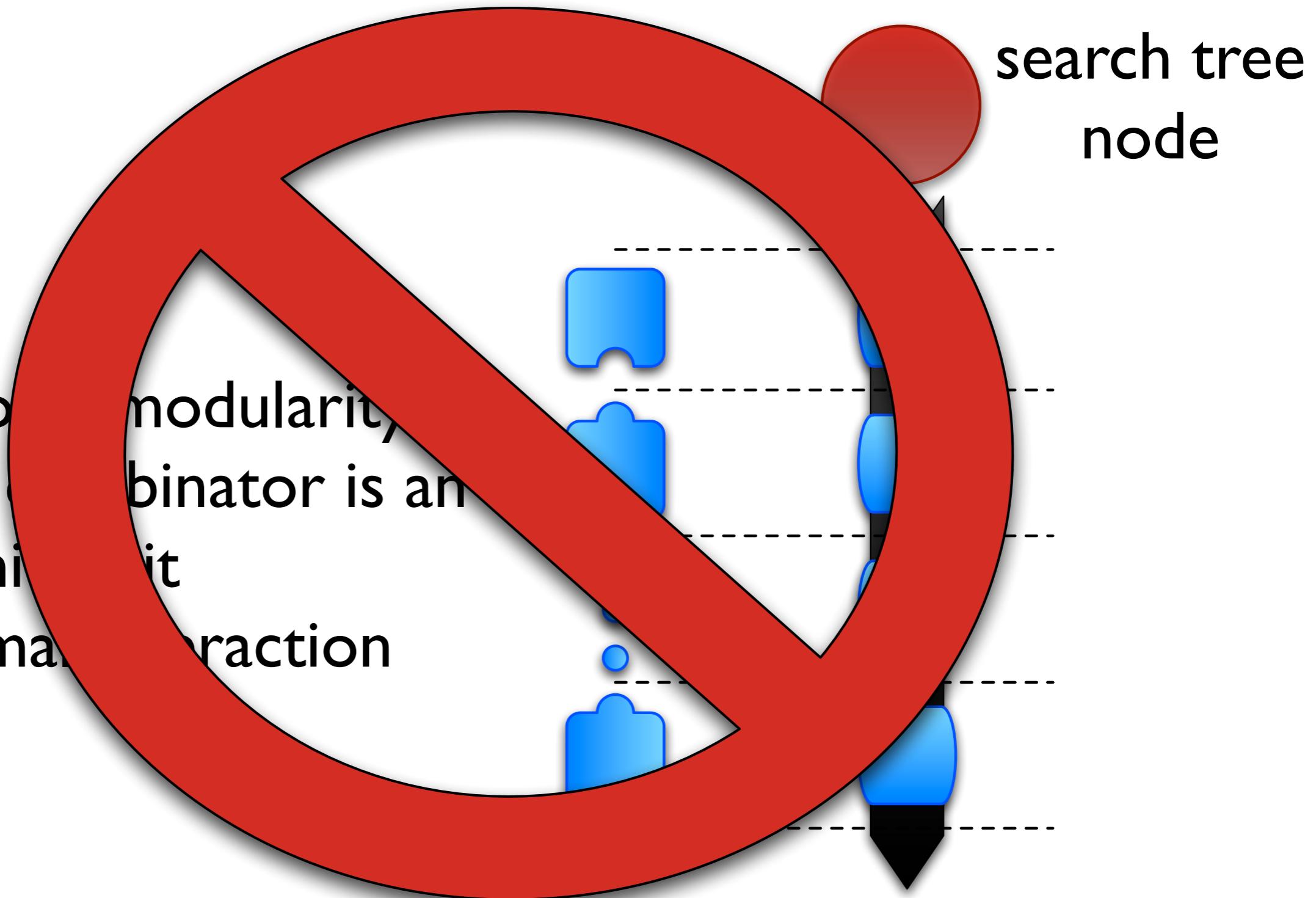
traditional modularity:

- each combinator is an atomic unit
- minimal interaction



Modular Semantics

- traditional
• each component is modular,
 atomically self-contained
• minimally interacting



Modular Semantics

- cross-cutting behavior
- highly entangled in monolithic code

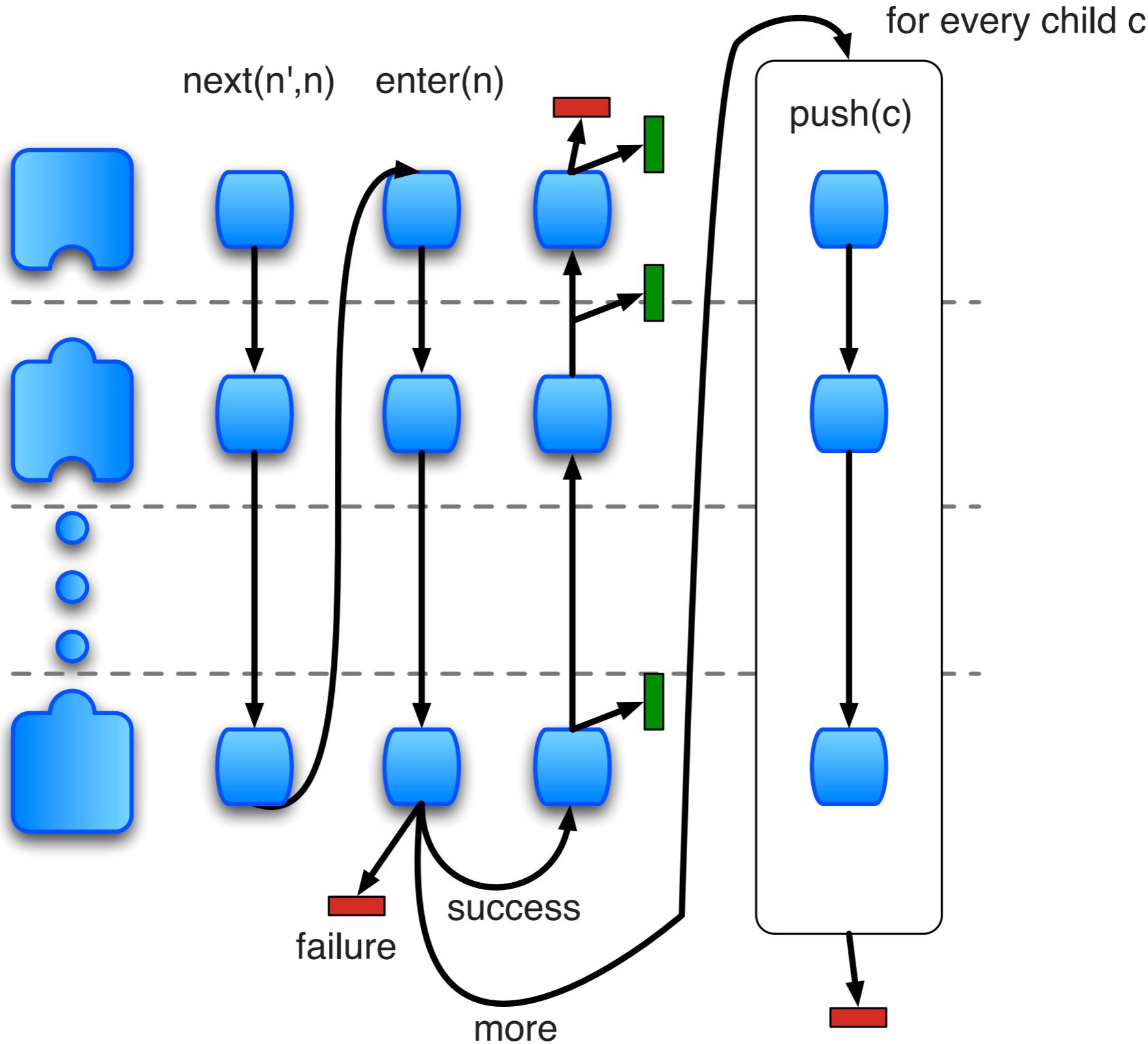


➡ “Aspect-Oriented Programming”

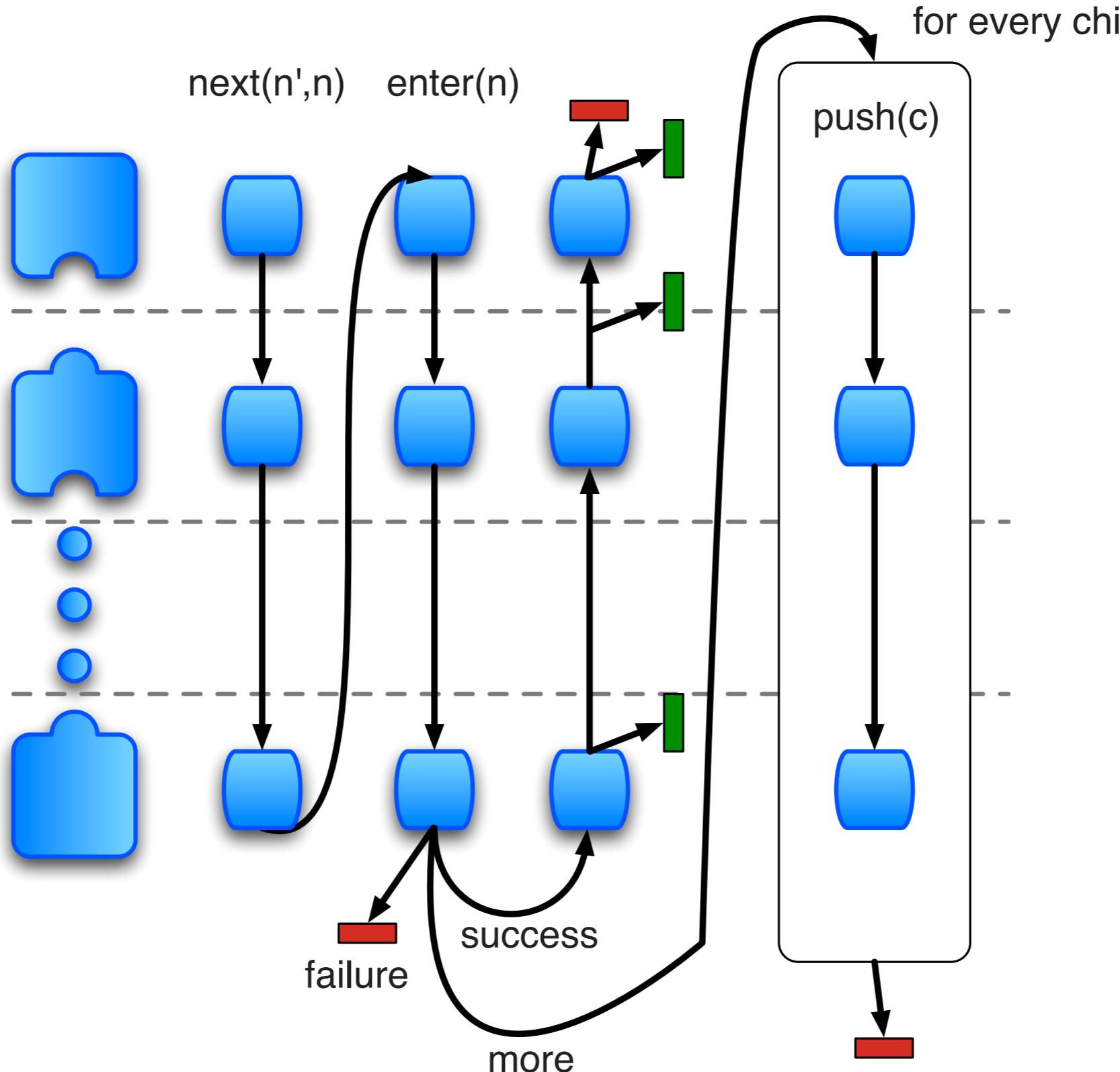
Functional Mixin Inheritance

- disciplined form of AOP
 - ★ meaningful and pre-defined set of interaction points
- no AOP system needed

Mixin-based Interaction



Mixin-based Interaction



Details:
see paper

Implementations

DSL

Haskell

C++

Compact Loop

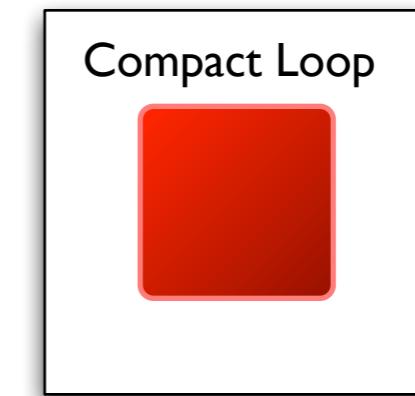
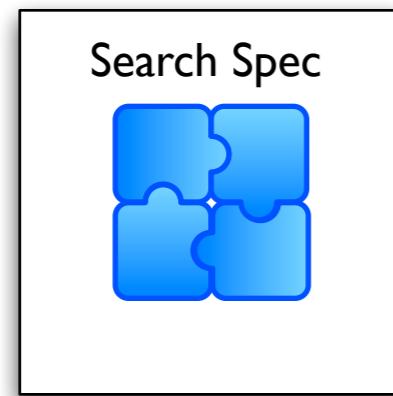


Implementations

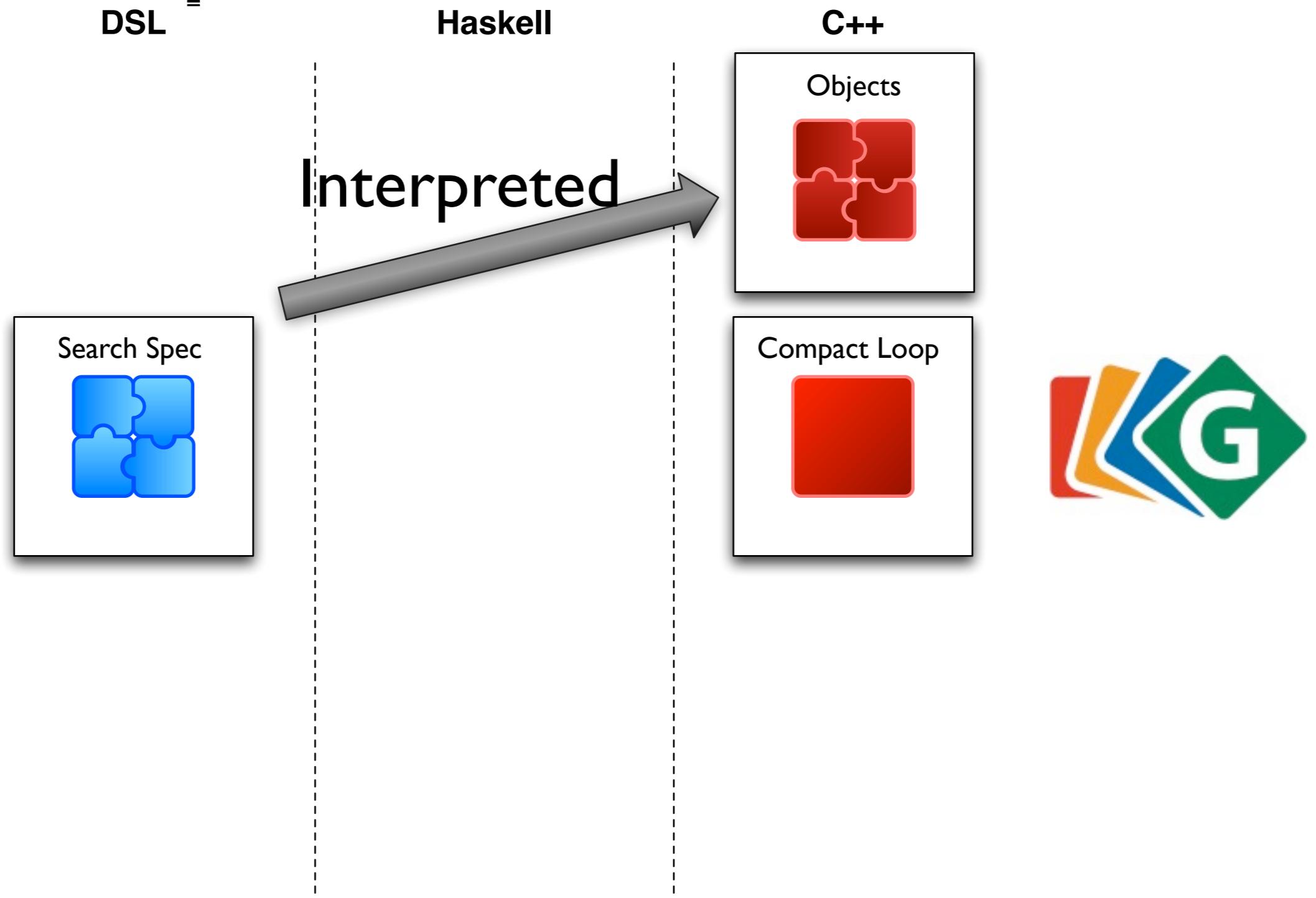
DSL

Haskell

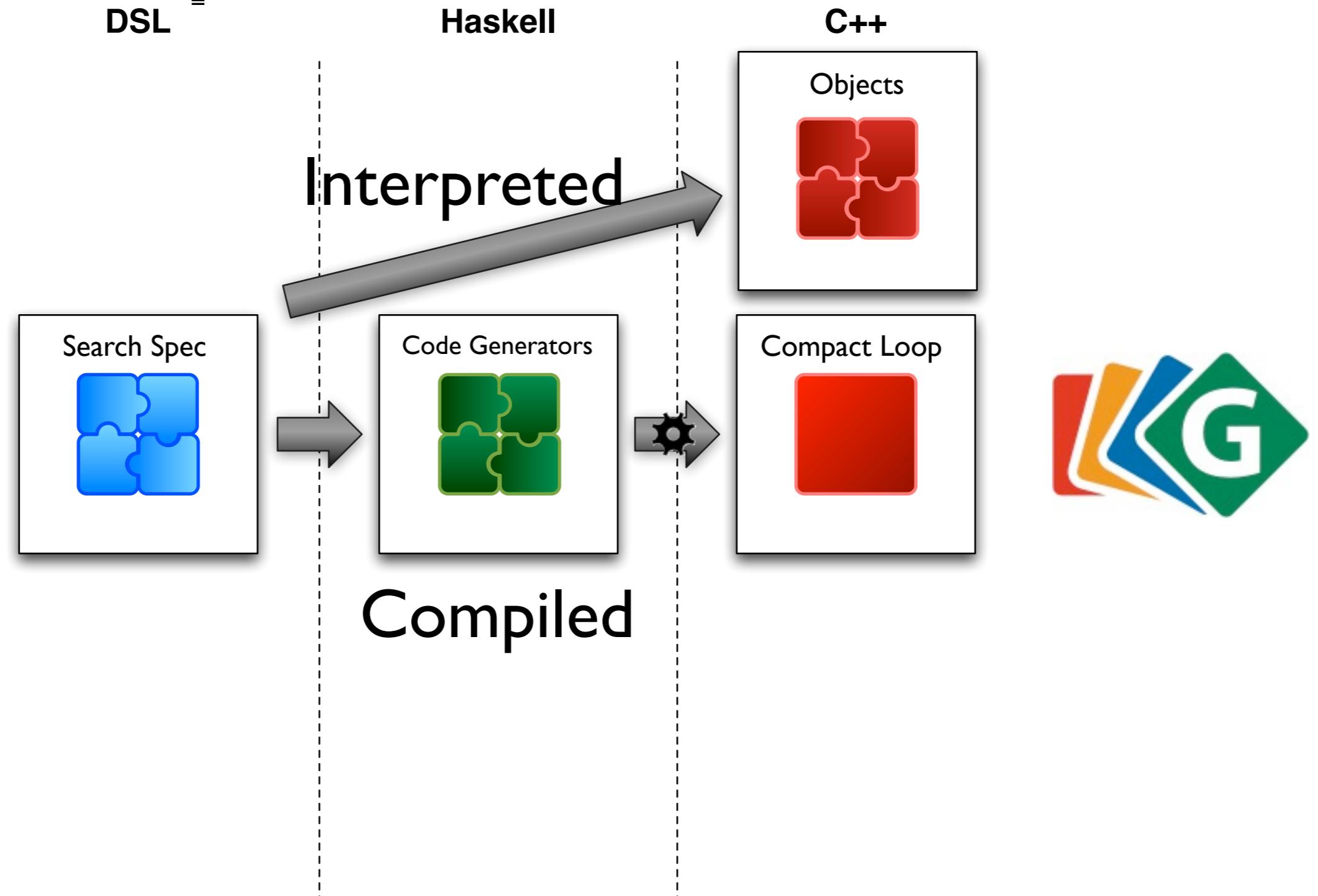
C++



Implementations

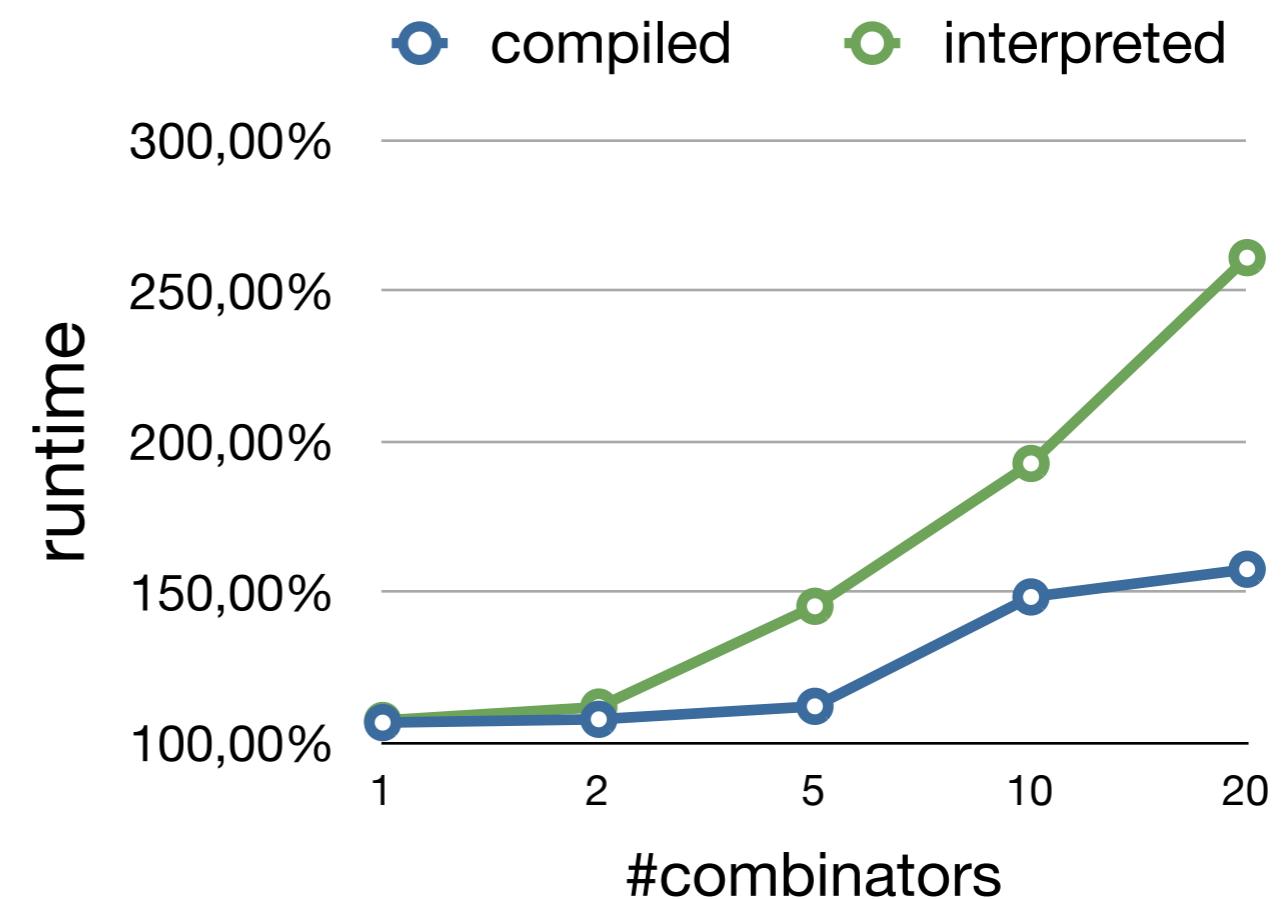
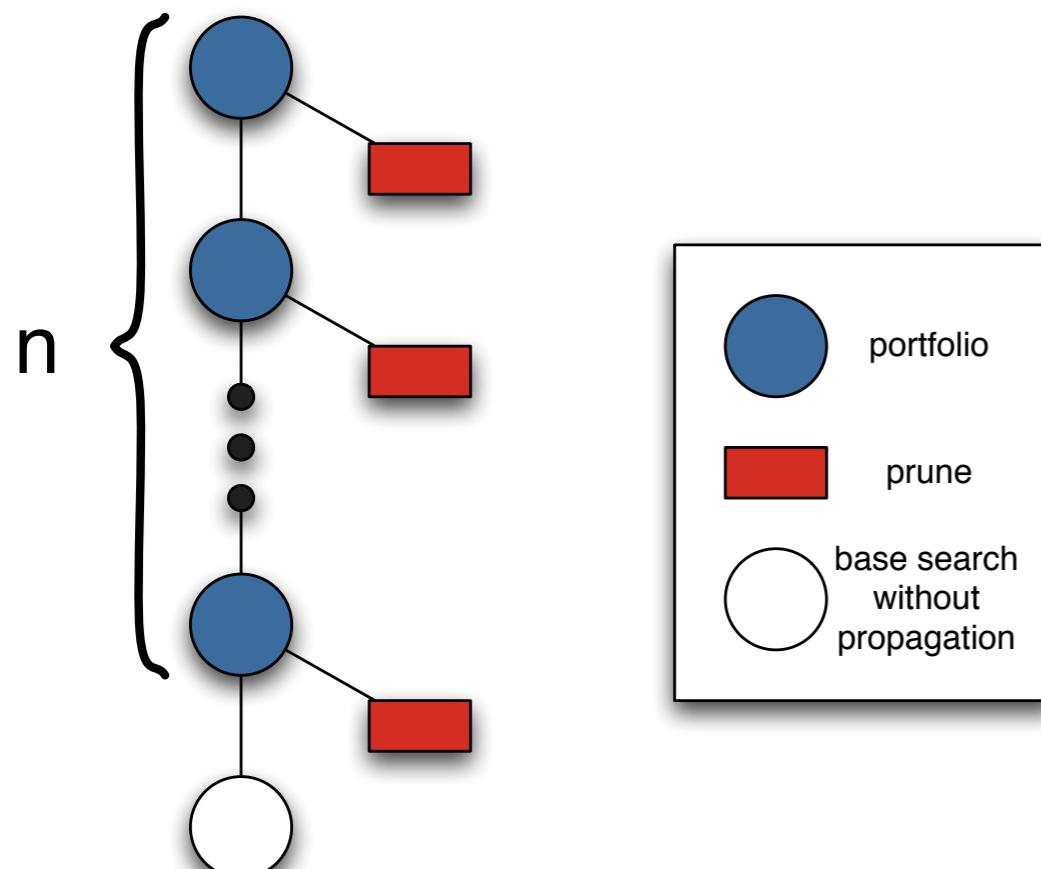


Implementations

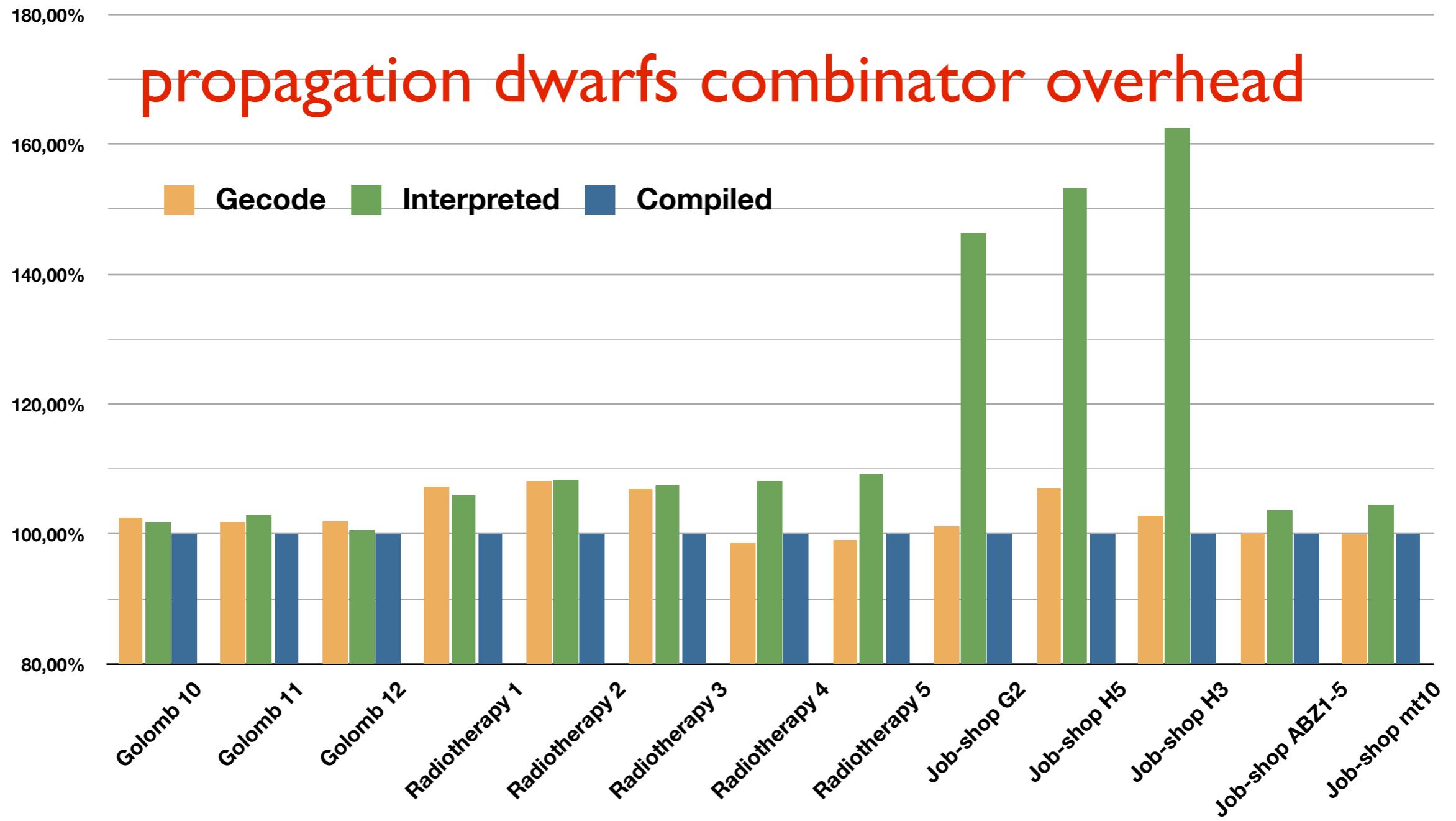


Combinator Overhead?

Worst-case Scenario



In Practice



Summary

- high-level modular modeling of search
- low-level modular implementation
- competitive performance compared to hand-coded algorithm

Future Work

- Optimizations
- MiniZinc integration
- Combinators for parallel search
- Extend other solving technology (e.g., LP)
 - ➡ Combinators for hybrid search

Thank You!



Combinator Overhead

