

- 94(a) If $P: \text{bin} \rightarrow \text{bin}$ is monotonic, prove $(\exists x \cdot P x) = P \top$ and $(\forall x \cdot P x) = P \perp$.
(b) If $P: \text{bin} \rightarrow \text{bin}$ is antimonotonic, prove $(\exists x \cdot P x) = P \perp$ and $(\forall x \cdot P x) = P \top$.

After trying the question, scroll down to the solution.

§(a)	\top	definition of monotonic specialize
=	$\forall x, y: \text{bin} \cdot (x \Rightarrow y \implies P x \Rightarrow P y)$	
\Rightarrow	$(\perp \Rightarrow \top \implies P \perp \Rightarrow P \top)$	base and identity laws
=	$P \perp \Rightarrow P \top$	laws of inclusion
=	$(P \perp \vee P \top) = P \top \wedge (P \perp \wedge P \top) = P \perp$	quantifier laws: one-element domain
=	$((\exists x: \perp \cdot P x) \vee (\exists x: \top \cdot P x)) = P \top \wedge ((\forall x: \perp \cdot P x) \wedge (\forall x: \top \cdot P x)) = P \perp$	quantifier laws: union domain
=	$(\exists x: \text{bin} \cdot P x) = P \top \wedge (\forall x: \text{bin} \cdot P x) = P \perp$	
§(b)	\top	definition of antimonotonic specialize: \perp for x , and \top for y
=	$\forall x, y: \text{bin} \cdot (x \Rightarrow y \implies P x \Leftarrow P y)$	
\Rightarrow	$(\perp \Rightarrow \top \implies P \perp \Leftarrow P \top)$	base and identity laws
=	$P \perp \Leftarrow P \top$	idempotent law
=	$(P \perp \Leftarrow P \top) \wedge (P \perp \Leftarrow P \top)$	laws of inclusion
=	$(P \perp \vee P \top) = P \perp \wedge (P \perp \wedge P \top) = P \top$	quantifier laws: one-element domain
=	$((\exists x: \perp \cdot P x) \vee (\exists x: \top \cdot P x)) = P \perp \wedge ((\forall x: \perp \cdot P x) \wedge (\forall x: \top \cdot P x)) = P \top$	quantifier laws: union domain
=	$(\exists x: \text{bin} \cdot P x) = P \perp \wedge (\forall x: \text{bin} \cdot P x) = P \top$	