

479 Let x and y be natural variables. Ignoring time, rewrite the following program as a program that does not use \parallel .

- (a) $x := x+1 \parallel \text{if } x=0 \text{ then } y := 1 \text{ else } ok \text{ fi}$
- (b) $\text{if } x > 0 \text{ then } y := x-1 \text{ else } ok \text{ fi} \parallel \text{if } x=0 \text{ then } x := y+1 \text{ else } ok \text{ fi}$

After trying the question, scroll down to the solution.

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|-----|--|----------------------------------|
| (a) | $x := x+1 \parallel \text{if } x=0 \text{ then } y := 1 \text{ else } ok \text{ fi}$ | |
| § | $x := x+1 \parallel \text{if } x=0 \text{ then } y := 1 \text{ else } ok \text{ fi}$ | expand assignments and <i>ok</i> |
| = | $x' = x+1 \parallel \text{if } x=0 \text{ then } y'=1 \text{ else } y'=y \text{ fi}$ | concurrent composition |
| = | $x' = x+1 \wedge \text{if } x=0 \text{ then } y'=1 \text{ else } y'=y \text{ fi}$ | distribution |
| = | $\text{if } x=0 \text{ then } x' = x+1 \wedge y'=1 \text{ else } x' = x+1 \wedge y'=y \text{ fi}$ | substitution law and identity |
| = | $\text{if } x=0 \text{ then } y := 1. \ x' = x+1 \wedge y'=y \text{ else } ok. \ x' = x+1 \wedge y'=y \text{ fi}$ | assignment |
| = | $\text{if } x=0 \text{ then } y := 1. \ x := x+1 \text{ else } ok. \ x := x+1 \text{ fi}$ | distribution |
| = | $\text{if } x=0 \text{ then } y := 1 \text{ else } ok \text{ fi}. \ x := x+1$ | |
| (b) | $\text{if } x>0 \text{ then } y := x-1 \text{ else } ok \text{ fi} \parallel \text{if } x=0 \text{ then } x := y+1 \text{ else } ok \text{ fi}$ | |
| § | $\text{if } x>0 \text{ then } y := x-1 \text{ else } ok \text{ fi} \parallel \text{if } x=0 \text{ then } x := y+1 \text{ else } ok \text{ fi}$ | asmts and <i>ok</i> |
| = | $\text{if } x>0 \text{ then } y' = x-1 \text{ else } y'=y \text{ fi} \parallel \text{if } x=0 \text{ then } x' = y+1 \text{ else } x'=x \text{ fi}$ | conc. comp. |
| = | $\text{if } x>0 \text{ then } y' = x-1 \text{ else } y'=y \text{ fi} \wedge \text{if } x=0 \text{ then } x' = y+1 \text{ else } x'=x \text{ fi}$ | <i>x</i> is natural |
| = | $\text{if } x>0 \text{ then } y' = x-1 \text{ else } y'=y \text{ fi} \wedge \text{if } \neg(x>0) \text{ then } x' = y+1 \text{ else } x'=x \text{ fi}$ | case revers. |
| = | $\text{if } x>0 \text{ then } y' = x-1 \text{ else } y'=y \text{ fi} \wedge \text{if } x>0 \text{ then } x'=x \text{ else } x' = y+1 \text{ fi}$ | case distributive |
| = | $\text{if } x>0 \text{ then } y' = x-1 \wedge x'=x \text{ else } x' = y+1 \wedge y'=y \text{ fi}$ | assignment twice |
| = | $\text{if } x>0 \text{ then } y := x-1 \text{ else } x := y+1 \text{ fi}$ | |