(input implementation) Let $W$ be “wait for input on channel $c$ and then read it”.

(a) \[ W = t := \max t (T_r + 1). \] $c$
Prove $W \iff \sqrt{c} \text{ if } c \text{? then } c \text{? else } t := t+1. \ \text{W fi}
$ assuming time is an extended integer.

(b) Now let time be an extended real, redefine $W$ appropriately, and reprove the refinement.
§ With integer time, we took communication transit time to be 1, and recursive call time to be 1, and the computation reads the input the instant it becomes available. If time is an extended real, we can't expect transit time to equal recursive call time, and we can't expect to read the input the instant it is available. The best we can do is to read it at the first opportunity after it is available. Let the communication transit time be $a$, and let the recursive call time be 1 (because that's what's in the refinement). I redefine
\[ W = t := \max t (t + \text{ceil} (T_r + a - t)). \] $c$
The input check $\sqrt{c}$ is $T_r + a \leq t$. Now I prove the refinement by cases. First
\[ \sqrt{c} \land (r := r+1) \]
\[ \equiv T_r + a \leq t \land (t := t, r := r+1) \]
\[ \equiv T_r + a \leq t \land (t := \max t (t + \text{ceil} (T_r + a - t)), r := r+1) \]
\[ \implies W \]
And the other case:
\[ \neg \sqrt{c} \land (t := t+1, W) \]
\[ \equiv T_r + a > t \land (t := \max (t+1) (t + 1 + \text{ceil} (T_r + a - t - 1)), r := r+1) \]
\[ \equiv T_r + a > t \land (t := \max (t+1) (t + \text{ceil} (T_r + a - t)), r := r+1) \]
Since $T_r + a > t$, then $\text{ceil} (T_r + a - t) \geq 1$, so $\max$ is its second argument, and it does no harm to decrease the first argument
\[ \equiv T_r + a > t \land (t := \max t (t + \text{ceil} (T_r + a - t)), r := r+1) \]
\[ \implies W \]