

299 Prove

- (a) the Precondition Law: A is a sufficient precondition for specification P to be refined by specification S if and only if $A \Rightarrow P$ is refined by S .
- (b) the Postcondition Law: A' is a sufficient postcondition for specification P to be refined by specification S if and only if $A' \Rightarrow P$ is refined by S .

After trying the question, scroll down to the solution.

(a) the Precondition Law: A is a sufficient precondition for specification P to be refined by specification S if and only if $A \Rightarrow P$ is refined by S .

$$\begin{aligned}
 \S & \quad (\forall \sigma. A \Rightarrow (\forall \sigma'. P \Leftarrow S)) = \forall \sigma, \sigma'. (A \Rightarrow P) \Leftarrow S && \text{distribution} \\
 = & \quad (\forall \sigma, \sigma'. A \Rightarrow (P \Leftarrow S)) = \forall \sigma, \sigma'. (A \Rightarrow P) \Leftarrow S && \text{portation twice} \\
 = & \quad (\forall \sigma, \sigma'. A \wedge S \Rightarrow P = \forall \sigma, \sigma'. A \wedge S \Rightarrow P) && \text{reflexivity} \\
 = & \quad \top
 \end{aligned}$$

(b) the Postcondition Law: A' is a sufficient postcondition for specification P to be refined by specification S if and only if $A' \Rightarrow P$ is refined by S .

$$\begin{aligned}
 \S & \quad (\forall \sigma'. A' \Rightarrow (\forall \sigma. P \Leftarrow S)) = \forall \sigma, \sigma'. (A' \Rightarrow P) \Leftarrow S && \text{distribution} \\
 = & \quad (\forall \sigma, \sigma'. A' \Rightarrow (P \Leftarrow S)) = \forall \sigma, \sigma'. (A' \Rightarrow P) \Leftarrow S && \text{portation twice} \\
 = & \quad (\forall \sigma, \sigma'. A' \wedge S \Rightarrow P = \forall \sigma, \sigma'. A' \wedge S \Rightarrow P) && \text{reflexivity} \\
 = & \quad \top
 \end{aligned}$$