- 287 (fast string searching)
- (a) Given list P, find list L such that for every index n of list P, Ln is the length of the longest list that is both a proper prefix and a proper suffix of P[0;..n+1]. Here is a program to find L.

$$A \iff i:= 0. \ L:= [\#P*0]. \ j:= 1. \ B$$

$$B \iff \text{if } j \ge \#P \text{ then } ok \text{ else } C. \ L:= j \rightarrow i \mid L. \ j:= j+1. \ B \text{ fi}$$

$$C \iff \text{if } P \ i = P \ j \text{ then } i:= i+1$$

else if $i=0 \text{ then } ok$
else $i:= L \ (i-1). \ C \text{ fi fi}$

Find specifications A, B, and C so that A is the problem and the three refinements are theorems.

(b) Given list S (subject), list P (pattern), and list L (as in part (a)), determine if P is a segment of S, and if so, where it occurs. Here is a program.

 $D \iff m := 0. \ n := 0. \ E$ $E \iff \text{if } m = \#P \text{ then } h := n - \#P \text{ else } F \text{ fi}$ $F \iff \text{if } n = \#S \text{ then } h := \infty$ else if P m = S n then $m := m + 1. \ n := n + 1. \ E$ else G fi fi

$$G \Leftarrow \text{if } m=0 \text{ then } n:=n+1. F \text{ else } m:=L(m-1). G \text{ fi}$$

Find specifications D, E, F, and G so that D is the problem and the four refinements are theorems.

no solution given