

254 (alternate Ackermann) For each of the following functions f , refine $n := f m n$, find a time bound (possibly involving f), and find a space bound.

- (a) $f 0 n = n+2$
 $f 1 0 = 0$
 $f (m+2) 0 = 1$
 $f (m+1) (n+1) = f m (f (m+1) n)$
- (b) $f 0 n = n \times 2$
 $f (m+1) 0 = 1$
 $f (m+1) (n+1) = f m (f (m+1) n)$
- (c) $f 0 n = n+1$
 $f 1 0 = 2$
 $f 2 0 = 0$
 $f (m+3) 0 = 1$
 $f (m+1) (n+1) = f m (f (m+1) n)$

After trying the question, scroll down to the solution.

§ Part (a) defines the same function *ack* as in Exercise 253. Parts (b) and (c) do not, but they define very similar, very closely related functions.

(a) $f\ 0\ n = n+2$
 $f\ 1\ 0 = 0$
 $f\ (m+2)\ 0 = 1$
 $f\ (m+1)\ (n+1) = f\ m\ (f\ (m+1)\ n)$

§ $n := f\ m\ n \iff$
if $m=0$ **then** $n := n+2$
else if $m=1 \wedge n=0$ **then** $n := 0$
else if $n=0$ **then** $n := 1$
else $n := n-1$. $n := f\ m\ n$. $m := m-1$. $n := f\ m\ n$. $m := m+1$ **fi fi fi**

For a time bound, we want a function *g* such that

$t' \leq t + g\ m\ n \wedge n' = f\ m\ n \wedge m' = m \iff$
if $m=0$ **then** $n := n+2$
else if $m=1 \wedge n=0$ **then** $n := 0$
else if $n=0$ **then** $n := 1$
else $n := n-1$. $t := t+1$. $t' \leq t + g\ m\ n \wedge n' = f\ m\ n \wedge m' = m$.
 $m := m-1$. $t' \leq t + g\ m\ n \wedge n' = f\ m\ n \wedge m' = m$. $m := m+1$
fi fi fi

In the last alternative, I put $t := t+1$ before the first recursive call, but not before the second. The one occurrence ensures that every loop includes a time increment. But I could have put another one in. Using Refinement by Cases, and throwing away the unnecessary pieces, we need *f* to satisfy four things.

$t' \leq t + g\ m\ n \iff m=0 \wedge t'=t$
 $t' \leq t + g\ m\ n \iff m=1 \wedge n=0 \wedge t'=t$
 $t' \leq t + g\ m\ n \iff m>1 \wedge n=0 \wedge t'=t$
 $t' \leq t + g\ m\ n \iff m>0 \wedge n>0 \wedge t' \leq t + 1 + g\ m\ (n-1) + g\ (m-1)\ (f\ m\ (n-1))$

Simplifying,

$g\ 0\ n \geq 0$
 $g\ m\ 0 \geq 0$
 $g\ (m+1)\ (n+1) \geq g\ (m+1)\ n + g\ m\ (f\ (m+1)\ n) + 1$

These are the constraints on *g*. So replace \geq by $=$ and we have a definition of *g* that gives the exact execution time (in terms of *f*).

SPACE BOUND NOT DONE YET

(b) $f\ 0\ n = n \times 2$
 $f\ (m+1)\ 0 = 1$
 $f\ (m+1)\ (n+1) = f\ m\ (f\ (m+1)\ n)$

not done

(c) $f\ 0\ n = n+1$
 $f\ 1\ 0 = 2$
 $f\ 2\ 0 = 0$
 $f\ (m+3)\ 0 = 1$
 $f\ (m+1)\ (n+1) = f\ m\ (f\ (m+1)\ n)$

not done