

170 Let s and n be *nat* variables. Here is a refinement.

$$s' = s + 2^n - 1 \iff \mathbf{if\ } n=0 \mathbf{\ then\ } ok \mathbf{\ else\ } n:=n-1. \ s:=s+2^n. \ s' = s + 2^n - 1 \mathbf{\ fi}$$

- (a) Prove it.
- (b) Insert appropriate time increments according to the recursive measure, and write appropriate timing specifications.
- (c) Prove the timing refinement.

After trying the question, scroll down to the solution.

(a) Prove it.

§ By cases. First case:

$$\begin{aligned} & (s' = s + 2^n - 1 \Leftarrow n=0 \wedge ok) && \text{expand } ok \\ = & (s' = s + 2^n - 1 \Leftarrow n=0 \wedge s'=s \wedge n'=n) && \text{context} \\ = & (s = s + 2^0 - 1 \Leftarrow n=0 \wedge s'=s \wedge n'=n) && \text{simplify and specialize} \\ \Rightarrow & \top \end{aligned}$$

Last case, right side:

$$\begin{aligned} & n \neq 0 \wedge (n := n-1. s := s + 2^n. s' = s + 2^n - 1) && \text{substitution law twice} \\ = & n \neq 0 \wedge s' = s + 2^{n-1} + 2^{n-1} - 1 && \text{simplify and specialize} \\ \Rightarrow & s' = s + 2^n - 1 \end{aligned}$$

(b) Insert appropriate time increments according to the recursive measure, and write appropriate timing specifications.

§ $t' = t+n \Leftarrow \mathbf{if } n=0 \mathbf{ then } ok \mathbf{ else } n := n-1. s := s + 2^n. t := t+1. t' = t+n \mathbf{ fi}$

(c) Prove the timing refinement.

§ By cases. First case:

$$\begin{aligned} & (t' = t+n \Leftarrow n=0 \wedge ok) && \text{expand } ok \\ = & (t' = t+n \Leftarrow n=0 \wedge s'=s \wedge n'=n \wedge t'=t) && \text{context} \\ = & (t = t+0 \Leftarrow n=0 \wedge s'=s \wedge n'=n \wedge t'=t) && \text{simplify and specialize} \\ \Rightarrow & \top \end{aligned}$$

Last case, right side:

$$\begin{aligned} & n \neq 0 \wedge (n := n-1. s := s + 2^n. t := t+1. t' = t+n) && \text{substitution law 3 times} \\ = & n \neq 0 \wedge t' = t+1 + n-1 && \text{simplify and specialize} \\ \Rightarrow & t' = t+n \end{aligned}$$