

135 Let  $x$  and  $n$  be natural variables. Find a specification  $P$  such that both the following refinements can be proven:

$$x = x' \times 2^{n'} \Leftarrow n := 0. P$$

$$P \Leftarrow \mathbf{if\ even\ } x \mathbf{\ then\ } x := x/2. \ n := n+1. \ P \mathbf{\ else\ ok\ fi}$$

After trying the question, scroll down to the solution.

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$$P = x = x' \times 2^{n'-n}$$

Proof:

$$n:=0. x = x' \times 2^{n'-n}$$

Substitution Law

$$= x = x' \times 2^{n'}$$

The second one is proven by cases. First,

$$\text{even } x \wedge (x:=x/2. n:=n+1. x = x' \times 2^{n'-n})$$

use Substitution Law twice

$$= \text{even } x \wedge x/2 = x' \times 2^{n'-(n+1)}$$

number theory

$$= \text{even } x \wedge x = x' \times 2^{n'-n}$$

specialization

$$\Rightarrow P$$

Now the other case:

$$(x = x' \times 2^{n'-n} \Leftarrow \neg \text{even } x \wedge \text{ok})$$

expand *ok*

$$= (x = x' \times 2^{n'-n} \Leftarrow \neg \text{even } x \wedge x'=x \wedge n'=n)$$

use antecedent as context in consequent

$$= (x = x \times 2^{n'-n} \Leftarrow \neg \text{even } x \wedge x'=x \wedge n'=n)$$

number theory

$$= (\top \Leftarrow \neg \text{even } x \wedge x'=x \wedge n'=n)$$

base law for  $\Leftarrow$

$$= \top$$