a Practical Theory of Programming

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The cover picture is an inukshuk, which is a human-like figure made of piled stones. Inukshuks are found throughout arctic Canada. They are built by the Inuit people, who use them to mean “You are on the right path.”.

11.3 Laws

11.3.0 Binary

Let a , b , c , d , and e be binary.

Binary
$$\top$$
$$\neg\bot$$

Excluded Middle (Tertium non Datur)
$$a \vee \neg a$$

Noncontradiction
$$\neg(a \wedge \neg a)$$

Base
$$\neg(a \wedge \bot)$$
$$a \vee \top$$
$$a \Rightarrow \top$$
$$\bot \Rightarrow a$$

Identity
$$\top \wedge a = a$$
$$\bot \vee a = a$$
$$\top \Rightarrow a = a$$
$$\top = a = a$$

Idempotent
$$a \wedge a = a$$
$$a \vee a = a$$

Reflexive
$$a \Rightarrow a$$
$$a = a$$

Indirect Proof
$$\neg a \Rightarrow \bot = a \quad \text{(Reductio ad Absurdum)}$$
$$\neg a \Rightarrow a = a$$

Specialization
$$a \wedge b \Rightarrow a$$

Associative
$$a \wedge (b \wedge c) = (a \wedge b) \wedge c$$
$$a \vee (b \vee c) = (a \vee b) \vee c$$
$$a = (b = c) = (a = b) = c$$
$$a \neq (b \neq c) = (a \neq b) \neq c$$
$$a = (b \neq c) = (a = b) \neq c$$

Mirror
$$a \Leftarrow b = b \Rightarrow a$$

Double Negation
$$\neg\neg a = a$$

Duality (deMorgan)
$$\neg(a \wedge b) = \neg a \vee \neg b$$
$$\neg(a \vee b) = \neg a \wedge \neg b$$

Exclusion
$$a \Rightarrow \neg b = b \Rightarrow \neg a \quad \text{(Contrapositive)}$$
$$a = \neg b = a \neq b = \neg a = b$$

Inclusion
$$a \Rightarrow b = \neg a \vee b \text{ (Material Implication)}$$
$$a \Rightarrow b = (a \wedge b = a)$$
$$a \Rightarrow b = (a \vee b = b)$$

Absorption
$$a \wedge (a \vee b) = a$$
$$a \vee (a \wedge b) = a$$

Direct Proof
$$(a \Rightarrow b) \wedge a \Rightarrow b \quad \text{(Modus Ponens)}$$
$$(a \Rightarrow b) \wedge \neg b \Rightarrow \neg a \quad \text{(Modus Tollens)}$$
$$(a \vee b) \wedge \neg a \Rightarrow b \text{(Disjunctive Syllogism)}$$

Transitive
$$(a \wedge b) \wedge (b \wedge c) \Rightarrow (a \wedge c)$$
$$(a \Rightarrow b) \wedge (b \Rightarrow c) \Rightarrow (a \Rightarrow c)$$
$$(a = b) \wedge (b = c) \Rightarrow (a = c)$$
$$(a \Rightarrow b) \wedge (b = c) \Rightarrow (a \Rightarrow c)$$
$$(a = b) \wedge (b \Rightarrow c) \Rightarrow (a \Rightarrow c)$$

Distributive (Factoring)
$$a \wedge (b \wedge c) = (a \wedge b) \wedge (a \wedge c)$$
$$a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$$
$$a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$$
$$a \vee (b \vee c) = (a \vee b) \vee (a \vee c)$$
$$a \vee (b \Rightarrow c) = (a \vee b) \Rightarrow (a \vee c)$$
$$a \vee (b = c) = (a \vee b) = (a \vee c)$$
$$a \Rightarrow (b \wedge c) = (a \Rightarrow b) \wedge (a \Rightarrow c)$$
$$a \Rightarrow (b \vee c) = (a \Rightarrow b) \vee (a \Rightarrow c)$$
$$a \Rightarrow (b \Rightarrow c) = (a \Rightarrow b) \Rightarrow (a \Rightarrow c)$$
$$a \Rightarrow (b = c) = (a \Rightarrow b) = (a \Rightarrow c)$$

Symmetry (Commutative)

$a \wedge b = b \wedge a$
$a \vee b = b \vee a$
$a = b = b = a$
$a \neq b = b \neq a$

Antisymmetry (Double Implication)

$(a \Rightarrow b) \wedge (b \Rightarrow a) = a = b$

Discharge

$a \wedge (a \Rightarrow b) = a \wedge b$
$a \Rightarrow (a \wedge b) = a \Rightarrow b$

Antimonotonic

$a \Rightarrow b \Longrightarrow (b \Rightarrow c) \Rightarrow (a \Rightarrow c)$

Monotonic

$a \Rightarrow b \Longrightarrow c \wedge a \Rightarrow c \wedge b$
$a \Rightarrow b \Longrightarrow c \vee a \Rightarrow c \vee b$
$a \Rightarrow b \Longrightarrow (c \Rightarrow a) \Rightarrow (c \Rightarrow b)$

Generalization

$a \Rightarrow a \vee b$

Antidistributive

$a \wedge b \Rightarrow c = (a \Rightarrow c) \vee (b \Rightarrow c)$
$a \vee b \Rightarrow c = (a \Rightarrow c) \wedge (b \Rightarrow c)$

Portation

$a \wedge b \Rightarrow c = a \Rightarrow (b \Rightarrow c)$
$a \wedge b \Rightarrow c = a \Rightarrow \neg b \vee c$

Conflation

$(a \Rightarrow b) \wedge (c \Rightarrow d) \Longrightarrow a \wedge c \Rightarrow b \wedge d$
$(a \Rightarrow b) \wedge (c \Rightarrow d) \Longrightarrow a \vee c \Rightarrow b \vee d$

Contrapositive

$a \Rightarrow b = \neg b \Rightarrow \neg a$

Equality and Difference

$a=b = (a \wedge b) \vee (\neg a \wedge \neg b)$
$a \neq b = (a \wedge \neg b) \vee (\neg a \wedge b)$

Resolution

$a \wedge c \Longrightarrow (a \vee b) \wedge (\neg b \vee c) = (a \wedge \neg b) \vee (b \wedge c) \Longrightarrow a \vee c$

Case Creation

a = **if** b **then** $b \Rightarrow a$ **else** $\neg b \Rightarrow a$ **fi**
a = **if** b **then** $b \wedge a$ **else** $\neg b \wedge a$ **fi**
a = **if** b **then** $b = a$ **else** $b \neq a$ **fi**

Case Analysis

if a **then** b **else** c **fi** $= (a \wedge b) \vee (\neg a \wedge c)$
if a **then** b **else** c **fi** $= (a \Rightarrow b) \wedge (\neg a \Rightarrow c)$

One Case

if a **then** $\top$ **else** b **fi** $= a \vee b$
if a **then** $\bot$ **else** b **fi** $= \neg a \wedge b$
if a **then** b **else** $\top$ **fi** $= a \Rightarrow b$
if a **then** b **else** $\bot$ **fi** $= a \wedge b$
if a **then** b **else** $\neg b$ **fi** $= a = b$
if a **then** $\neg b$ **else** b **fi** $= a \neq b$

Case Absorption

if a **then** b **else** c **fi** = **if** a **then** $a \wedge b$ **else** c **fi**
if a **then** b **else** c **fi** = **if** a **then** $a \Rightarrow b$ **else** c **fi**
if a **then** b **else** c **fi** = **if** a **then** $a = b$ **else** c **fi**
if a **then** b **else** c **fi** = **if** a **then** b **else** $\neg a \wedge c$ **fi**
if a **then** b **else** c **fi** = **if** a **then** b **else** $a \vee c$ **fi**
if a **then** b **else** c **fi** = **if** a **then** b **else** $a \neq c$ **fi**

Case Distributive (Case Factoring)

$\neg$ **if** a **then** b **else** c **fi** = **if** a **then** $\neg b$ **else** $\neg c$ **fi**
if a **then** b **else** c **fi** $\wedge\ d$ = **if** a **then** $b \wedge d$ **else** $c \wedge d$ **fi**
and similarly replacing $\wedge$ by any of $\vee = \neq \Rightarrow \Leftarrow$
if a **then** $b \wedge c$ **else** $d \wedge e$ **fi** = **if** a **then** b **else** d **fi** $\wedge$ **if** a **then** c **else** e **fi**
and similarly replacing $\wedge$ by any of $\vee = \neq \Rightarrow \Leftarrow$

End of Binary

11.3.1 Generic

The operators $=$ $\neq$ **if then else fi** apply to every type of expression (but the first operand of **if then else fi** must be binary), with the laws

$x{=}x$	reflexivity	**if** $\top$ **then** x **else** y **fi** $= x$	case base
$x{=}y \;=\; y{=}x$	symmetry	**if** $\bot$ **then** x **else** y **fi** $= y$	case base
$x{=}y \wedge y{=}z \Rightarrow x{=}z$	transitivity	**if** a **then** x **else** x **fi** $= x$	case idempotent
$x{=}y \Rightarrow f\,x = f\,y$	transparency	**if** a **then** x **else** y **fi** $=$ **if** $\neg a$ **then** y **else** x **fi**	
$x{\neq}y \;=\; \neg(x{=}y)$	unequality		case reversal

The operators $\uparrow$ $\downarrow$ $<$ $\leq$ $>$ $\geq$ apply to numbers, characters, strings, and lists, with the laws

$x{\leq}x$	reflexivity	$\neg\, x{<}x$	irreflexivity
$\neg(x{<}y \wedge x{=}y)$	exclusivity	$\neg(x{>}y \wedge x{=}y)$	exclusivity
$\neg(x{<}y \wedge x{>}y)$	exclusivity	$x{\leq}y \;=\; x{<}y \vee x{=}y$	inclusivity
$x{\leq}y \wedge y{\leq}z \Rightarrow x{\leq}z$	transitivity	$x{<}y \wedge y{\leq}z \Rightarrow x{<}z$	transitivity
$x{<}y \wedge y{<}z \Rightarrow x{<}z$	transitivity	$x{\leq}y \wedge y{<}z \Rightarrow x{<}z$	transitivity
$x{>}y \;=\; y{<}x$	mirror	$x{\geq}y \;=\; y{\leq}x$	mirror
$\neg\, x{<}y \;=\; x{\geq}y$	totality	$\neg\, x{\leq}y \;=\; x{>}y$	totality
$x{\leq}y \wedge y{\leq}x = x{=}y$	antisymmetry	$x{<}y \vee x{=}y \vee x{>}y$	totality, trichotomy
$x{\uparrow}x = x$	idempotence	$x{\downarrow}x = x$	idempotence
$x{\uparrow}y = y{\uparrow}x$	symmetry	$x{\downarrow}y = y{\downarrow}x$	symmetry
$x{\uparrow}(y{\uparrow}z) = (x{\uparrow}y){\uparrow}z$	associativity	$x{\downarrow}(y{\downarrow}z) = (x{\downarrow}y){\downarrow}z$	associativity
$x{\uparrow}(y{\downarrow}z) = (x{\uparrow}y){\downarrow}(x{\uparrow}z)$	distributivity	$x{\downarrow}(y{\uparrow}z) = (x{\downarrow}y){\uparrow}(x{\downarrow}z)$	distributivity
$x{\uparrow}y \;=$ **if** $x{\geq}y$ **then** x **else** y **fi**		$x{\downarrow}y \;=$ **if** $x{\leq}y$ **then** x **else** y **fi**	
$x{\downarrow}y \leq x \leq x{\uparrow}y$			

End of Generic

11.3.2 Numbers

Let d be a sequence of (zero or more) digits, and let x, y, and z be numbers.

$d0{+}1 = d1$	$d5{+}1 = d6$	counting
$d1{+}1 = d2$	$d6{+}1 = d7$	counting
$d2{+}1 = d3$	$d7{+}1 = d8$	counting
$d3{+}1 = d4$	$d8{+}1 = d9$	counting
$d4{+}1 = d5$	$d9{+}1 = (d{+}1)0$	counting (see Exercise 32)
$x{+}0 = x$		identity
$x{+}y = y{+}x$		symmetry
$x{+}(y{+}z) \;=\; (x{+}y){+}z$		associativity
$-\infty{<}x{<}\infty \Rightarrow (x{+}y = x{+}z \;=\; y{=}z)$		cancellation
$-\infty{<}x \Rightarrow\; \infty{+}x = \infty$		absorption
$x{<}\infty \;\Rightarrow\; -\infty + x = -\infty$		absorption
$-x \;=\; 0{-}x$		negation
$-\,-x = x$		self-inverse
$-(x{+}y) = -x + -y$		distributivity
$-(x{-}y) = y{-}x$		antisymmetry
$-(x{\times}y) = -x \times y$		semi-distributivity
$-(x/y) = -x \,/\, y$		semi-distributivity
$x{-}0 = x$		identity
$x{-}y \;=\; x + -y$		subtraction
$x{+}(y{-}z) = (x{+}y){-}z$		associativity
$-\infty{<}x{<}\infty \;\Rightarrow\; (x{-}y = x{-}z \;=\; y{=}z)$		cancellation
$-\infty{<}x{<}\infty \;\Rightarrow\; x{-}x = 0$		inverse
$x{<}\infty \;\Rightarrow\; \infty{-}x = \infty$		absorption
$-\infty{<}x \;\Rightarrow\; -\infty - x = -\infty$		absorption

$-\infty<x<\infty \Rightarrow x\times 0 = 0$		base
$x\times 1 = x$		identity
$x\times y = y\times x$		symmetry
$x\times(y+z) = x\times y + x\times z$		distributivity
$x\times(y\times z) = (x\times y)\times z$		associativity
$-\infty<x<\infty \wedge x\neq 0 \Rightarrow (x\times y = x\times z = y=z)$		cancellation
$0<x \Rightarrow x\times\infty = \infty$		absorption
$0<x \Rightarrow x\times -\infty = -\infty$		absorption
$x/1 = x$		identity
$-\infty<x<\infty \wedge x\neq 0 \Rightarrow 0/x = 0$		base
$-\infty<x<\infty \wedge x\neq 0 \Rightarrow x/x = 1$		base
$z\neq 0 \Rightarrow x\times(y/z) = (x\times y)/z = x/(z/y)$		multiplication-division
$y\neq 0 \Rightarrow (x/y)/z = x/(y\times z)$		multiplication-division
$-\infty<y<\infty \wedge y\neq 0 \Rightarrow (x/y)\times y = x$		multiplication-division
$-\infty<x<\infty \Rightarrow x/\infty = 0 = x/-\infty$		annihilation
$-\infty<x<\infty \Rightarrow x^0 = 1$		base
$x^1 = x$		identity
$x^{y+z} = x^y \times x^z$		exponents
$-\infty<0<1<\infty$		direction
$x<y = -y<-x$		reflection
$-\infty<x<\infty \Rightarrow (x+y < x+z = y<z)$		cancellation, translation
$0<x<\infty \Rightarrow (x\times y < x\times z = y<z)$		cancellation, scale
$x<y \vee x=y \vee x>y$		trichotomy
$-\infty \le x \le \infty$		extremes
$x\uparrow\infty = \infty$	$x\downarrow -\infty = -\infty$	base
$x\uparrow -\infty = x$	$x\downarrow\infty = x$	identity
$-(x\uparrow y) = -x\downarrow -y$	$-(x\downarrow y) = -x\uparrow -y$	duality
$x + y\uparrow z = (x+y)\uparrow(x+z)$	$x - y\uparrow z = (x-y)\downarrow(x-z)$	distributivity
$x\ge 0 \Rightarrow x\times(y\uparrow z) = (x\times y)\uparrow(x\times z)$	$x\ge 0 \Rightarrow x\times(y\downarrow z) = (x\times y)\downarrow(x\times z)$	distributivity
$x\le 0 \Rightarrow x\times(y\uparrow z) = (x\times y)\downarrow(x\times z)$	$x\ge 0 \Rightarrow x\times(y\downarrow z) = (x\times y)\uparrow(x\times z)$	distributivity

End of Numbers

11.3.3 Bunches

Let x and y be elements (binaries, numbers, characters, sets, strings and lists of elements).

$x: y = x=y$	elementary
$x: A, B = x: A \vee x: B$	compound
$A, A = A$	idempotence
$A, B = B, A$	symmetry
$A, (B, C) = (A, B), C$	associativity
$A`A = A$	idempotence
$A`B = B`A$	symmetry
$A`(B`C) = (A`B)`C$	associativity
$A, B: C = A: C \wedge B: C$	antidistributivity
$A: B`C = A: B \wedge A: C$	distributivity
$A: A, B$	generalization
$A`B: A$	specialization
$A: A$	reflexivity
$A: B \wedge B: A = A=B$	antisymmetry
$A: B \wedge B: C \Rightarrow A: C$	transitivity

$¢\ null\ =\ 0$	size
$¢\ x\ =\ 1$	size
$¢\ nat\ =\ \infty$	size
$¢(A, B) + ¢(A`B)\ =\ ¢A + ¢B$	size
$\neg\ x{:}\ A\ =\!=\ ¢(A`x) = 0$	size
$A{:}\ B\ \Rightarrow\ ¢A \le ¢B$	size
$A, (A`B)\ =\ A$	absorption
$A`(A, B)\ =\ A$	absorption
$A{:}\ B\ =\!=\ A, B = B\ =\!=\ A = A`B$	inclusion
$A, (B, C)\ =\ (A, B), (A, C)$	distributivity
$A,(B`C)\ =\ (A,B)`(A,C)$	distributivity
$A`(B, C)\ =\ (A`B), (A`C)$	distributivity
$A`(B`C)\ =\ (A`B)`(A`C)$	distributivity
$A{:}\ B\ \wedge\ C{:}\ D\ \Rightarrow\ A,C{:}\ B,D$	conflation, monotonicity
$A{:}\ B\ \wedge\ C{:}\ D\ \Rightarrow\ A`C{:}\ B`D$	conflation, monotonicity
$null{:}\ A$	induction
$A, null = A$	identity
$A\ `\ null = null$	base
$¢A = 0\ =\!=\ A = null$	size
$x, y{:}\ xint\ \wedge\ x \le y\ \Rightarrow\ (i{:}\ x,..y\ =\!=\ i{:}\ xint\ \wedge\ x \le i < y)$	interval
$x, y{:}\ xint\ \wedge\ x \le y\ \Rightarrow\ ¢(x,..y)\ =\ y – x$	interval
$nat = 0,..\infty$	interval
$–null\ =\ null$	distribution
$–(A, B)\ =\ –A, –B$	distribution
$A+null\ =\ null+A\ =\ null$	distribution
$(A, B)+(C, D)\ =\ A+C, A+D, B+C, B+D$	distribution

and similarly for many other operators (see the final page of the book)

End of Bunches

11.3.4 Sets

Let S be a set, and let A and B be anything.

$\{\sim S\}\ =\ S$	$\{A\}{:}\ ♯B\ =\!=\ A{:}\ B$
$\sim\{A\}\ =\ A$	$\$\{A\} = ¢A$
$\{A\}\ ≠\ A$	$\{A\} \cup \{B\}\ =\ \{A, B\}$
$A \in \{B\}\ =\!=\ A{:}\ B$	$\{A\} \cap \{B\}\ =\ \{A\ `\ B\}$
$\{A\} \subseteq \{B\}\ =\!=\ A{:}\ B$	$\{A\} = \{B\}\ =\!=\ A = B$
	$\{A\} ≠ \{B\}\ =\!=\ A ≠ B$

End of Sets

11.3.5 Strings

Let S, T, and U be strings; let i and j be items (binary values, numbers, characters, sets, lists, functions); let n and m be extended natural; let x, y, and z be extended integers, $x \le y \le z$; let A and B be bunches (anything).

$S; nil\ =\ S\ =\ nil; S$	$S_{(T_U)}\ =\ (S_T)_U$
$S; (T; U)\ =\ (S; T); U$	$S_{nil}\ =\ nil$
$↔nil\ =\ 0$	$S_{T;\,U}\ =\ S_T; S_U$
$↔i\ =\ 1$	$S_{\{A\}}\ =\ \{S_A\}$
$↔(S; T)\ =\ ↔S + ↔T$	$↔S < \infty\ \Rightarrow\ nil \le S < S; i; T$
$¢\ nil\ =\ 1$	$↔S < \infty\ \Rightarrow\ (i < j\ \Rightarrow\ S; i; T < S; j; U)$

$\not\subset(A; B) \le \not\subset A \times \not\subset B$
$\leftrightarrow S{<}\infty \Rightarrow (S; i; T)_{\leftrightarrow S} = i$
$\leftrightarrow S{<}\infty \Rightarrow S;i;T \triangleleft \leftrightarrow S \triangleright j = S;j;T$
$0{*}S = nil$
$(n{+}1){*}S = n{*}S; S$
$*S = **S = nat{*}S$

$\leftrightarrow S{<}\infty \Rightarrow (i{=}j = S; i; T = S; j; T)$
$(S \triangleleft n \triangleright i)_m =$ **if** $n{=}m$ **then** i **else** S_m **fi**
$x;..x = nil$
$x;..x{+}1 = x$
$(x;..y) \ ; \ (y;..z) = x;..z$
$\leftrightarrow(x;..y) = y–x$

End of Strings

11.3.6 Lists

Let S and T be strings; let i be an item (binary value, number, character, set, list, function); let L, M, and N be lists; let n and m be extended natural.

$[S] + S = \sim[S]$
$[\sim L] = L$
$[S];;[T] = [S; T]$
$[S] = [T] = S = T$
$[S] < [T] = S < T$
$[A]: [B] = A: B$
$\#[S] = \leftrightarrow S$
$nil \rightarrow i \mid L = i$
$n \rightarrow i \mid [S] = [S \triangleleft n \triangleright i]$
$(n \rightarrow i \mid L)m =$ **if** $n{=}m$ **then** i **else** $L\ m$ **fi**
$(S;T) \rightarrow i \mid L = S \rightarrow (T \rightarrow i \mid L@S) \mid L$

$\Box L = 0,..\#L$
$[S]\ T = S_T$
$S_{[T]} = [S_T]$
$[S]\ [T] = [S_T]$
$L\ \{A\} = \{L\ A\}$
$L\ [S] = [L\ S]$
$(L\ M)\ N = L\ (M\ N)$
$\#L = \not\subset \Box L$
$L@nil = L$
$L@i = L\ i$
$L@(S; T) = L@S@T$

End of Lists

11.3.7 Functions

Renaming — if v and w do not appear in D and w does not appear in b
$\langle v: D \rightarrow b \rangle = \langle w: D \rightarrow \langle v: D \rightarrow b \rangle\ w \rangle$

Domain
$\Box \langle v: D \rightarrow b \rangle = D$

Application — if element $x: D$
$\langle v: D \rightarrow b \rangle\ x =$ (substitute x for v in b)

Size
$\#f = \not\subset \Box f$

Functional Union
$\Box(f, g) = \Box f \ ` \ \Box g$
$(f, g)\ x = f\ x, g\ x$

Distributive
$f\ null = null$
$f(A, B) = f\ A, f\ B$
$f(\S g) = \S y: f(\Box g) \cdot \exists x: \Box g \cdot f\ x = y \ \wedge \ g\ x$
f **if** b **then** x **else** y **fi** = **if** b **then** $f\ x$ **else** $f\ y$ **fi**
if b **then** f **else** g **fi** $x =$ **if** b **then** $f\ x$ **else** $g\ x$ **fi**

Function Composition — if $\neg f: \Box g$
$\Box(g\ f) = \S x: \Box f \cdot f\ x: \Box g$
$(g\ f)\ x = g\ (f\ x)$
$f\ (g\ h) = (f\ g)\ h$

Functional Intersection
$\Box(f \ ` g) = \Box f, \Box g$
$(f \ ` g)x = (f \mid g)x \ ` \ (g \mid f)x$

Arrow
$f: null \rightarrow A$
$A \rightarrow B: \ (A ` C) \rightarrow (B, D)$
$(A, B) \rightarrow C = A \rightarrow C \mid B \rightarrow C$
$f: A \rightarrow B = A: \Box f \ \wedge \ \forall a: A \cdot f\ a: B$

Selective Union
$\Box(f \mid g) = \Box f, \Box g$
$(f \mid g)x =$ **if** $x: \Box f$ **then** $f\ x$ **else** $g\ x$ **fi**
$f \mid f = f$
$f \mid (g \mid h) = (f \mid g) \mid h$
$(g \mid h)f = g\ f \mid h\ f$

Extension

$f \;=\; \langle v{:}\ \Box f \rightarrow f\, v\rangle$

Function Inclusion and Equality

$f{:}\ g \;=\; \Box g{:}\ \Box f \wedge \forall x{:}\ \Box g{\cdot}\ f\, x{:}\ g\, x$

$f = g \;=\; \Box f = \Box g \wedge \forall x{:}\ \Box f{\cdot}\ f\, x = g\, x$

End of Functions

11.3.8 Quantifiers

Let x be an element, let a, b and c be binary, let n and m be numeric, let f and g be functions, and let p be a predicate.

$\forall v{:}\ null{\cdot}\ b \;=\; \top$
$\forall v{:}\ x{\cdot}\ b \;=\; \langle v{:}\ x \rightarrow b\rangle\, x$

$\forall v{:}\ A, B{\cdot}\ b \;=\; (\forall v{:}\ A{\cdot}\ b) \wedge (\forall v{:}\ B{\cdot}\ b)$
$\forall v{:}\ (\S v{:}\ D{\cdot}\ b){\cdot}\ c \;=\; \forall v{:}\ D{\cdot}\ b \Rightarrow c$

$\exists v{:}\ null{\cdot}\ b \;=\; \bot$
$\exists v{:}\ x{\cdot}\ b \;=\; \langle v{:}\ x \rightarrow b\rangle\, x$

$\exists v{:}\ A, B{\cdot}\ b \;=\; (\exists v{:}\ A{\cdot}\ b) \vee (\exists v{:}\ B{\cdot}\ b)$
$\exists v{:}\ (\S v{:}\ D{\cdot}\ b){\cdot}\ c \;=\; \exists v{:}\ D{\cdot}\ b \wedge c$

$\Sigma v{:}\ null{\cdot}\ n \;=\; 0$
$\Sigma v{:}\ x{\cdot}\ n \;=\; \langle v{:}\ x \rightarrow n\rangle\, x$

$(\Sigma v{:}\ A, B{\cdot}\ n) + (\Sigma v{:}\ A\text{`}B{\cdot}\ n) \;=\; (\Sigma v{:}\ A{\cdot}\ n) + (\Sigma v{:}\ B{\cdot}\ n)$
$\Sigma v{:}\ (\S v{:}\ D{\cdot}\ b){\cdot}\ n \;=\; \Sigma v{:}\ D{\cdot}$ **if** b **then** n **else** 0 **fi**

$\Pi v{:}\ null{\cdot}\ n \;=\; 1$
$\Pi v{:}\ x{\cdot}\ n \;=\; \langle v{:}\ x \rightarrow n\rangle\, x$

$(\Pi v{:}\ A, B{\cdot}\ n) \times (\Pi v{:}\ A\text{`}B{\cdot}\ n) \;=\; (\Pi v{:}\ A{\cdot}\ n) \times (\Pi v{:}\ B{\cdot}\ n)$
$\Pi v{:}\ (\S v{:}\ D{\cdot}\ b){\cdot}\ n \;=\; \Pi v{:}\ D{\cdot}$ **if** b **then** n **else** 1 **fi**

$\Downarrow v{:}\ null{\cdot}\ n \;=\; \infty$
$\Downarrow v{:}\ x{\cdot}\ n \;=\; \langle v{:}\ x \rightarrow n\rangle\, x$

$\Downarrow v{:}\ A, B{\cdot}\ n \;=\; (\Downarrow v{:}\ A{\cdot}\ n) \downarrow (\Downarrow v{:}\ B{\cdot}\ n)$
$\Downarrow v{:}\ (\S v{:}\ D{\cdot}\ b){\cdot}\ n \;=\; \Downarrow v{:}\ D{\cdot}$ **if** b **then** n **else** ∞ **fi**

$\Uparrow v{:}\ null{\cdot}\ n \;=\; -\infty$
$\Uparrow v{:}\ x{\cdot}\ n \;=\; \langle v{:}\ x \rightarrow n\rangle\, x$

$\Uparrow v{:}\ A, B{\cdot}\ n \;=\; (\Uparrow v{:}\ A{\cdot}\ n) \uparrow (\Uparrow v{:}\ B{\cdot}\ n)$
$\Uparrow v{:}(\S v{:}\ D{\cdot}\ b){\cdot}\ n \;=\; \Uparrow v{:}\ D{\cdot}$ **if** b **then** n **else** $-\infty$ **fi**

$\S v{:}\ null{\cdot}\ b \;=\; null$
$\S v{:}\ x{\cdot}\ b \;=\;$ **if** $\langle v{:}\ x \rightarrow b\rangle\, x$ **then** x **else** $null$ **fi**
$\S v{:}\ A,B{\cdot}\ b \;=\; (\S v{:}\ A{\cdot}\ b), (\S v{:}\ B{\cdot}\ b)$
$\S v{:}\ A\text{`}B{\cdot}\ b \;=\; (\S v{:}\ A{\cdot}\ b)\ \text{`}\ (\S v{:}\ B{\cdot}\ b)$
$\S v{:}\ (\S v{:}\ D{\cdot}\ b){\cdot}\ c \;=\; \S v{:}\ D{\cdot}\ b \wedge c$

Inclusion

$A{:}\ B \;=\; \forall x{:}\ A{\cdot}\ x{:}\ B$

Cardinality

$\not\subset A \;=\; \Sigma\, (A \rightarrow 1)$

Identity

$\forall v{\cdot}\ \top$
$\neg \exists v{\cdot}\ \bot$

Change of Variable — if d does not appear in b

$\forall r{:}\ f\, D{\cdot}\ b \;=\; \forall d{:}\ D{\cdot}\ \langle r{:}\ f\, D \rightarrow b\rangle\, (f\, d)$
$\exists r{:}\ f\, D{\cdot}\ b \;=\; \exists d{:}\ D{\cdot}\ \langle r{:}\ f\, D \rightarrow b\rangle\, (f\, d)$
$\Downarrow r{:}\ f\, D{\cdot}\ n \;=\; \Downarrow d{:}\ D{\cdot}\ \langle r{:}\ f\, D \rightarrow n\rangle\, (f\, d)$
$\Uparrow r{:}\ f\, D{\cdot}\ n \;=\; \Uparrow d{:}\ D{\cdot}\ \langle r{:}\ f\, D \rightarrow n\rangle\, (f\, d)$

Specialize and Generalize — if element $x{:}\ \Box f$

$\Downarrow f \le f\, x \le \Uparrow f$

Bunch-Element Conversion

$A{:}\ B \;=\; \forall a{:}\ A{\cdot}\ \exists b{:}\ B{\cdot}\ a{=}b$
$f\, A{:}\ g\, B \;=\; \forall a{:}\ A{\cdot}\ \exists b{:}\ B{\cdot}\ f\, a = g\, b$

Distributive — if $D{\neq}null$ and v does not appear in a

$a \wedge \forall v{:}\ D{\cdot}\ b \;=\; \forall v{:}\ D{\cdot}\ a \wedge b$
$a \wedge \exists v{:}\ D{\cdot}\ b \;=\; \exists v{:}\ D{\cdot}\ a \wedge b$
$a \vee \forall v{:}\ D{\cdot}\ b \;=\; \forall v{:}\ D{\cdot}\ a \vee b$
$a \vee \exists v{:}\ D{\cdot}\ b \;=\; \exists v{:}\ D{\cdot}\ a \vee b$
$a \Rightarrow \forall v{:}\ D{\cdot}\ b \;=\; \forall v{:}\ D{\cdot}\ a \Rightarrow b$
$a \Rightarrow \exists v{:}\ D{\cdot}\ b \;=\; \exists v{:}\ D{\cdot}\ a \Rightarrow b$

Idempotent — if $D{\neq}null$ and v does not appear in b

$\forall v{:}\ D{\cdot}\ b \;=\; b$
$\exists v{:}\ D{\cdot}\ b \;=\; b$

Absorption — if x: D

$\langle v\colon D \to b\rangle\, x \wedge \exists v\colon D\cdot b = \langle v\colon D \to b\rangle\, x$
$\langle v\colon D \to b\rangle\, x \vee \forall v\colon D\cdot b = \langle v\colon D \to b\rangle\, x$
$\langle v\colon D \to b\rangle\, x \wedge \forall v\colon D\cdot b = \forall v\colon D\cdot b$
$\langle v\colon D \to b\rangle\, x \vee \exists v\colon D\cdot b = \exists v\colon D\cdot b$

Specialization — if element x: $\square p$

$\forall p \implies p\, x$

One-Point — if x: D
and v does not appear in x

$\forall v\colon D\cdot v{=}x \Rightarrow b = \langle v\colon D \to b\rangle\, x$
$\exists v\colon D\cdot v{=}x \wedge b = \langle v\colon D \to b\rangle\, x$

Duality

$\neg \forall v\cdot b = \exists v\cdot \neg b$ (deMorgan)
$\neg \exists v\cdot b = \forall v\cdot \neg b$ (deMorgan)
$-\Uparrow v\cdot n = \Downarrow v\cdot -n$
$-\Downarrow v\cdot n = \Uparrow v\cdot -n$

Solution

$\S v\colon D\cdot \top = D$
$(\S v\colon D\cdot b)\colon D$
$\S v\colon D\cdot \bot = null$
$(\S v\cdot b)\colon (\S v\cdot c) = \forall v\cdot b{\Rightarrow}c$
$(\S v\cdot b), (\S v\cdot c) = \S v\cdot b \vee c$
$(\S v\cdot b)\ `\ (\S v\cdot c) = \S v\cdot b \wedge c$
$x\colon \S p = x\colon \square p \wedge p\, x$
$\forall f = (\S f){=}(\square f)$
$\exists f = (\S f){\neq}null$

Bounding — if $D{\neq}null$
and v does not appear in n

$n > (\Uparrow v\colon D\cdot m) \Rightarrow (\forall v\colon D\cdot n{>}m)$
$n < (\Downarrow v\colon D\cdot m) \Rightarrow (\forall v\colon D\cdot n{<}m)$
$n \geq (\Uparrow v\colon D\cdot m) = (\forall v\colon D\cdot n{\geq}m)$
$n \leq (\Downarrow v\colon D\cdot m) = (\forall v\colon D\cdot n{\leq}m)$
$n \geq (\Downarrow v\colon D\cdot m) \Leftarrow (\exists v\colon D\cdot n{\geq}m)$
$n \leq (\Uparrow v\colon D\cdot m) \Leftarrow (\exists v\colon D\cdot n{\leq}m)$
$n > (\Downarrow v\colon D\cdot m) = (\exists v\colon D\cdot n{>}m)$
$n < (\Uparrow v\colon D\cdot m) = (\exists v\colon D\cdot n{<}m)$

Distributive — if $D{\neq}null$ and v does not appear in n

$n{\uparrow}(\Uparrow v\colon D\cdot m) = (\Uparrow v\colon D\cdot n{\uparrow}m)$
$n{\uparrow}(\Downarrow v\colon D\cdot m) = (\Downarrow v\colon D\cdot n{\uparrow}m)$
$n + (\Uparrow v\colon D\cdot m) = (\Uparrow v\colon D\cdot n{+}m)$
$n - (\Uparrow v\colon D\cdot m) = (\Downarrow v\colon D\cdot n{-}m)$
$(\Uparrow v\colon D\cdot m) - n = (\Uparrow v\colon D\cdot m{-}n)$
$n{\geq}0 \Rightarrow n{\times}(\Uparrow v\colon D\cdot m) = (\Uparrow v\colon D\cdot n{\times}m)$
$n{\leq}0 \Rightarrow n{\times}(\Uparrow v\colon D\cdot m) = (\Downarrow v\colon D\cdot n{\times}m)$
$n{\times}(\Sigma v\colon D\cdot m) = (\Sigma v\colon D\cdot n{\times}m)$

Antidistributive — if $D{\neq}null$
and v does not appear in a

$a \Leftarrow \exists v\colon D\cdot b = \forall v\colon D\cdot a \Leftarrow b$
$a \Leftarrow \forall v\colon D\cdot b = \exists v\colon D\cdot a \Leftarrow b$

Generalization — if element x: $\square p$

$p\, x \implies \exists p$

Splitting — for any fixed domain

$\forall v\cdot a \wedge b = (\forall v\cdot a) \wedge (\forall v\cdot b)$
$\exists v\cdot a \wedge b \implies (\exists v\cdot a) \wedge (\exists v\cdot b)$
$\forall v\cdot a \vee b \Longleftarrow (\forall v\cdot a) \vee (\forall v\cdot b)$
$\exists v\cdot a \vee b = (\exists v\cdot a) \vee (\exists v\cdot b)$
$\forall v\cdot a \Rightarrow b \implies (\forall v\cdot a) \Rightarrow (\forall v\cdot b)$
$\forall v\cdot a \Rightarrow b \implies (\exists v\cdot a) \Rightarrow (\exists v\cdot b)$
$\forall v\cdot a = b \implies (\forall v\cdot a) = (\forall v\cdot b)$
$\forall v\cdot a = b \implies (\exists v\cdot a) = (\exists v\cdot b)$

Commutative

$\forall v\cdot \forall w\cdot b = \forall w\cdot \forall v\cdot b$
$\exists v\cdot \exists w\cdot b = \exists w\cdot \exists v\cdot b$

Semicommutative (Skolem)

$\exists v\cdot \forall w\cdot b \implies \forall w\cdot \exists v\cdot b$
$\forall x\cdot \exists y\cdot p\, x\, y = \exists f\cdot \forall x\cdot p\, x\, (f\, x)$

Domain Change

$A\colon B \implies (\forall v\colon A\cdot b) \Leftarrow (\forall v\colon B\cdot b)$
$A\colon B \implies (\exists v\colon A\cdot b) \Rightarrow (\exists v\colon B\cdot b)$
$\forall v\colon A\cdot v\colon B \Rightarrow p = \forall v\colon A\text{`}B\cdot p$
$\exists v\colon A\cdot v\colon B \wedge p = \exists v\colon A\text{`}B\cdot p$

Extreme

$(\Downarrow n\colon int\cdot n) = (\Downarrow n\colon real\cdot n) = -\infty$
$(\Uparrow n\colon int\cdot n) = (\Uparrow n\colon real\cdot n) = \infty$

Connection (Galois)

$n{\leq}m = \forall k\cdot k{\leq}n \Rightarrow k{\leq}m$
$n{\leq}m = \forall k\cdot k{<}n \Rightarrow k{<}m$
$n{\leq}m = \forall k\cdot m{\leq}k \Rightarrow n{\leq}k$
$n{\leq}m = \forall k\cdot m{<}k \Rightarrow n{<}k$

$n{\downarrow}(\Downarrow v\colon D\cdot m) = (\Downarrow v\colon D\cdot n{\downarrow}m)$
$n{\downarrow}(\Uparrow v\colon D\cdot m) = (\Uparrow v\colon D\cdot n{\downarrow}m)$
$n + (\Downarrow v\colon D\cdot m) = (\Downarrow v\colon D\cdot n{+}m)$
$n - (\Downarrow v\colon D\cdot m) = (\Uparrow v\colon D\cdot n{-}m)$
$(\Downarrow v\colon D\cdot m) - n = (\Downarrow v\colon D\cdot m{-}n)$
$n{\geq}0 \Rightarrow n{\times}(\Downarrow v\colon D\cdot m) = (\Downarrow v\colon D\cdot n{\times}m)$
$n{\leq}0 \Rightarrow n{\times}(\Downarrow v\colon D\cdot m) = (\Uparrow v\colon D\cdot n{\times}m)$
$(\Pi v\colon D\cdot m)^n = (\Pi v\colon D\cdot m^n)$

End of Quantifiers

11.3.9 Limits

$(\Uparrow m{\cdot}\ \Downarrow n{\cdot}\ f\,(m{+}n)) \ \le\ \Updownarrow f \ \le\ (\Downarrow m{\cdot}\ \Uparrow n{\cdot}\ f\,(m{+}n))$
$\exists m{\cdot}\ \forall n{\cdot}\ p\,(m{+}n) \ \Longrightarrow\ \Updownarrow p \ \Longrightarrow\ \forall m{\cdot}\ \exists n{\cdot}\ p\,(m{+}n)$
$\Updownarrow n{\cdot}\ n \ =\ \infty$

End of Limits

11.3.10 Specifications and Programs

For specifications P, Q, R, and S, and binary b,

$ok \ =\ x'{=}x \ \wedge\ y'{=}y \ \wedge\ ...$
$x{:=}\,e \ =\ x'{=}e \ \wedge\ y'{=}y \ \wedge\ ...$
$P.\,Q \ =\ \exists x'', y'', ...{\cdot}\ \langle x', y', ...{\rightarrow}P\rangle\ x''\ y''\ ...\ \wedge\ \langle x, y, ...{\rightarrow}Q\rangle\ x''\ y''\ ...$
$P \| Q \ =\ \exists t_P, t_Q{\cdot}\ \langle t'{\rightarrow}P\rangle\ t_P \ \wedge\ \langle t'{\rightarrow}Q\rangle\ t_Q \ \wedge\ t' = t_P \uparrow t_Q$
if b **then** P **else** Q **fi** $\ =\ b \wedge P \vee \neg b \wedge Q \ =\ (b \Rightarrow P) \wedge (\neg b \Rightarrow Q)$
var $x{:}\ T{\cdot}\ P \ =\ \exists x, x'{:}\ T{\cdot}\ P$
frame $x{\cdot}\ P \ =\ P \ \wedge\ y'{=}y \ \wedge\ ...$
while b **do** P **od** $\ =\ t'{\ge}t \wedge$ **if** b **then** $P.\ t{:=}\,t{+}1.$ **while** b **do** P **od else** ok **fi**
$\forall \sigma, \sigma'{\cdot}$ **if** b **then** $P.\ W$ **else** ok **fi** $\ \Leftarrow\ W \ \Longrightarrow\ \forall \sigma, \sigma'{\cdot}$ **while** b **do** P **od** $\ \Leftarrow\ W$

To prove $\quad F\,m \ \Leftarrow$ **for** $i{:=}\,m;..n$ **do** P **od**
prove $\quad F\,i \ \Leftarrow\ i{:}\ m,..n \ \wedge\ (P.\ F(i{+}1))$
and $\quad F\,n \ \Leftarrow\ ok$

$A\,m \Rightarrow A'n \ \Leftarrow$ **for** $i{:=}\,m;..n$ **do** $i{:}\ m,..n \wedge A\,i \ \Rightarrow\ A'(i{+}1)$ **od**
wait until $w \ =\ t{:=}\,t{\uparrow}w$
assert $b \ =$ **if** b **then** ok **else** *screen*! "error". **wait until** ∞ **fi**
ensure $b \ =\ b \wedge ok$
$P.\ (P$ **result** $e){=}e$ but do not double-prime or substitute in (P **result** e)
$c? \ =\ \boldsymbol{r}{:=}\,\boldsymbol{r}{+}1$
$c \ =\ \mathcal{M}c_{\boldsymbol{r}c-1}$
$c!\,e \ =\ \mathcal{M}c_{\boldsymbol{w}c} = e \ \wedge\ \mathcal{T}c_{\boldsymbol{w}c} = t \ \wedge\ (wc{:=}\,wc{+}1)$
$\surd c \ =\ \mathcal{T}c_{\boldsymbol{r}c}$ + (transit time) $\le t$
ivar $x{:}\ T{\cdot}\ S \ =\ \exists x{:}\ time{\rightarrow}T{\cdot}\ S$
chan $c{:}\ T{\cdot}\ P \ =\ \exists \mathcal{M}c{:}\ \infty{*}T{\cdot}\ \exists \mathcal{T}c{:}\ \infty{*}xreal{\cdot}\ \exists \boldsymbol{r}c, \boldsymbol{r}c', \boldsymbol{w}c, \boldsymbol{w}c'{:}\ xnat{\cdot}$
$\qquad (\forall i, j{:}\ nat{\cdot}\ i{\le}j \ \Rightarrow\ t \le \mathcal{T}c_i \le \mathcal{T}c_j \le t') \wedge \boldsymbol{r}c{=}\boldsymbol{w}c{=}0 \wedge P$

$ok.\,P \ =\ P.\,ok \ =\ P$	identity
$P.\,(Q.\,R) \ =\ (P.\,Q).\,R$	associativity
$P{\vee}Q.\,R{\vee}S \ =\ (P.\,R) \vee (P.\,S) \vee (Q.\,R) \vee (Q.\,\mathrm{S})$	distributivity
if b **then** P **else** Q **fi**$.\,R \ =$ **if** b **then** $P.\,R$ **else** $Q.\,R$ **fi**	distributivity (unprimed b)
$P.$ **if** b **then** Q **else** R **fi** $\ =$ **if** $P.\,b$ **then** $P.\,Q$ **else** $P.\,R$ **fi**	distributivity (unprimed b)
$P \,\|\, Q \ =\ Q \,\|\, P$	symmetry
$P \,\|\, (Q \,\|\, R) \ =\ (P \,\|\, Q) \,\|\, R$	associativity
$P \,\|\, t'{=}t \ =\ P \ =\ t'{=}t \,\|\, P$	identity
$P \,\|\, Q{\vee}R \ =\ (P \,\|\, Q) \vee (P \,\|\, R)$	distributivity
$P \,\|$ **if** b **then** Q **else** R **fi** $\ =$ **if** b **then** $P \,\|\, Q$ **else** $P \,\|\, R$ **fi**	distributivity
if b **then** $P \| Q$ **else** $R \| S$ **fi** $\ =$ **if** b **then** P **else** R **fi** $\,\|$ **if** b **then** Q **else** S **fi**	distributivity
$x{:=}$ **if** b **then** e **else** f **fi** $\ =$ **if** b **then** $x{:=}\,e$ **else** $x{:=}\,f$ **fi**	functional-imperative

End of Specifications and Programs

11.3.11 Substitution

Let x and y be different boundary state variables, let e and f be expressions of the prestate, and let P be a specification.

$x{:=}\,e\,.\,P \;=\;$ (for x substitute e in P)

$(x{:=}\,e \parallel y{:=}\,f).\;P \;=\;$ (for x substitute e and concurrently for y substitute f in P)

End of Substitution

11.3.12 Assertions

Let P and Q be specifications. Let A be an assertion and let A' be the same as A but with primes on all the variables.

$A \wedge (P.\,Q) \;\Leftarrow\; A \wedge P.\,Q$

$A \Rightarrow (P.Q) \;\Leftarrow\; A \Rightarrow P.\,Q$

$(P.Q) \wedge A' \;\Leftarrow\; P.\,Q \wedge A'$

$(P.Q) \Leftarrow A' \;\Leftarrow\; P.\,Q \Leftarrow A'$

$P.\,A \wedge Q \;\Leftarrow\; P \wedge A'.\,Q$

$P.\,Q \;\Leftarrow\; P \wedge A'.\;A \Rightarrow Q$

A is a sufficient precondition for P to be refined by S
if and only if $A \Rightarrow P$ is refined by S.

A' is a sufficient postcondition for P to be refined by S
if and only if $A' \Rightarrow P$ is refined by S.

End of Assertions

11.3.13 Refinement

Refinement by Steps (Stepwise Refinement) (monotonicity, transitivity)

If $A \Leftarrow$ **if** b **then** C **else** D **fi** and $C \Leftarrow E$ and $D \Leftarrow F$ are theorems,
then $A \Leftarrow$ **if** b **then** E **else** F **fi** is a theorem.

If $A \Leftarrow B.C$ and $B \Leftarrow D$ and $C \Leftarrow E$ are theorems, then $A \Leftarrow D.E$ is a theorem.

If $A \Leftarrow B \| C$ and $B \Leftarrow D$ and $C \Leftarrow E$ are theorems, then $A \Leftarrow D \| E$ is a theorem.

If $A \Leftarrow B$ and $B \Leftarrow C$ are theorems, then $A \Leftarrow C$ is a theorem.

Refinement by Parts (monotonicity, conflation)

If $A \Leftarrow$ **if** b **then** C **else** D **fi** and $E \Leftarrow$ **if** b **then** F **else** G **fi** are theorems,
then $A \wedge E \Leftarrow$ **if** b **then** $C \wedge F$ **else** $D \wedge G$ **fi** is a theorem.

If $A \Leftarrow B.C$ and $D \Leftarrow E.F$ are theorems, then $A \wedge D \Leftarrow B \wedge E.\,C \wedge F$ is a theorem.

If $A \Leftarrow B \| C$ and $D \Leftarrow E \| F$ are theorems, then $A \wedge D \Leftarrow B \wedge E \parallel C \wedge F$ is a theorem.

If $A \Leftarrow B$ and $C \Leftarrow D$ are theorems, then $A \wedge C \Leftarrow B \wedge D$ is a theorem.

Refinement by Cases

$P \Leftarrow$ **if** b **then** Q **else** R **fi** is a theorem if and only if
$P \Leftarrow b \wedge Q$ and $P \Leftarrow \neg b \wedge R$ are theorems.

End of Refinement

End of Laws

11.4 Names

abs: *xreal*→§*r*: *xreal*· *r*≥0	*abs r* = **if** *r*≥0 **then** *r* **else** –*r* **fi**
bin (the binary values)	*bin* = ⊤, ⊥
ceil: *real*→*int*	*r* ≤ *ceil r* < *r*+1
char (the characters)	*char* = ..., "a", "A", ...
div: *real*→(§*r*: *real*· *r*>0)→*int*	*div x y* = *floor* (*x*/*y*)
divides: (*nat*+1)→*int*→*bin*	*divides n i* = *i*/*n*: *int*
entro: *prob*→§*r*: *xreal*· *r*≥0	*entro p* = *p* × *info p* + (1–*p*) × *info* (1–*p*)
even: *int*→*bin*	*even i* = *i*/2: *int*
	even = *divides* 2
floor: *real*→*int*	*floor r* ≤ *r* < *floor r* + 1
info: *prob*→§*r*: *xreal*· *r*≥0	*info p* = – *log p*
int (the integers)	*int* = *nat*, –*nat*
log: (§*r*: *xreal*· *r*≥0)→*xreal*	*log* (2^x) = *x*
	log (*x*×*y*) = *log x* + *log y*
mod: *real*→(§*r*: *real*· *r*>0)→*real*	0 ≤ *mod a d* < *d*
	a = *div a d* × *d* + *mod a d*
nat (the naturals)	0, *nat*+1: *nat*
	0, *B*+1: *B* ⇒ *nat*: *B*
nil (the empty string)	↔*nil* = 0
	nil; *S* = *S* = *S*; *nil*
	nil ≤ *S*
null (the empty bunch)	¢*null* = 0
	null, *A* = *A* = *A*, *null*
	null: *A*
odd: *int*→*bin*	*odd i* = ¬ *i*/2: *int*
	odd = ¬*even*
ok (the empty program)	*ok* = σ′=σ
	ok.*P* = *P* = *P*.*ok*
prob (probability)	*prob* = §*r*: *real*· 0≤*r*≤1
rand (random number)	*rand n*: 0,..*n*
rat (the rationals)	*rat* = *int*/(*nat*+1)
real (the reals)	*r*: *real* = *r*: *xreal* ∧ –∞<*r*<∞
suc: *nat*→(*nat*+1)	*suc n* = *n*+1
xint (the extended integers)	*xint* = –∞, *int*, ∞
xnat (the extended naturals)	*xnat* = *nat*, ∞
xrat (the extended rationals)	*xrat* = –∞, *rat*, ∞
xreal (the extended reals)	*x*: *xreal* = ∃*f*: *nat*→*rat*· *x* = ⇳*f*

End of Names

11.5 Symbols

symbol	page	pronunciation
$\top$	3	true
$\bot$	3	false
$\neg$	3	not
$\wedge$	3	and
$\vee$	3	or
$\Rightarrow$	3	implies
$\Longrightarrow$	3	implies
$\Leftarrow$	3	follows from, is implied by
$\Longleftarrow$	3	follows from, is implied by
$=$	3	equals, if and only if
$\equiv$	3	equals, if and only if
$\neq$	3	differs from, is unequal to
$<$	13	less than
$>$	13	greater than
$\leq$	13	less than or equal to
$\geq$	13	greater than or equal to
$+$	12	plus
$-$	12	minus
$\times$	12	times, multiplication
$/$	12	divided by
$\uparrow$	12	maximum
$\downarrow$	12	minimum
,	14	bunch union
,..	16	union from (including) to (excluding)
‘	14	bunch intersection
;	17	string join
;;	20	list join
;..	19	join from (including) to (excluding)
:	14	is in, are in, bunch inclusion
::	92	includes
:=	36	assignment
∞	78	label, target of **go to**
.	36	sequential composition
·	26	quantifier abbreviation
!	136	output
?	136	input

symbol	page	pronunciation
$\surd$	136	input check
( )	4	parentheses for precedence
{ }	17	set brackets
[ ]	20	list brackets
$\langle\ \rangle$	23	function (scope) brackets
$\not$	17	power
$\not\in$	14	bunch size, cardinality
$\$$	17	set size, cardinality
$\leftrightarrow$	18	string size (length)
#	20,23	list size (length), function size
$\mid$	20,24	selective union, otherwise
$\parallel$	121	concurrent (parallel) composition
~	17,20	contents of a set or list
*	18	repetition of a string
$\Box$	20,23	domain of a list or function
$\rightarrow$	23	function arrow
$\in$	17	element of a set
$\subseteq$	17	subset
$\cup$	17	set union
$\cap$	17	set intersection
@	22	index with a pointer
$\forall$	26	for all, universal quantifier
$\exists$	26	there exists, existential quantifier
Σ	26	sum of, summation quantifier
Π	26	product of, product quantifier
$\Uparrow$	26	maximum (lub) quantifier
$\Downarrow$	26	minimum (glb) quantifier
$\Updownarrow$	33	limit quantifier
§	28	those, solution quantifier
′	34	x' is final value of state variable x
“ ”	13,19	“hi” is a text or string of characters
a^b	12	exponentiation
a_b	18	string indexing
$a\ b$	20,31	indexing, application, composition
$\triangleleft\ \triangleright$	18	string modification
∞	12	infinity

assert	79	**if then else fi**	4
chan	141	**ivar**	129
do od	73	**or**	80
ensure	80	**result**	81
exit when	73	**var**	68,83
for do od	76	**wait until**	79
frame	69	**while do od**	71
go to	78		

End of Symbols

11.6 Precedence

0 ⊤ ⊥ ( ) { } [ ] ⟨ ⟩ **if fi** **do od** number text name superscript subscript
1 @ adjacency
2 prefix– ¢ $ ↔ # * ~ ↯ □ → √ ∀ ∃ Σ Π ⇑ ⇓ ⇳ §
3 × / ∩ ↑ ↓
4 + infix– ∪
5 ; ;.. ;; ‘
6 , ,.. | ◅ ▻
7 = ≠ < > ≤ ≥ : :: ∈ ⊆
8 ¬
9 ∧
10 ∨
11 ⇒ ⇐
12 := ! ?
13 **exit when** **go to** **wait until** **assert** **ensure** **or**
14 . ‖ **result**
15 ∀· ∃· Σ· Π· ⇑· ⇓· ⇳· §· **var·** **ivar·** **chan·** **frame·**
16 ⩵ ⟹ ⟸

Superscripting and subscripting associate from right to left, and bracket what is in them.

Adjacency associates from left to right, so $a\ b\ c$ means the same as $(a\ b)\ c$. The infix operators @ / – associate from left to right. The infix operators * → associate from right to left. The infix operators × ∩ + ∪ ; ;; ‘ , | ∧ ∨ . ‖ are associative (they associate in both directions).

Quantifiers as prefix operators are on level 2, but in the abbreviated quantifier notation (with ·) they are on level 15.

On levels 7, 11, and 16 the operators are continuing. For example, $a=b=c$ neither associates to the left nor associates to the right, but means the same as $a=b \wedge b=c$. On any one of these levels, a mixture of continuing operators can be used. For example, $a \le b < c$ means the same as $a \le b \wedge b < c$.

The operators ⩵ ⟹ ⟸ are identical to = ⇒ ⇐ except for precedence.

End of Precedence

11.7 Distribution

The operators in the following expressions distribute over bunch union in any operand:

$\neg A \quad A \wedge B \quad A \vee B \quad -A \quad A+B \quad A-B \quad A \times B \quad A/B \quad A^B \quad A \uparrow B \quad A \downarrow B \quad A,B$
$A‘B \quad \$A \quad A \cup B \quad A \cap B \quad \sim A \quad A;B \quad \leftrightarrow A \quad A_B \quad [A] \quad A;;B \quad A\ B \quad \#A \quad A@B$

The operator in $A*B$ distributes over bunch union in its left operand only.

End of Distribution

End of Reference

End of a Practical Theory of Programming